

Alloy & Applications



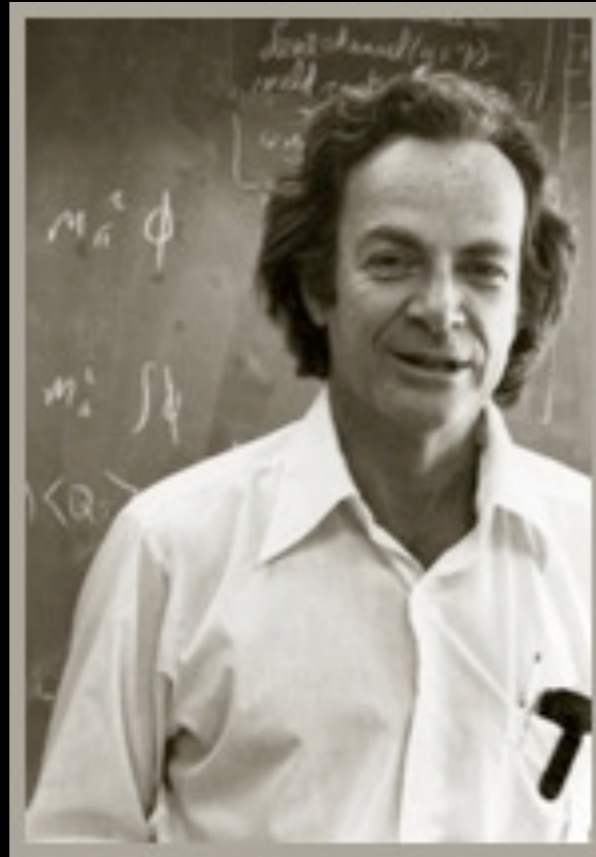
University of Wisconsin, Madison · November 4, 2009

Daniel Jackson, MIT

a puzzle



why doesn't pencil & paper design work?



The first principle is that you must not fool yourself, and you are the easiest person to fool.

Richard P. Feynman

early formalization

early formalization

formalize design

early formalization

formalize design

Z, VDM, B, ...

early formalization

formalize design

Z, VDM, B, ...

exercise design

early formalization

formalize design

Z, VDM, B, ...

exercise design

write proofs

early formalization

formal methods

formalize design

Z, VDM, B, ...

exercise design

write proofs

early formalization

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Z, VDM, B, ...

not code

exercise design

write proofs

very painful

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Java, ML, Ruby, C, ...

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formal methods

formalize design

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agile methods

formalize design

Java, ML, Ruby, C, ...

exercise design

early formalization

formal methods

formalize design

Z, VDM, B, ...

not code

exercise design

write proofs

very painful

agile methods

formalize design

Java, ML, Ruby, C, ...

exercise design

write tests

early formalization

formal methods

formalize design

Z, VDM, B, ...

not code

exercise design

write proofs

very painful

agile methods

formalize design

Java, ML, Ruby, C, ...

verbose

exercise design

write tests

painful

poor coverage

a dream?

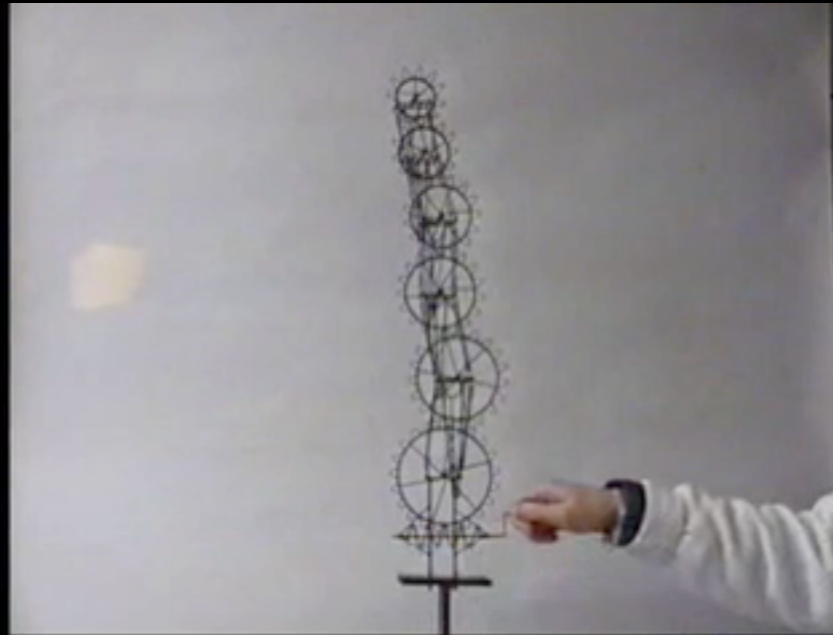
succinct, expressive models
animation & property checking
high coverage without proofs
no unit tests to write
code synthesized automatically

a heresy

Small Tower of 6 Gears, Arthur Ganson

a heresy

what matters most?



Small Tower of 6 Gears, Arthur Ganson

a heresy

what matters most?
finding subtle bugs?



Small Tower of 6 Gears, Arthur Ganson

a heresy

what matters most?
finding subtle bugs?
instant feedback



Small Tower of 6 Gears, Arthur Ganson

a heresy

what matters most?
finding subtle bugs?
instant feedback



Small Tower of 6 Gears, Arthur Ganson

transatlantic alloy

motivation

structural description + automation

transatlantic alloy

motivation

structural description + automation



Oxford, home of Z

transatlantic alloy

motivation

structural description + automation



Oxford, home of Z



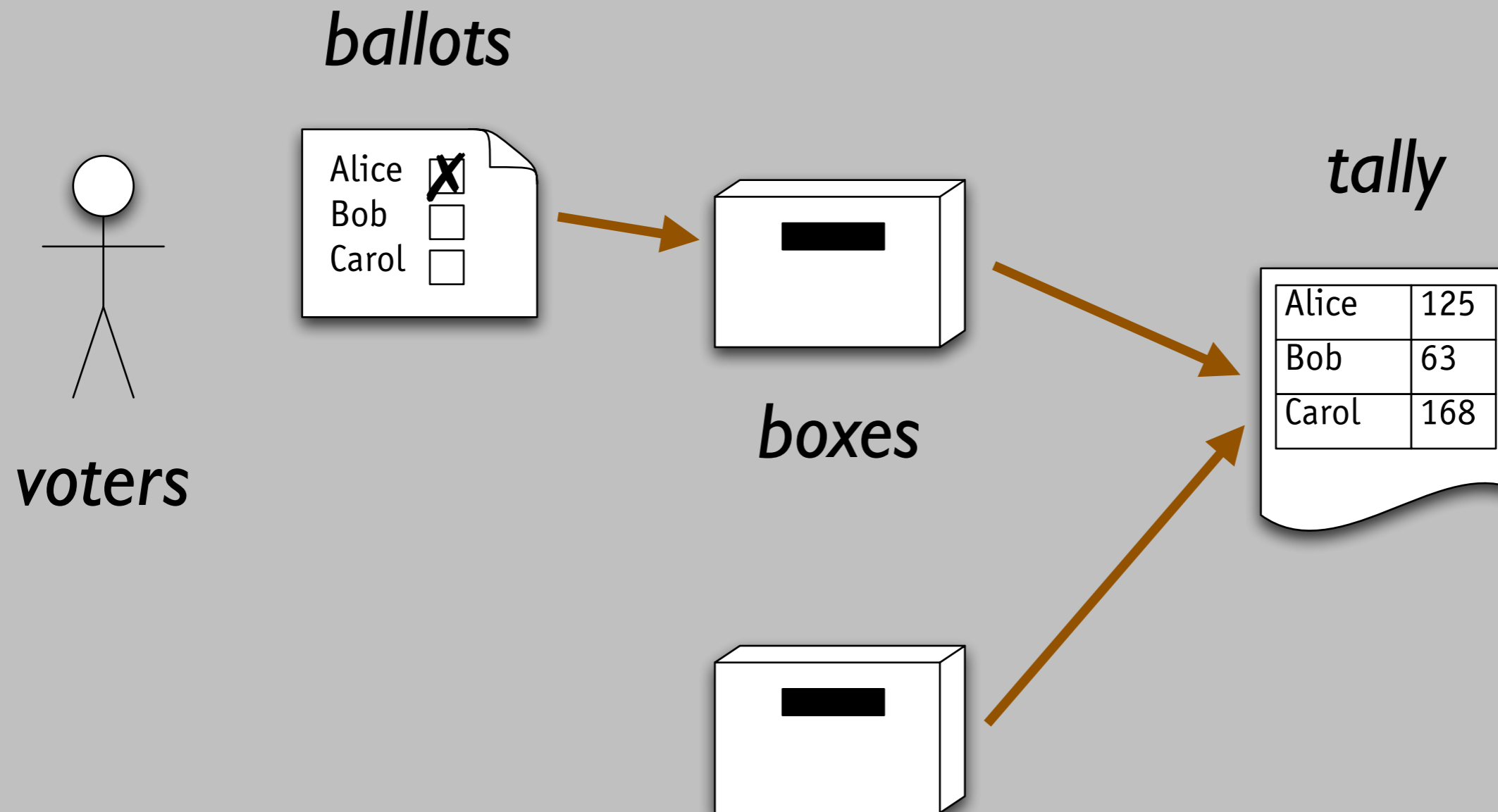
Pittsburgh, home of SMV

how we got here

<i>version</i>	<i>language</i>	<i>analysis</i>	<i>sample case study</i>
Nitpick (1995)	relational calculus subset of Z	relation enumeration	IPv6 routing
Alloy 1 (1999)	+ navigation exps quantifiers	WalkSAT, DP	intentional naming
Alloy 2 (2001)	+ relational ops higher arity	Chaff, Berkmin symmetry, sharing	key management, Unison filesync
Alloy 3 (2004)	+ subtyping, overloading	atomization (bad)	Mondex electronic purse
Alloy 4 (2007)	+ meta, sequences, arithmetic	sparse matrices better sharing	flash file systems, code analyses

example

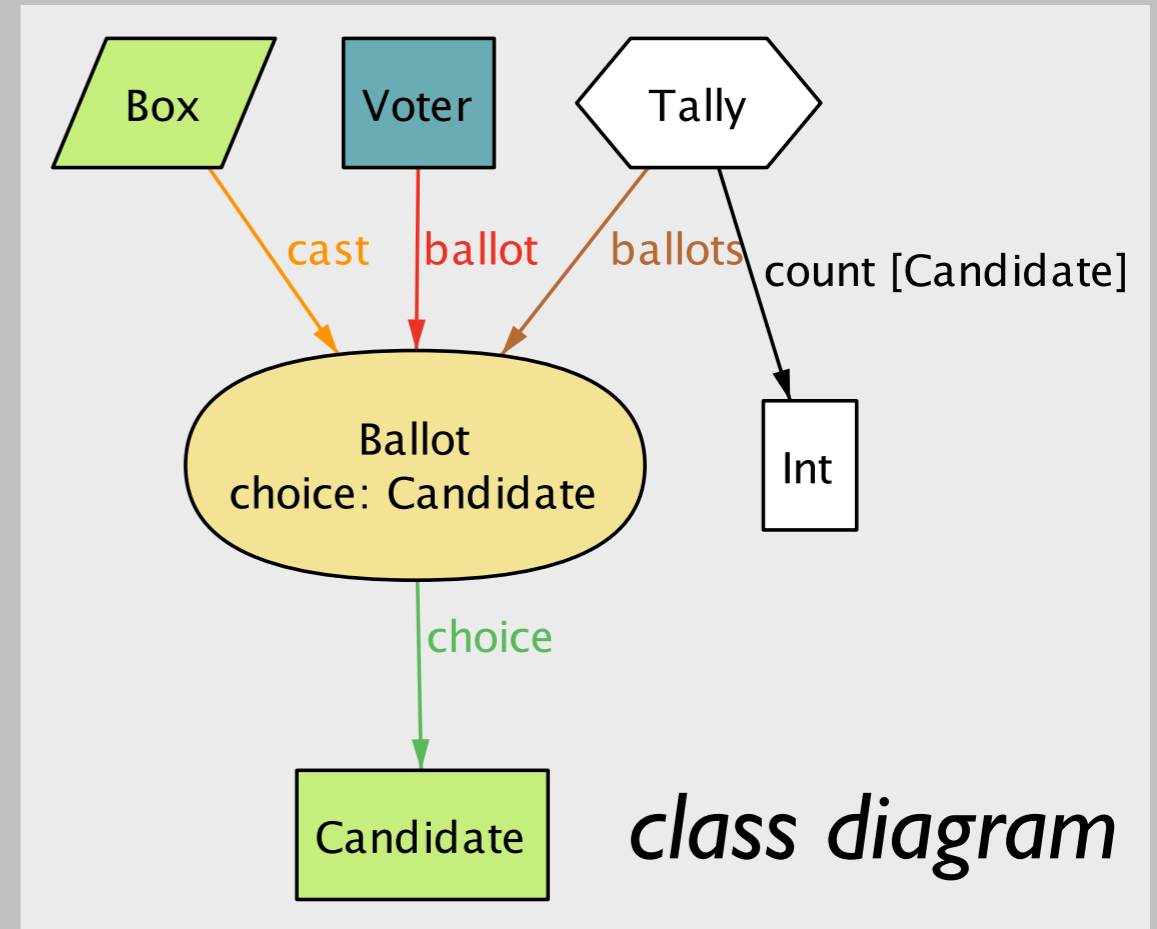
simplified voting



tally should correspond to voter's choices
what properties establish this requirement?

basic structure

```
sig Candidate {}  
sig Voter { ballot: lone Ballot }  
sig Ballot { choice: lone Candidate}  
sig Box { cast: set Ballot }  
one sig Tally {  
  ballots: set Ballot,  
  count: Candidate -> one Int  
} {  
all c: Candidate | count[c] = #(ballots & choice.c)  
}
```



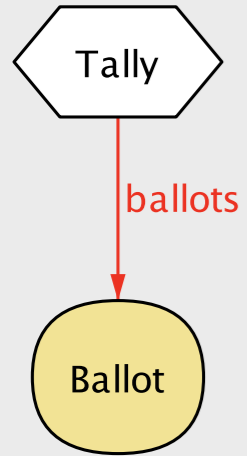
animation

animation

```
run {}
```

animation

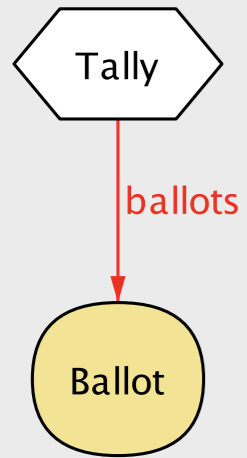
instance



`run {}`

animation

instance

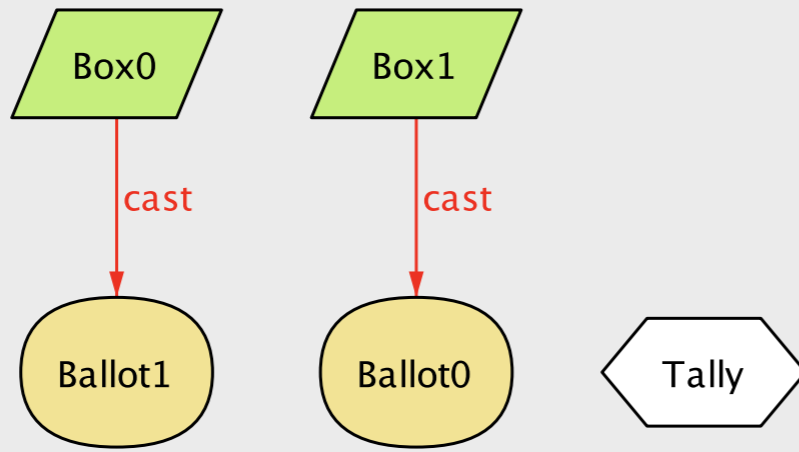
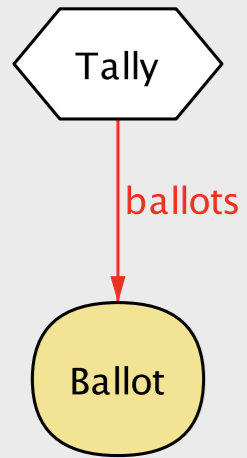


`run { }`

`run { some cast }`

animation

instance

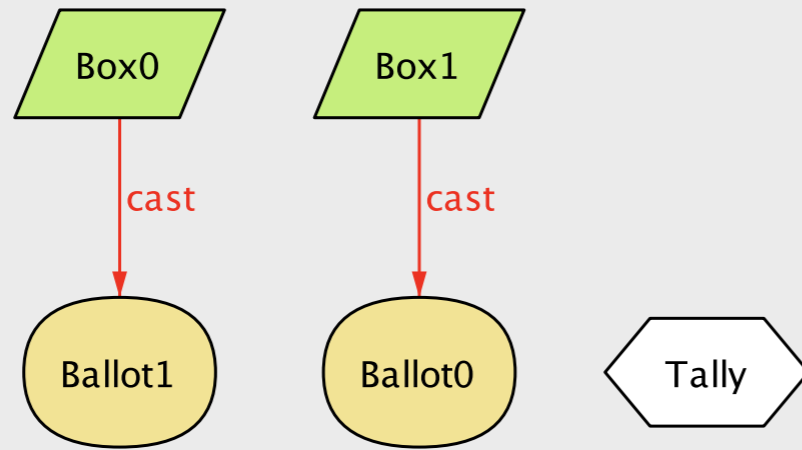
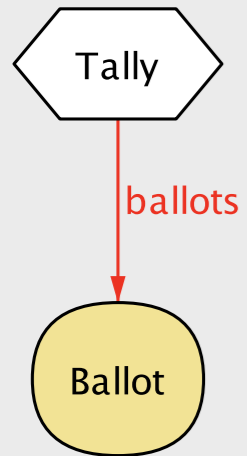


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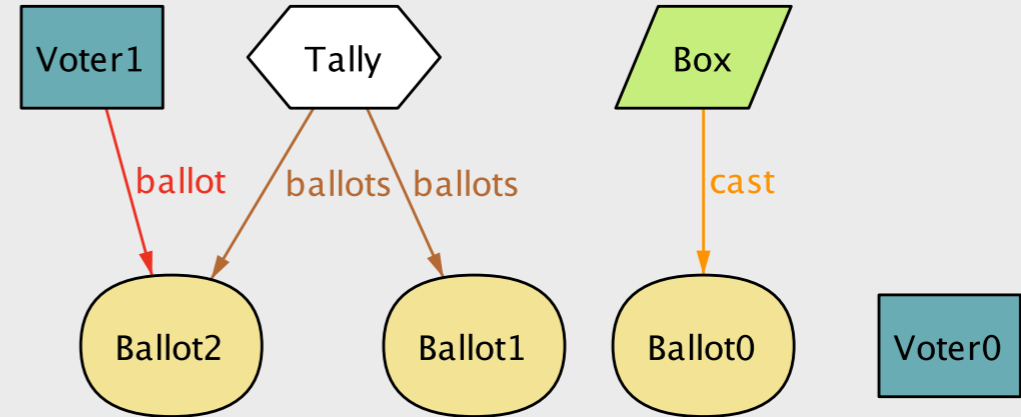
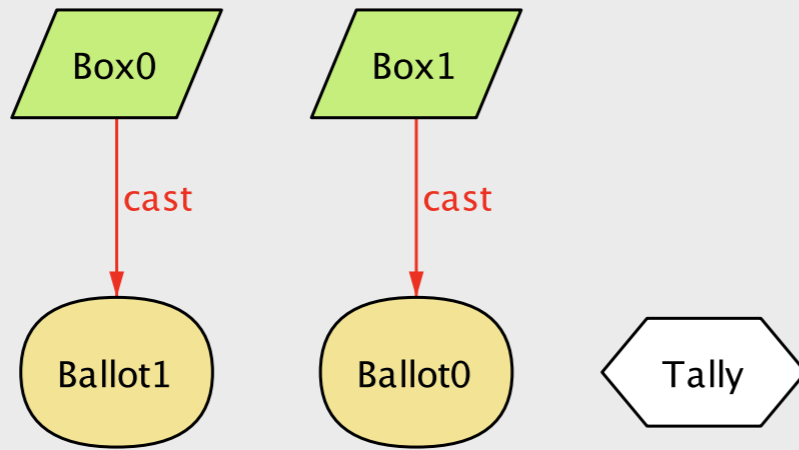
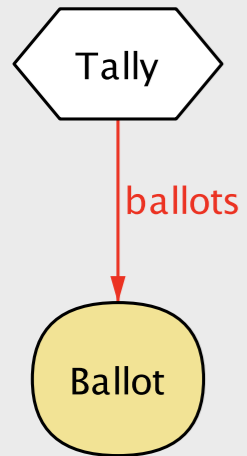
run { }

run { **some** cast }

... show next

animation

instance



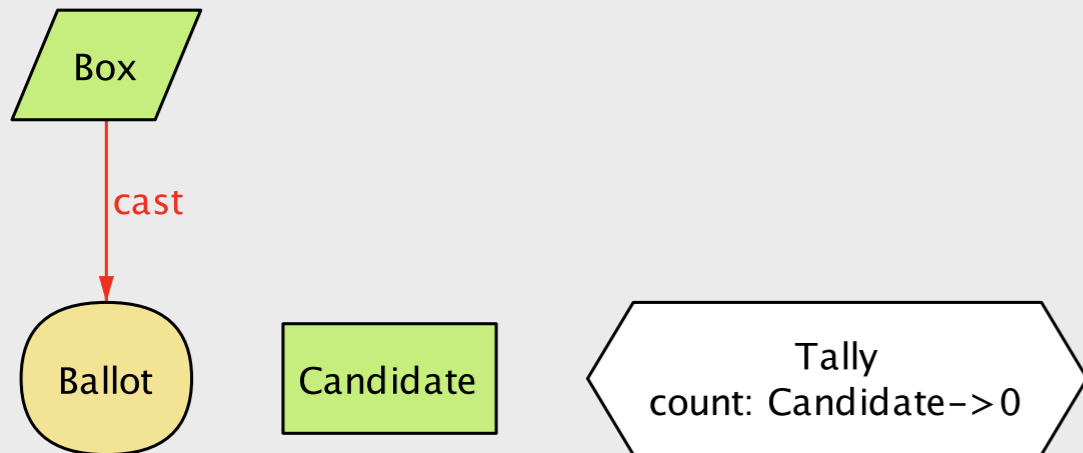
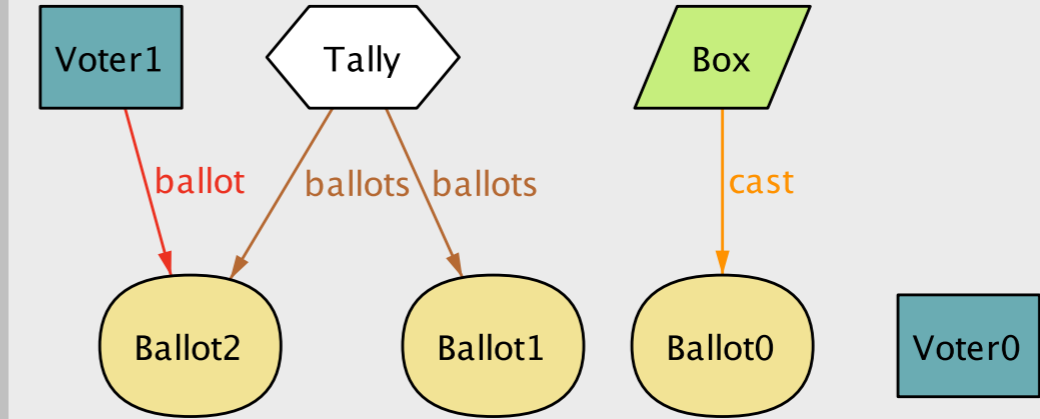
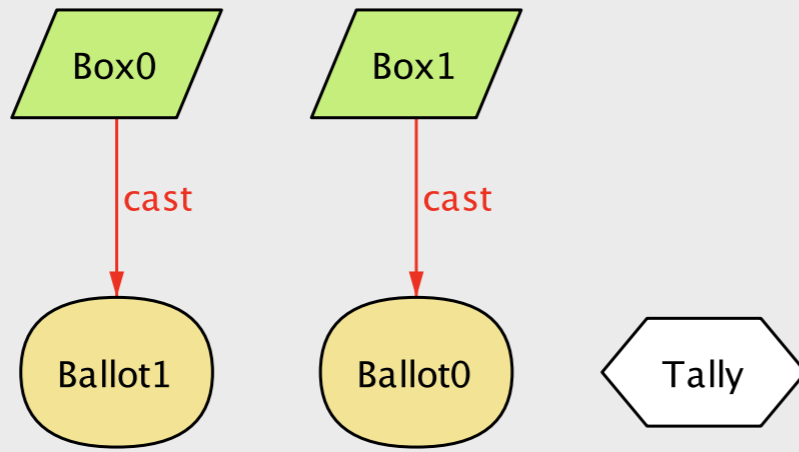
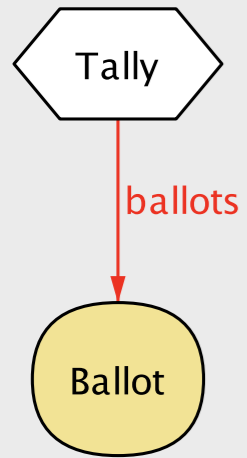
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animation

instance



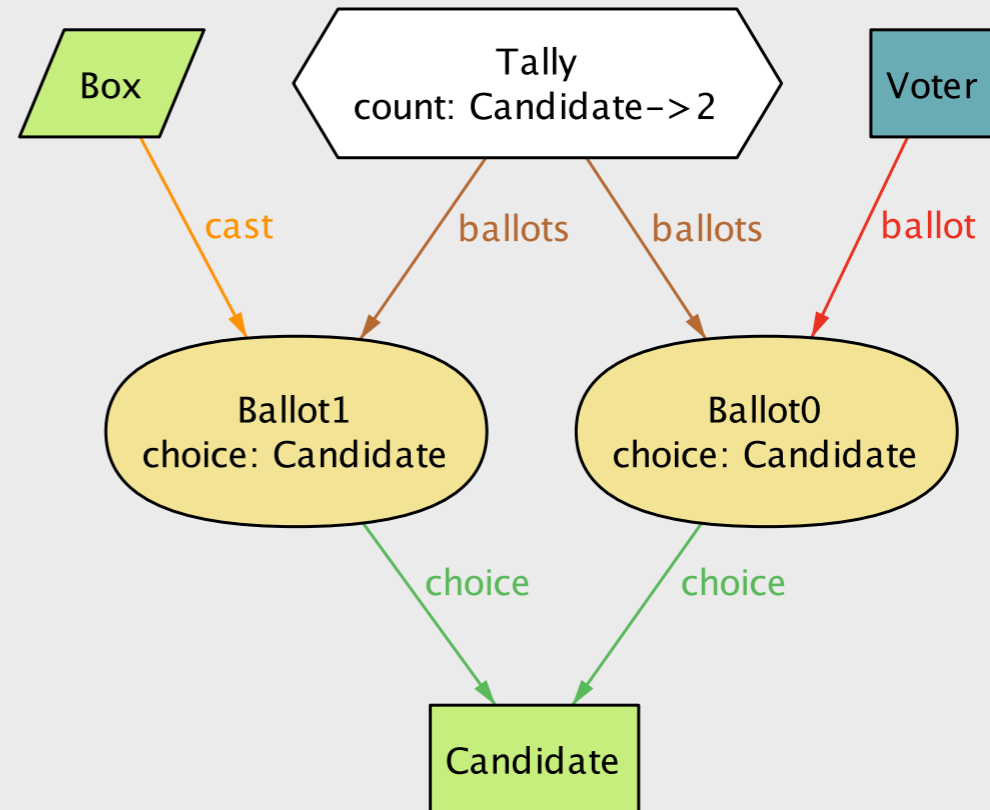
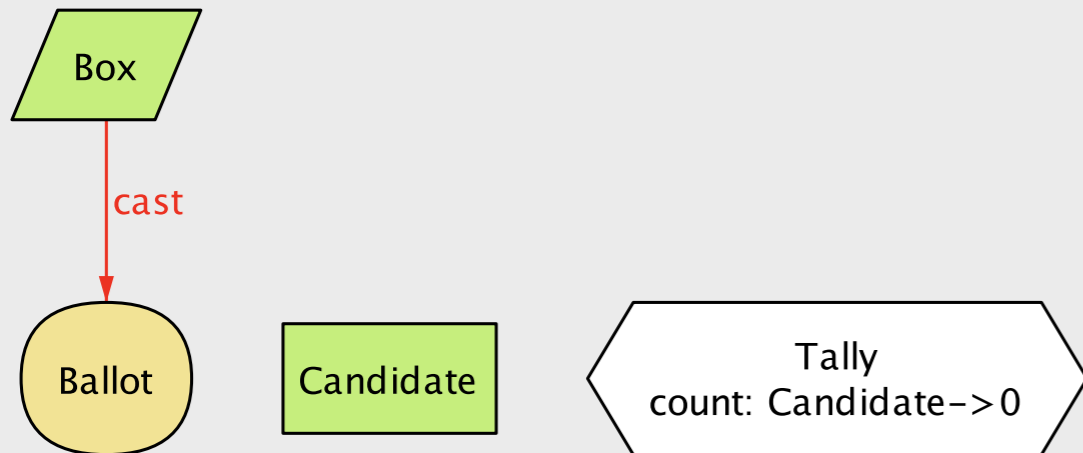
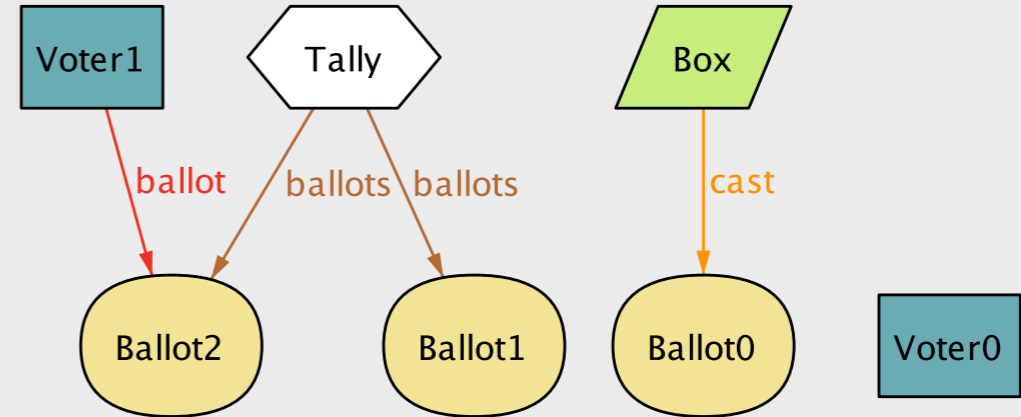
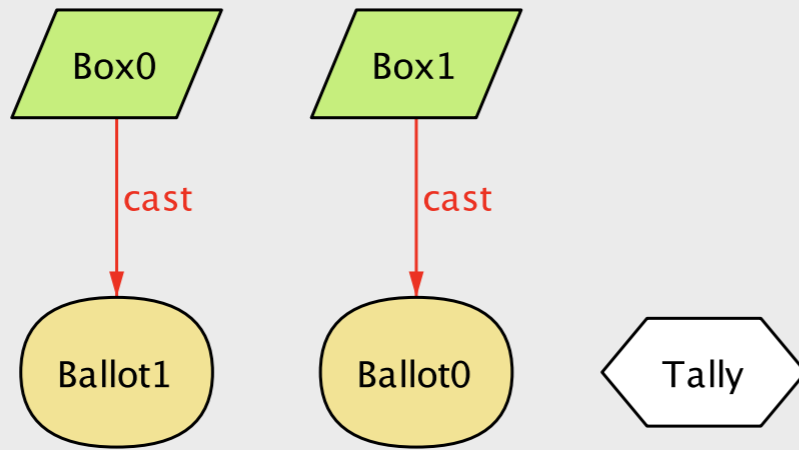
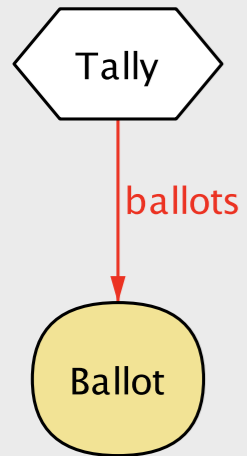
`run { }`

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... show next

animation

instance



`run { }`

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... show next

adding facts

sig Candidate { }

sig Voter { ballot: **lone** Ballot }

sig Ballot { choice: **lone** Candidate }

sig Box { cast: **set** Ballot }

one sig Tally { ballots: **set** Ballot, count: Candidate -> **one** Int }

{ **all** c: Candidate | count[c] = #(ballots & choice.c) }

fact AllVotersCastBallots { **all** v: Voter | **some** x: Box | v.ballot **in** x.cast }

fact NoBoxStuffing { **all** x: Box, b: x.cast | **some** v: Voter | v.ballot = b }

fact NoTallyStuffing { **all** b: Tally.ballots | **some** x: Box | b **in** x.cast }

fact NoBallotsLost { **all** x: Box | x.cast **in** Tally.ballots }

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run { }

adding facts

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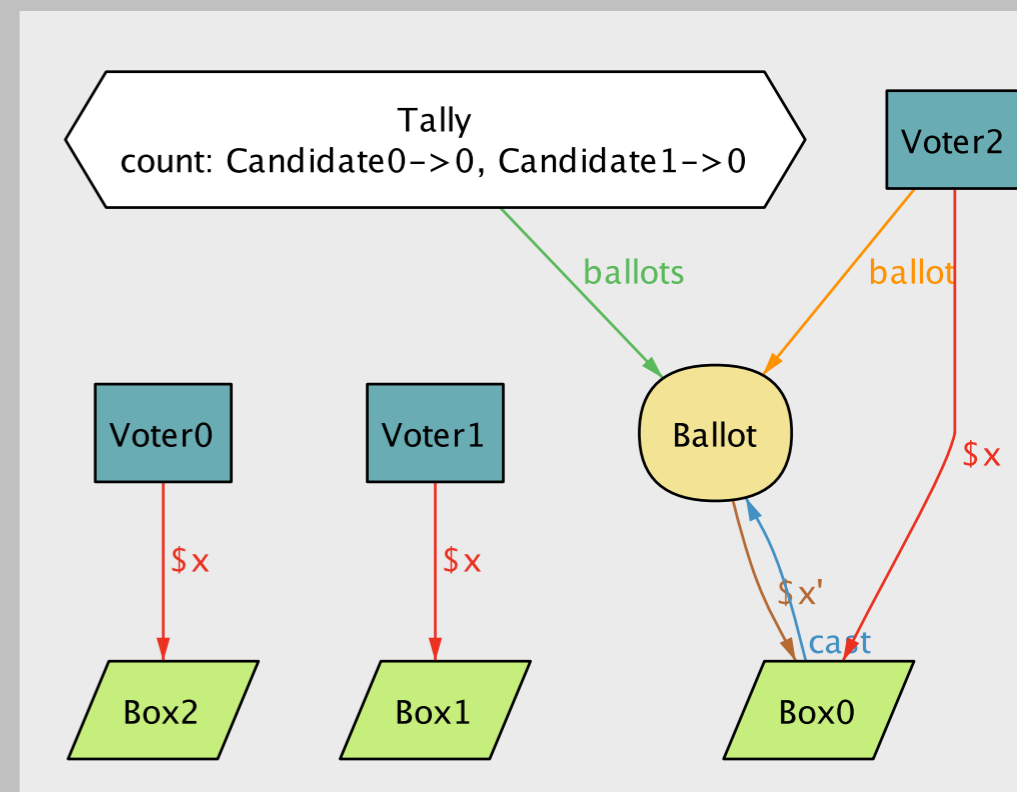
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fact AllVotersCastBallots {  
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```

```
run { }
```



checking an assertion

```
sig Candidate { }
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  ballots: set Ballot,
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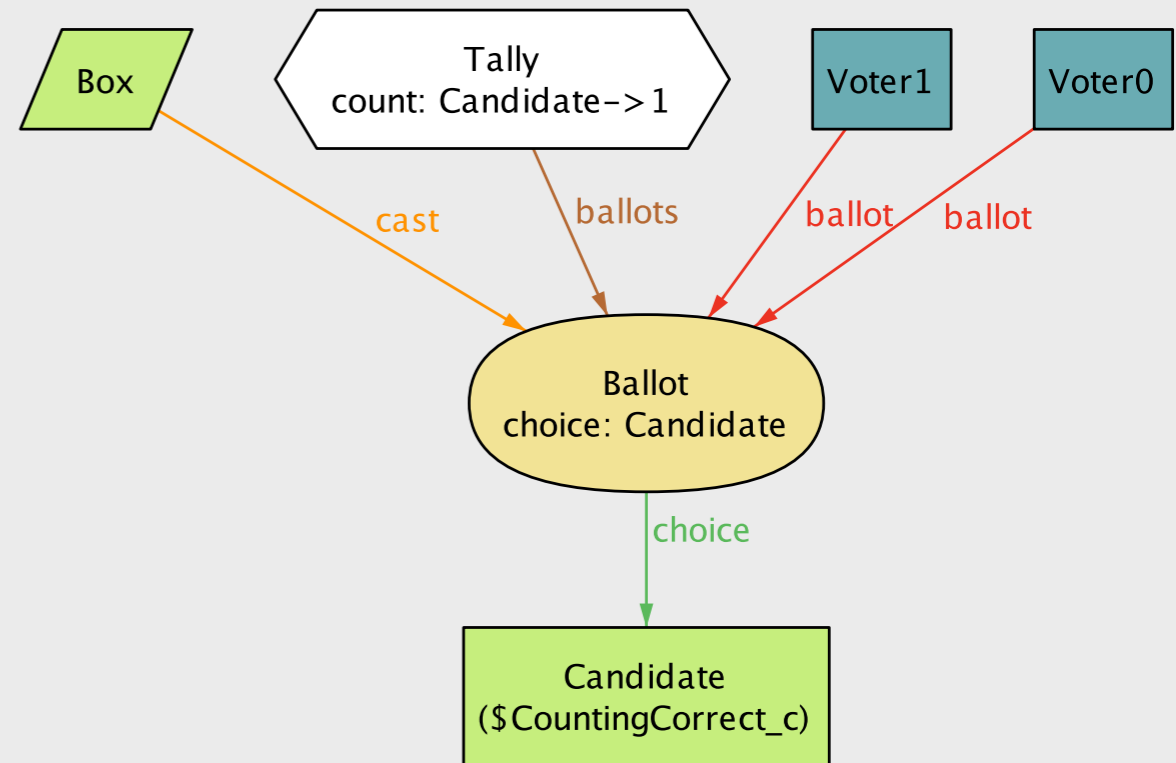
```
fact NoBallotsLost { all x: Box | x.cast in Tally.ballots }
```

```
pred CountingCorrect { all c: Candidate | Tally.count[c] = #ballot.choice.c }
```

```
check CountingCorrect
```


checking an assertion

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fact NoBallotsLost { all x: Box | x.cast in Tally.ballots }
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check CountingCorrect
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fixing & rechecking

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sig Candidate { }
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```

```
fact NoBallotsLost { all x: Box | x.cast in Tally.ballots }
```

```
fact NoSharedBallots { all b: Ballot | lone ballot.b }
```

```
check CountingCorrect for 5
```

fixing & rechecking

```
sig Car
sig Voter
sig Ballot
sig Box

Executing "Check CountingCorrect for 5"
Solver=minisat(jni) Bitwidth=4 MaxSeq=5 SkolemDepth=2 Symmetry=20
3693 vars. 235 primary vars. 8238 clauses. 180ms.
No counterexample found. Assertion may be valid. 92ms.
```

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one sig Tally {
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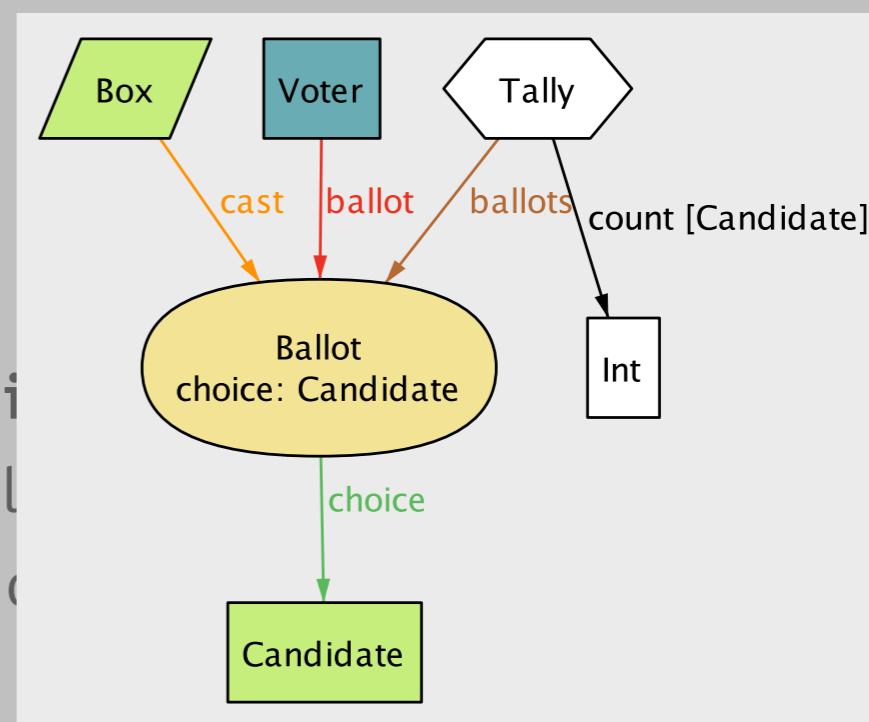
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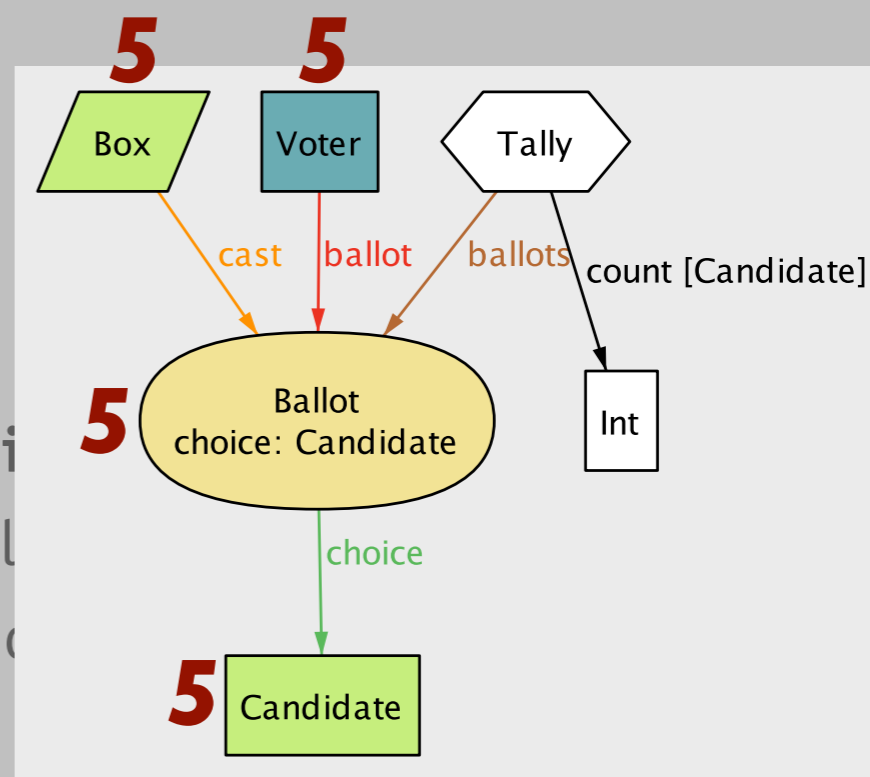
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fixing & rechecking

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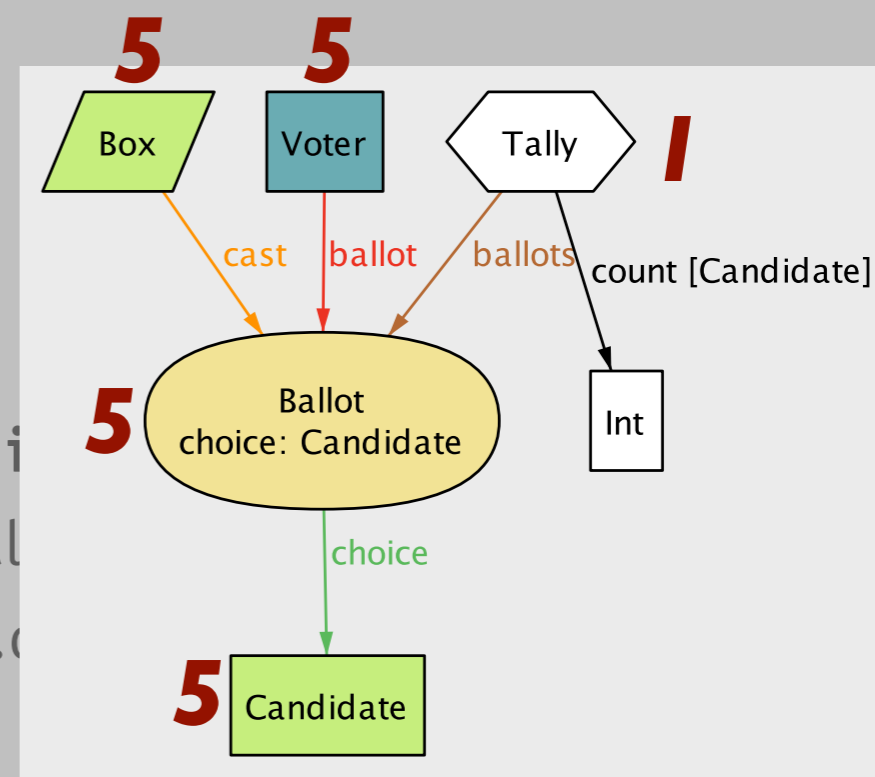
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fixing & rechecking

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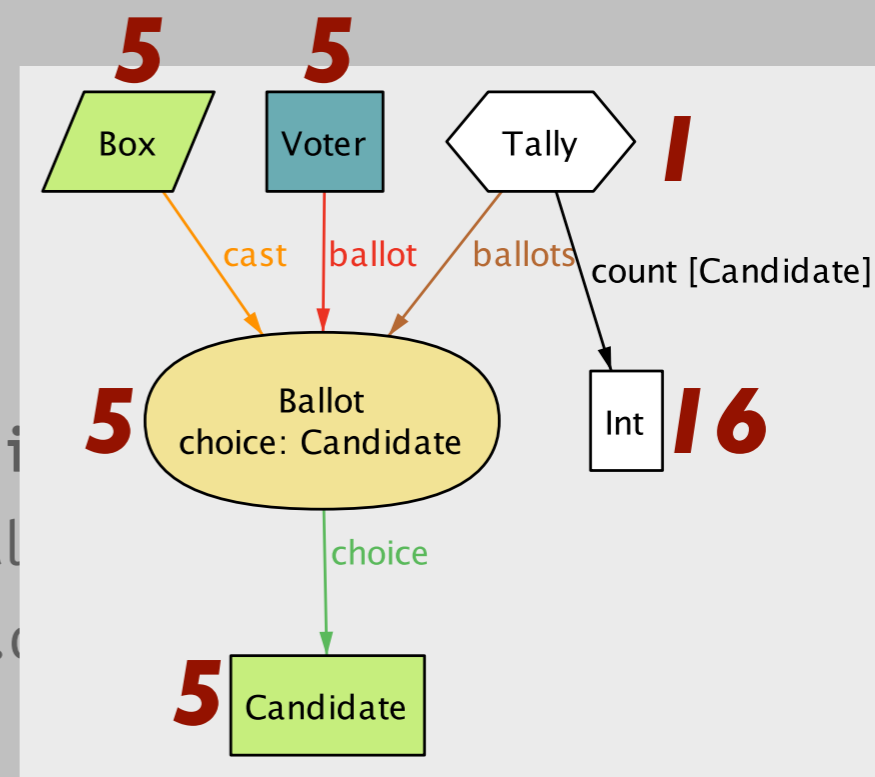
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{
all c: Candidate | count[c] = #(ballots & choice.c)
}
```

```
fact AllVotersCastBallots { all v: Voter | some x: Box | v.ballot in x.cast }
```

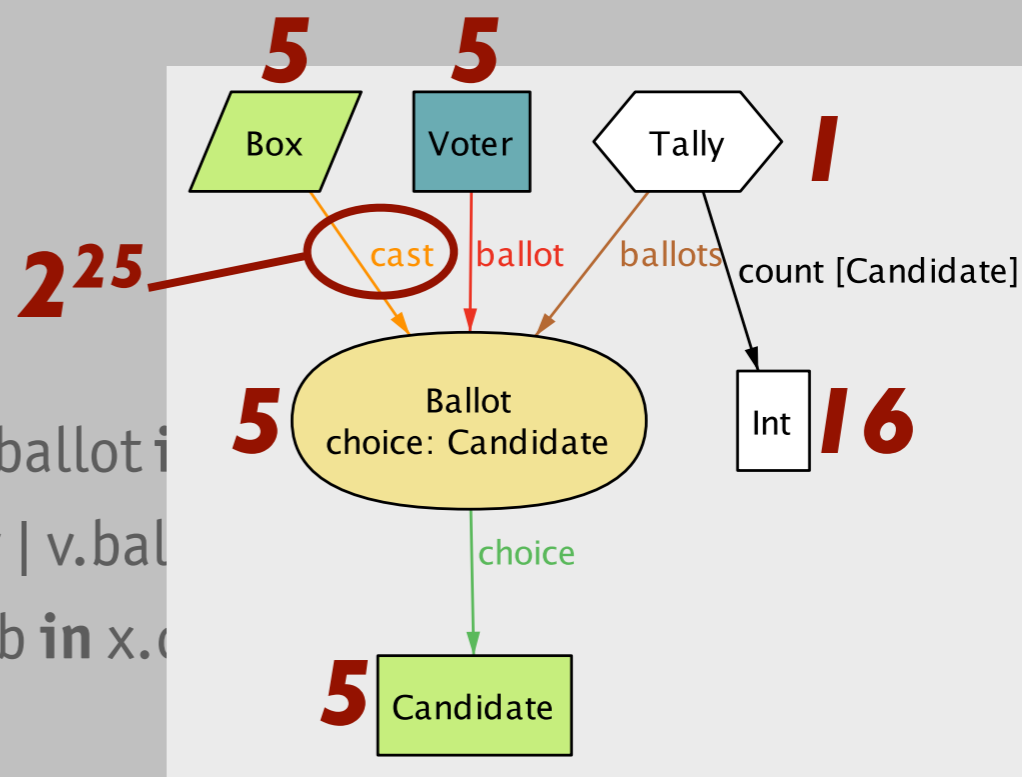
```
fact NoBoxStuffing { all x: Box, b: x.cast | some v: Voter | v.ballot = b }
```

```
fact NoTallyStuffing { all b: Tally.ballots | some x: Box | b in x.cast }
```

```
fact NoBallotsLost { all x: Box | x.cast in Tally.ballots }
```

```
fact NoSharedBallots { all b: Ballot | lone ballot.b }
```

```
check CountingCorrect for 5
```



fixing & rechecking

```
sig Car
sig Voter
sig Ballot
sig Box

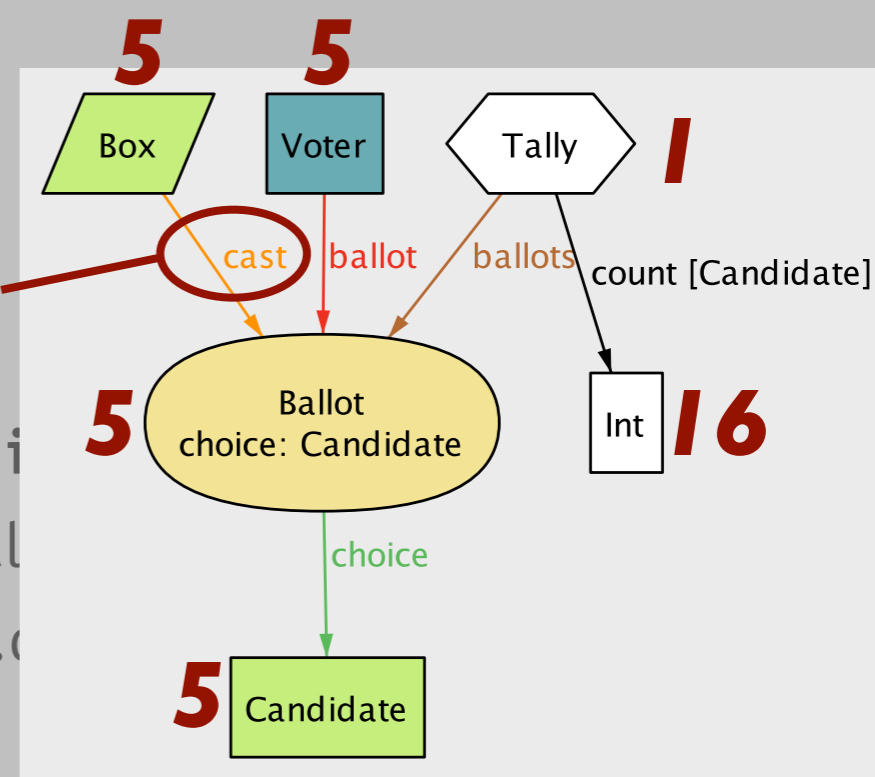
Executing "Check CountingCorrect for 5"
Solver=minisat(jni) Bitwidth=4 MaxSeq=5 SkolemDepth=2 Symmetry=20
3693 vars. 235 primary vars. 8238 clauses. 180ms.
No counterexample found. Assertion may be valid. 92ms.
```

```
one sig Tally {
  ballots: set Ballot,
  count: Candidate -> one Int }
{
all c: Candidate | count[c] = #(ballots & choice.c)
}
```

total space: 2^{235}

2^{25}

```
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fact NoBoxStuffing { all x: Box, b: x.cast | some v: Voter | v.ballot = b }
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```



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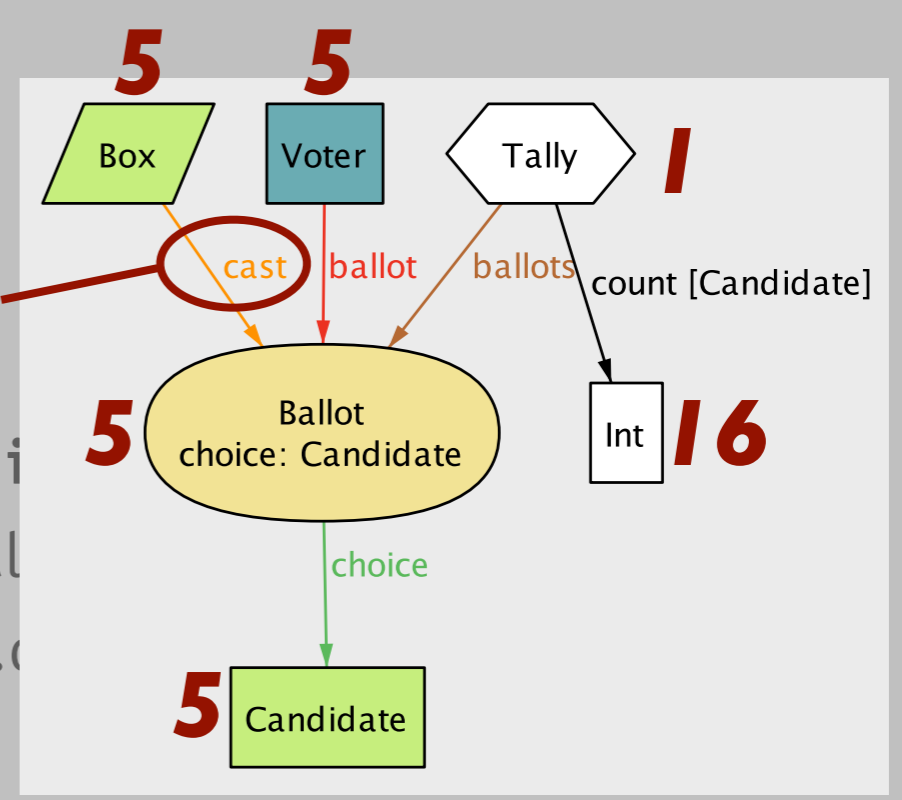
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```

total space: 2^{235}

2^{25}



fact NoSharedBallots { all b: Ballot | lone ballot.b }

check CountingCorrect for 5

what makes it work?

five ideas

analysis ideas

analysis as solving
small scope analysis
solving by SAT

language ideas

generalizing dot
funky idioms

idea: analysis = solving

idea: analysis = solving

does system S have property P ?

$$S(\text{beh}) \wedge \neg P(\text{beh})$$

idea: analysis = solving

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$$op(s, s') \wedge op(s, s'') \wedge s' \neq s''$$

idea: analysis = solving

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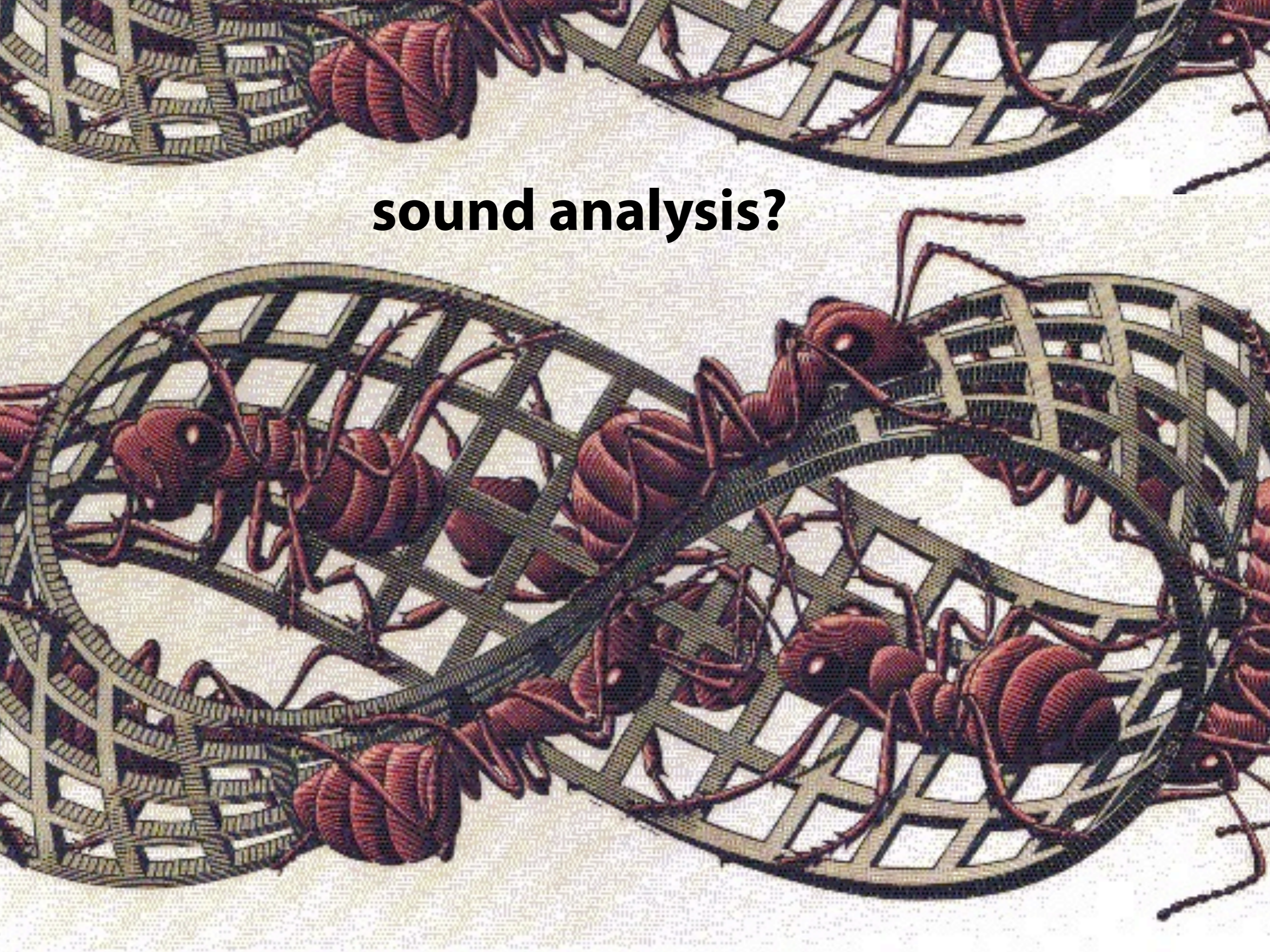
does method m meet pre , $post$?

$$m(s, s') \wedge pre(s) \wedge \neg post(s, s')$$

is op' a valid refactoring of op ?

$$op'(s, s') \wedge \neg op(s, s')$$

sound analysis?



traditional view

traditional view

analysis should be 'sound'

consider all cases

ensures no bugs missed

is a cast-iron guarantee

traditional view

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but sound \Rightarrow spurious errors

\Rightarrow *annoyance*

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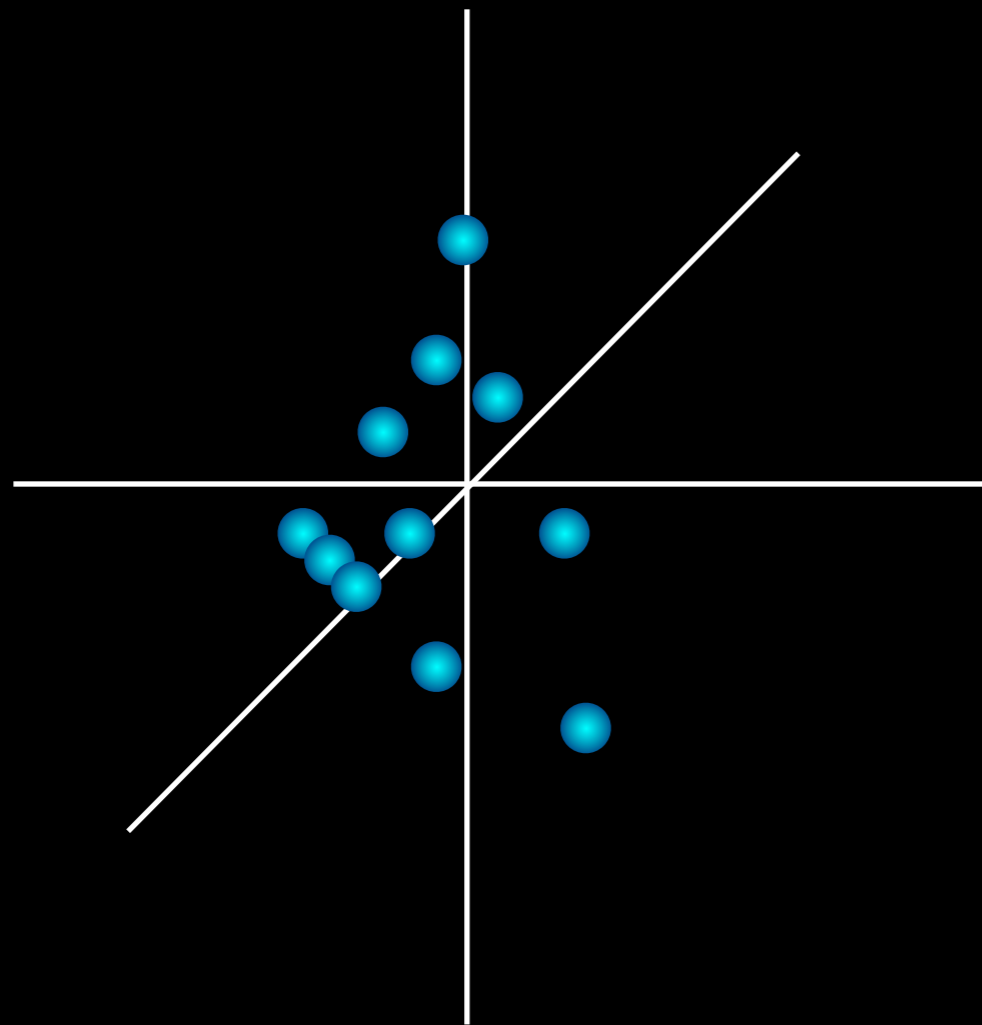
and not a guarantee anyway?

even theorem proving isn't watertight

inconsistent axioms may undermine proof

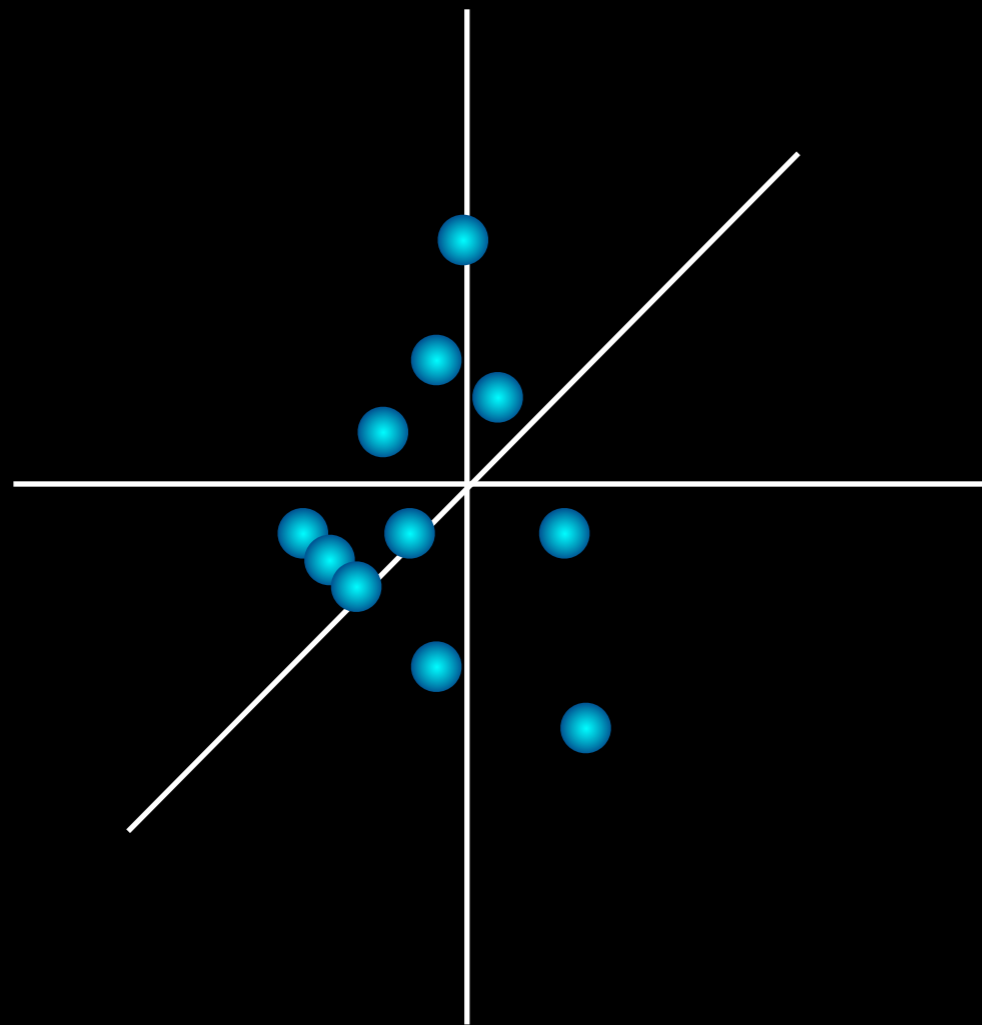
idea: small scope analysis

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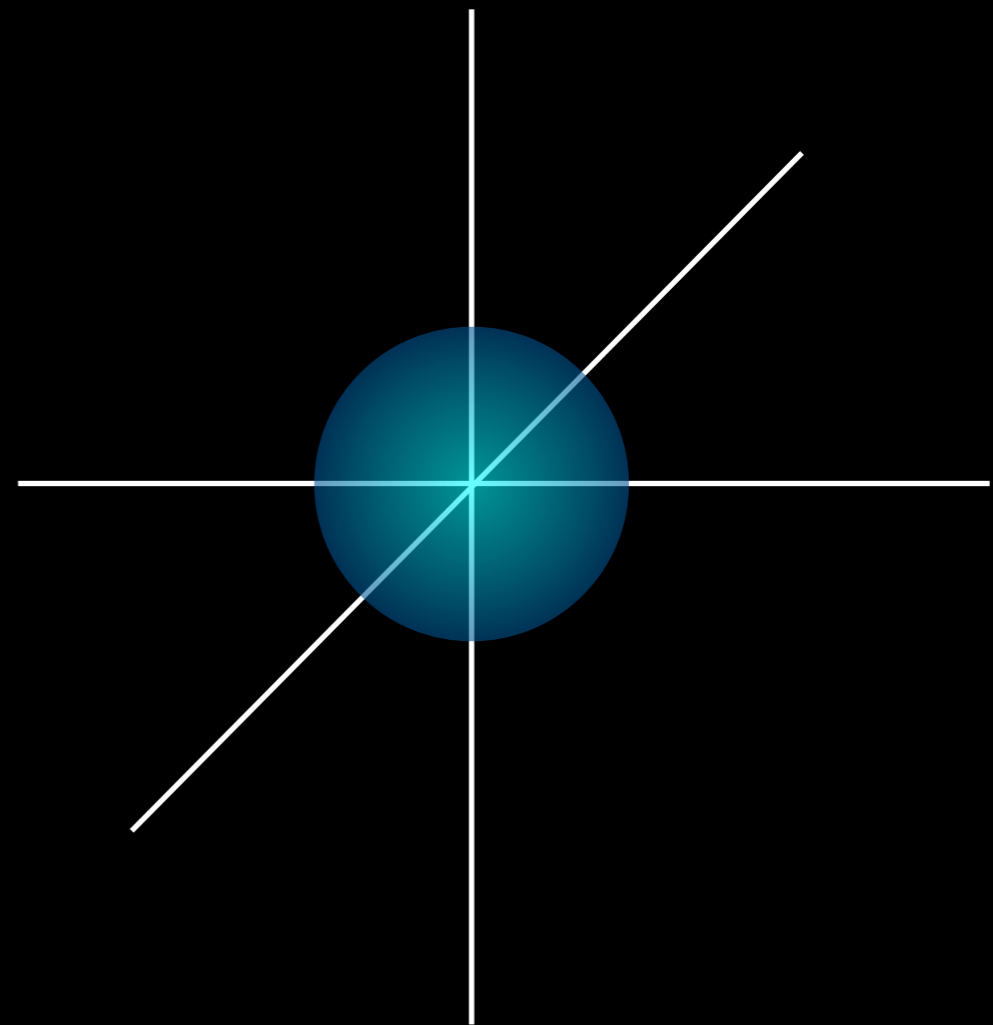


*random testing:
a few big tests*

idea: small scope analysis



*random testing:
a few big tests*



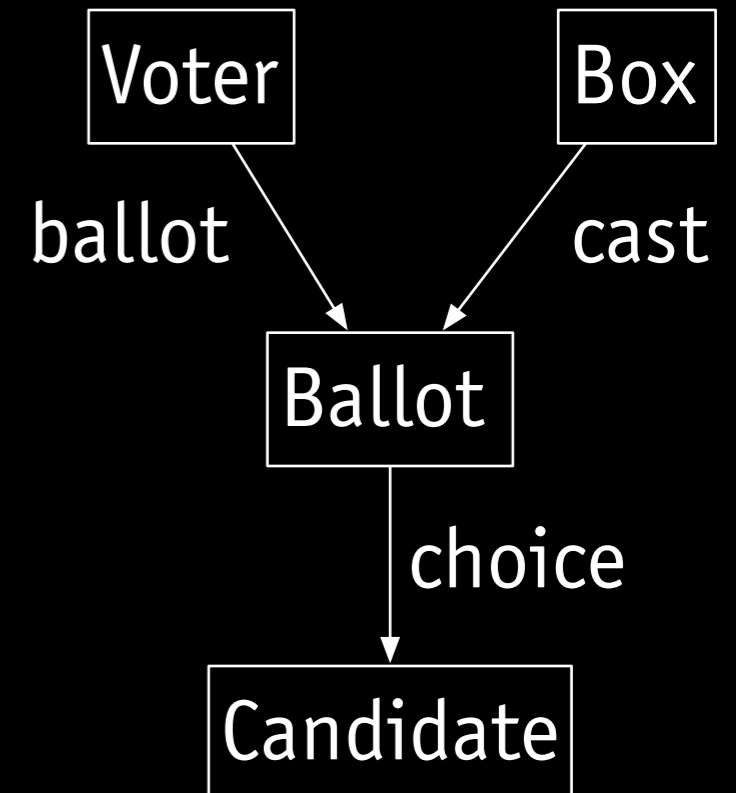
*small scope analysis:
'all small tests'*

idea: generalizing dot

x.f

idea: generalizing dot

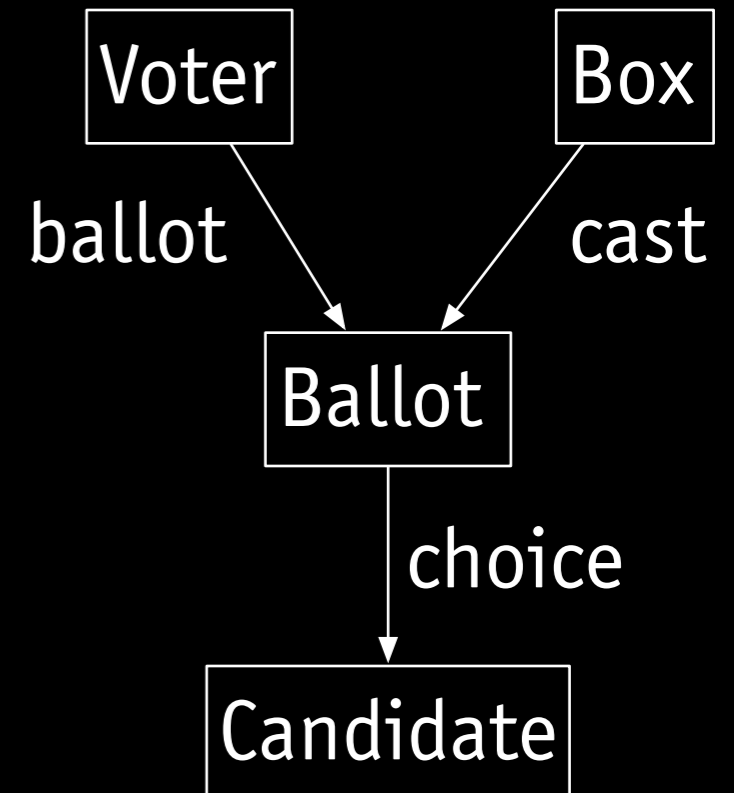
x.f



idea: generalizing dot

x.f

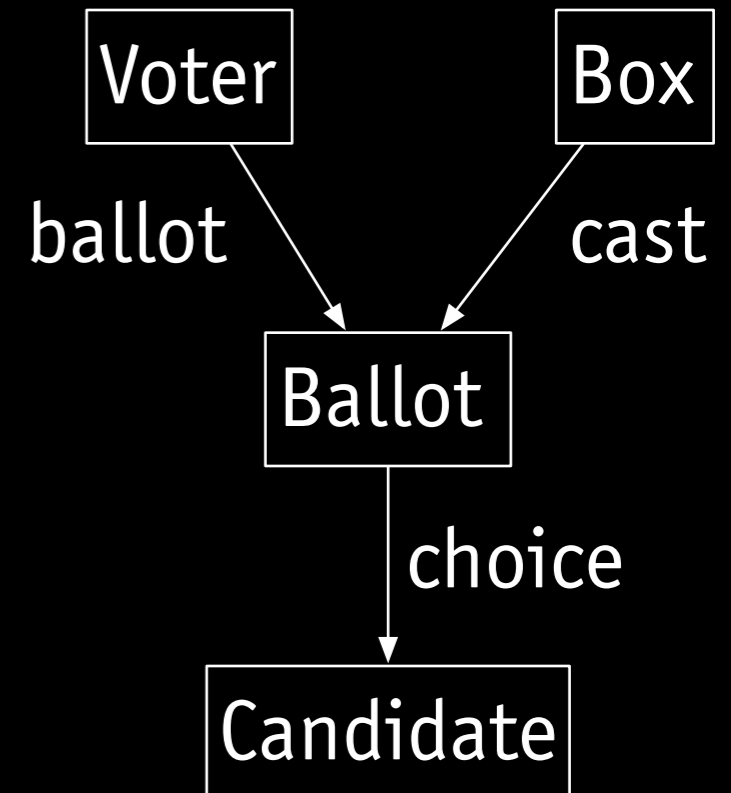
v.ballot



idea: generalizing dot



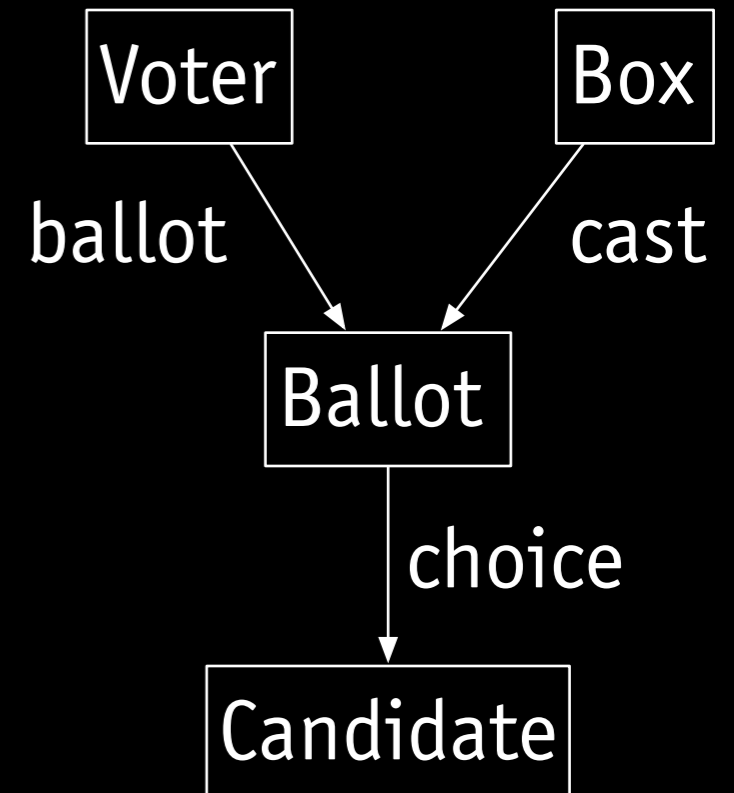
v.ballot
v.ballot.choice



idea: generalizing dot



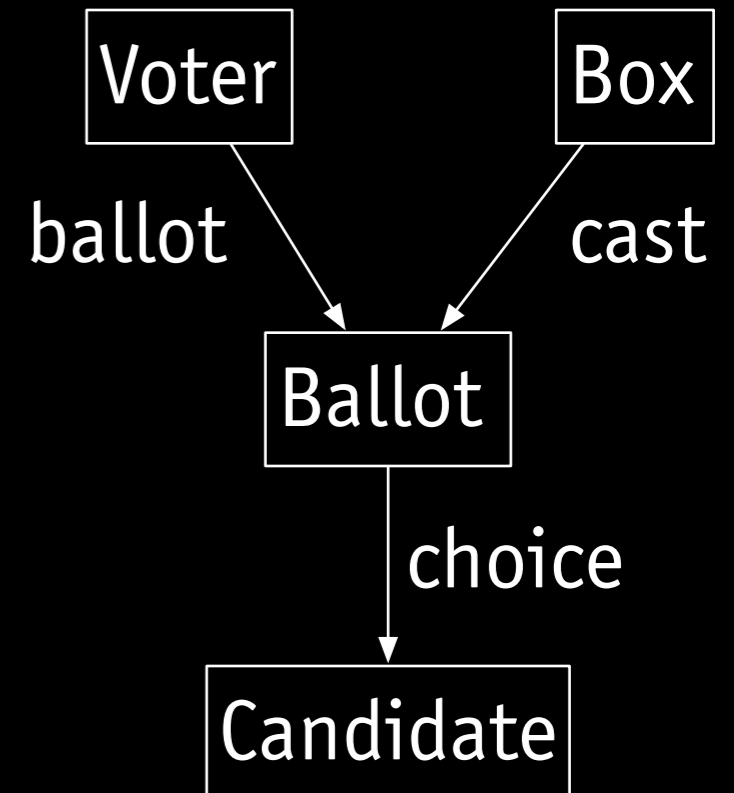
v.ballot
v.ballot.choice
b.cast



idea: generalizing dot



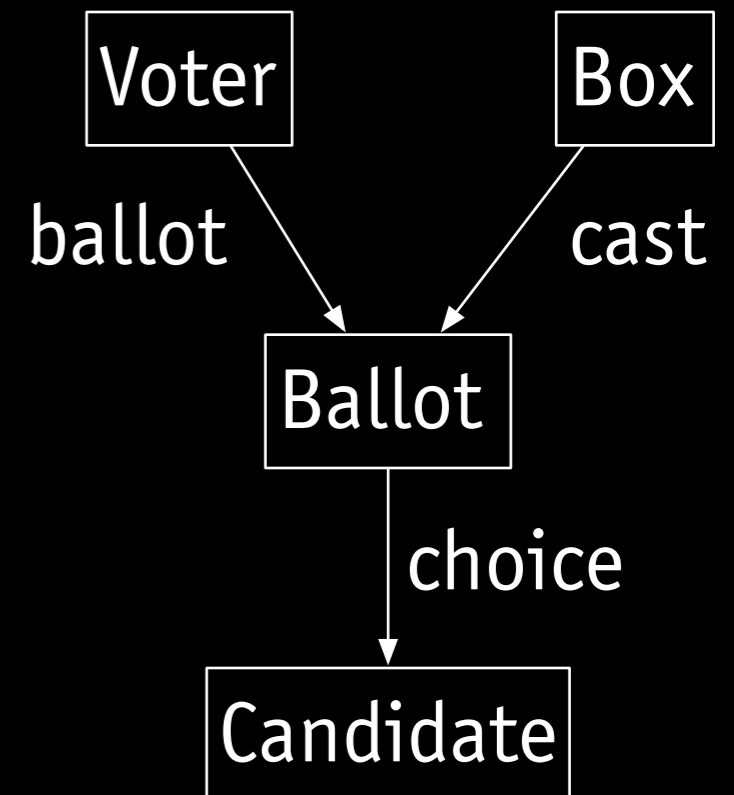
v.ballot
v.ballot.choice
b.cast
b.cast.choice



idea: generalizing dot



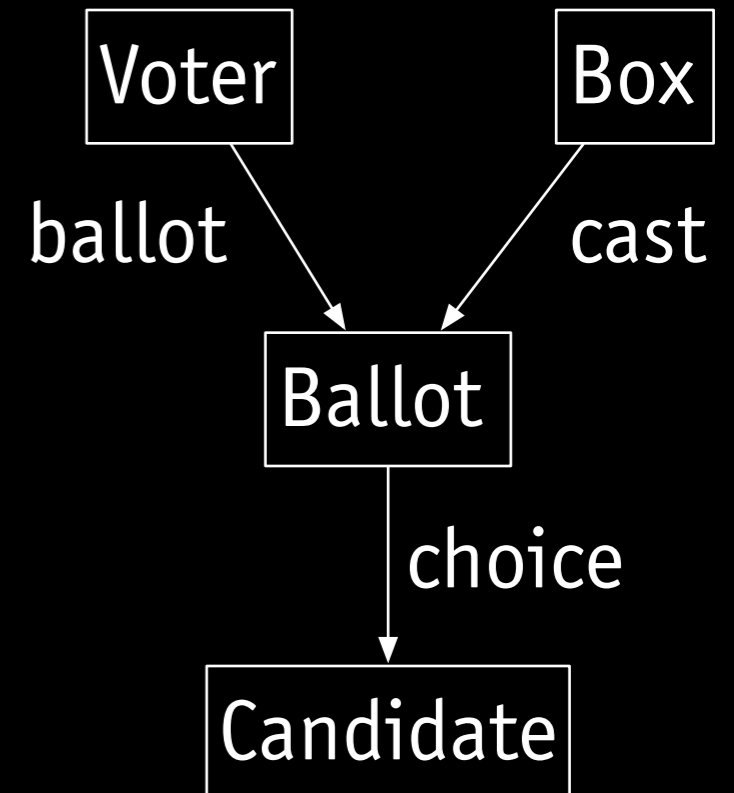
v.ballot
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b.cast
b.cast.choice
choice.c



idea: generalizing dot



v.ballot
v.ballot.choice
b.cast
b.cast.choice
choice.c
ballot.choice.c



a definition

generalized join

$$a.b = \{ (a_0, \dots, a_{n-1}, b_1, \dots, b_m) \mid (a_0, \dots, a_n) \in a \wedge (a_n, b_1, \dots, b_m) \in b \}$$

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v • *ballot* • *choice*

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V0

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v • *ballot* • *choice*

V0

V0	B0
V1	B2
V2	B1

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v • *ballot* • *choice*

V0

V0	B0
----	----

V1

B2

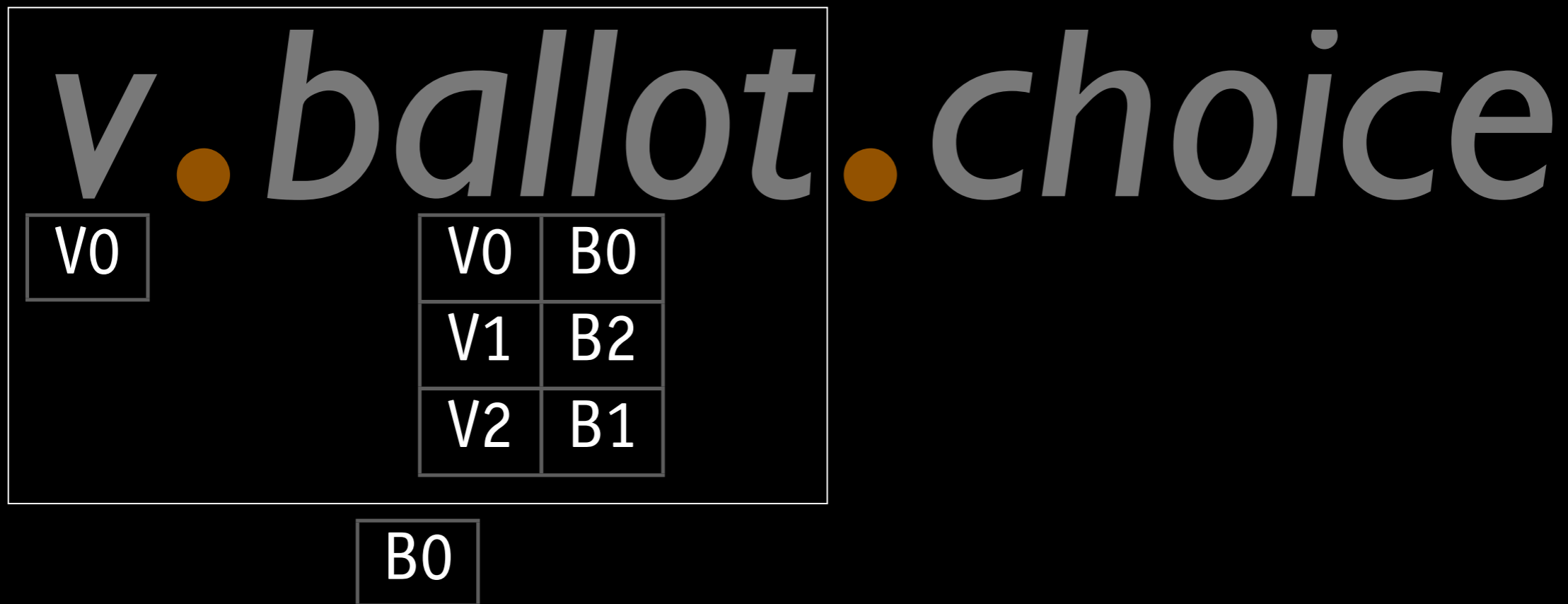
V2

B1

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v.*ballot*.*choice*

B0	C0
B1	C1
B2	C0

B0

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v.ballot.choice

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B0

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v.ballot.choice

right associating

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v. *ballot*. *choice*

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v

v0

ballot • *choice*

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v

v0

ballot.choice

v0

c0

v1

c0

v2

c1

c0

right associating

generalized join

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v • *ballot* • *choice*

navigating back

generalized join

$$a.b = \{ (a_0, \dots, a_{n-1}, b_1, \dots, b_m) \mid (a_0, \dots, a_n) \in a \wedge (a_n, b_1, \dots, b_m) \in b \}$$

ballot.choice.c

navigating back

generalized join

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ballot.choice.c

CO

navigating back

generalized join

$$a.b = \{ (a_0, \dots, a_{n-1}, b_1, \dots, b_m) \mid (a_0, \dots, a_n) \in a \wedge (a_n, b_1, \dots, b_m) \in b \}$$

ballot.choice.*c*

B0	C0
B1	C1
B2	C0

C0

navigating back

generalized join

$$a.b = \{ (a_0, \dots, a_{n-1}, b_1, \dots, b_m) \mid (a_0, \dots, a_n) \in a \wedge (a_n, b_1, \dots, b_m) \in b \}$$

ballot. *choice*. *c*

B0	C0
B1	C1
B2	C0

C0

navigating back

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ballot.choice.*c*

B0	C0
B1	C1
B2	C0

C0

B0
B2

navigating back

generalized join

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ballot.choice.c

B0

B2

navigating back

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ballot.*choice*.*c*

V0	B0
V1	B2
V2	B1

B0
B2

navigating back

generalized join

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ballot.choice.*c*

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B0
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ballot.choice.*c*

V0	B0
V1	B2
V2	B1

B0
B2

V0
V1

idea: solving by SAT

$v \cdot \text{ballot} = b$

idea: solving by SAT

v . ballot = b

v0

idea: solving by SAT

$v \cdot \text{ballot} = b$

v_0	1
v_1	0
v_2	0

v_0

idea: solving by SAT

$v \cdot \text{ballot} = b$

v_0	1
v_1	0
v_2	0

idea: solving by SAT

$v \cdot \text{ballot} = b$

v_0	1
v_1	0
v_2	0

B0

idea: solving by SAT

v . ballot = b

<i>V0</i>	1
<i>V1</i>	0
<i>V2</i>	0

<i>B0</i>	1
<i>B1</i>	0
<i>B2</i>	0

B0

idea: solving by SAT

$v \cdot \text{ballot} = b$

$V0$	1
$V1$	0
$V2$	0

$B0$	1
$B1$	0
$B2$	0

idea: solving by SAT

$v \cdot \text{ballot} = b$

$V0$	1
$V1$	0
$V2$	0

$B0$	1
$B1$	0
$B2$	0

$V0$	$B0$
$V1$	$B2$
$V2$	$B1$

idea: solving by SAT

$v \cdot \text{ballot} = b$

$V0$	1
$V1$	0
$V2$	0

	$V0$	$V1$	$V2$
$B0$	1	0	0
$B1$	0	0	1
$B2$	0	1	0

$B0$	1
$B1$	0
$B2$	0

$V0$	$B0$
$V1$	$B2$
$V2$	$B1$

idea: solving by SAT

$v \cdot \text{ballot} = b$

$V0$	1
$V1$	0
$V2$	0

	$V0$	$V1$	$V2$
$B0$	1	0	0
$B1$	0	0	1
$B2$	0	1	0

$B0$	1
$B1$	0
$B2$	0

idea: solving by SAT

v . ballot = b

v0 v1 v2

v0

b0

b0

v1

b1

b1

v2

b2

b2

idea: solving by SAT

v . ballot = b

v0 v1 v2

<i>v0</i>	<i>v0</i>
<i>v1</i>	<i>v1</i>
<i>v2</i>	<i>v2</i>

B0
B1
B2

B0
B1
B2

idea: solving by SAT

v . ballot = b

v_0 v_1 v_2

v_0	v_0
v_1	v_1
v_2	v_2

B_0
 B_1
 B_2

B_0	b_0
B_1	b_1
B_2	b_2

idea: solving by SAT

$v \cdot \text{ballot} = b$

V_0	v_0
V_1	v_1
V_2	v_2

	V_0	V_1	V_2
B_0	f_{00}	f_{01}	f_{02}
B_1	f_{10}	f_{11}	f_{12}
B_2	f_{20}	f_{21}	f_{22}

B_0	b_0
B_1	b_1
B_2	b_2

idea: solving by SAT

$v \cdot \text{ballot} = b$

$V0$	v_0
$V1$	v_1
$V2$	v_2

	$V0$	$V1$	$V2$
$B0$	f_{00}	f_{01}	f_{02}
$B1$	f_{10}	f_{11}	f_{12}
$B2$	f_{20}	f_{21}	f_{22}

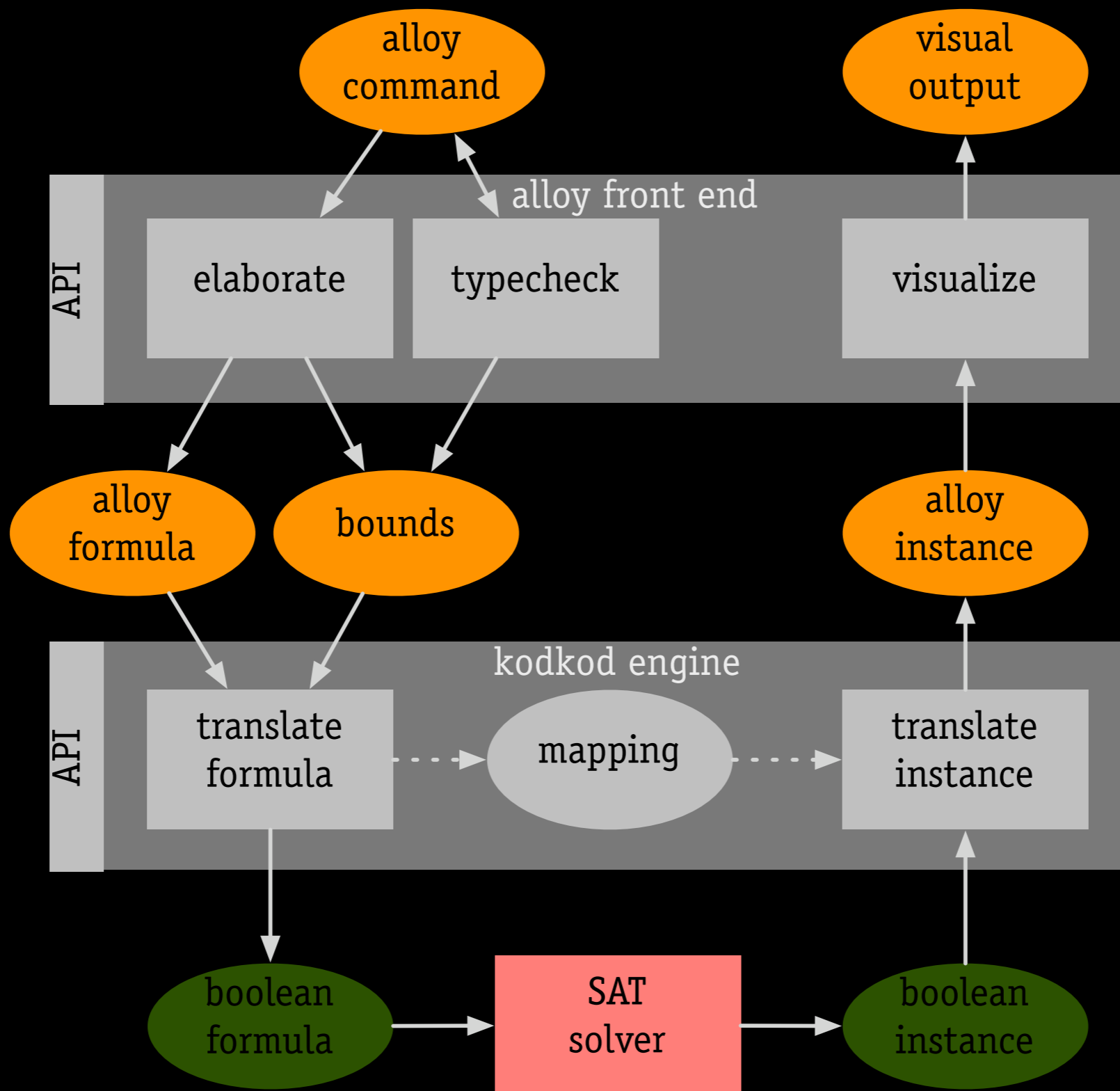
$B0$	b_0
$B1$	b_1
$B2$	b_2

$$b_0 = v_0 \wedge f_{00} \vee v_1 \wedge f_{10} \vee v_2 \wedge f_{20}$$

$$b_1 = v_0 \wedge f_{01} \vee v_1 \wedge f_{11} \vee v_2 \wedge f_{21}$$

$$b_2 = v_0 \wedge f_{02} \vee v_1 \wedge f_{12} \vee v_2 \wedge f_{22}$$

analyzer architecture



idea: **funky idioms**

idea: funky idioms

alloy is just a logic
with subtypes, navigation, etc
but no states or mutation

idea: funky idioms

alloy is just a logic
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for dynamic models
many ways to represent behaviour

idea: funky idioms

alloy is just a logic
with subtypes, navigation, etc
but no states or mutation

for dynamic models
many ways to represent behaviour

a surprise
least conventional most useful

funky idioms: stateful objects

```
sig Candidate {}
```

```
sig Voter { intent: Candidate }
```

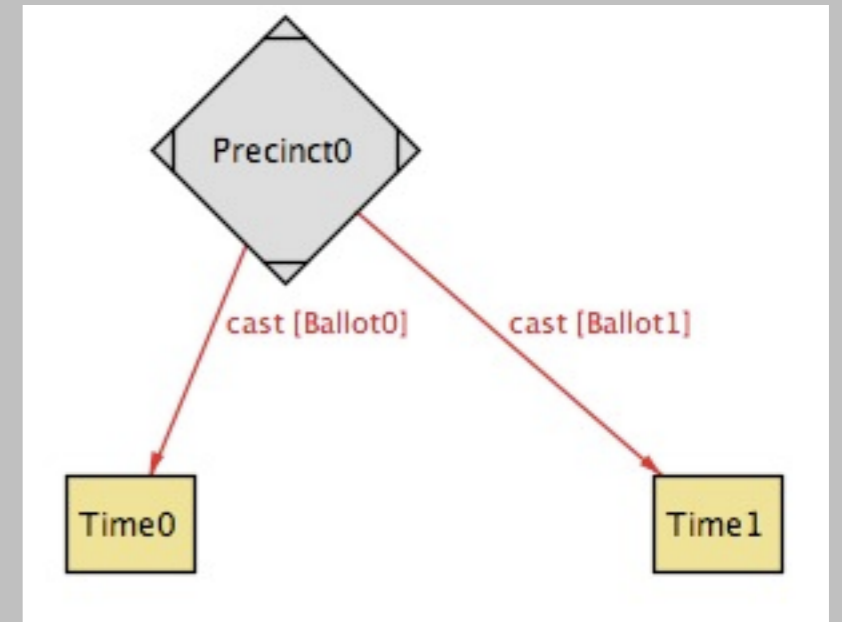
```
sig Ballot { choice: Candidate }
```

```
sig Tally { ballots: set Ballot, count: Candidate -> one Int }
```

```
sig Precinct {  
  registered: set Voter,  
  cast: Ballot -> Time,  
  voted: Voter -> Time  
}
```

funky idioms: stateful objects

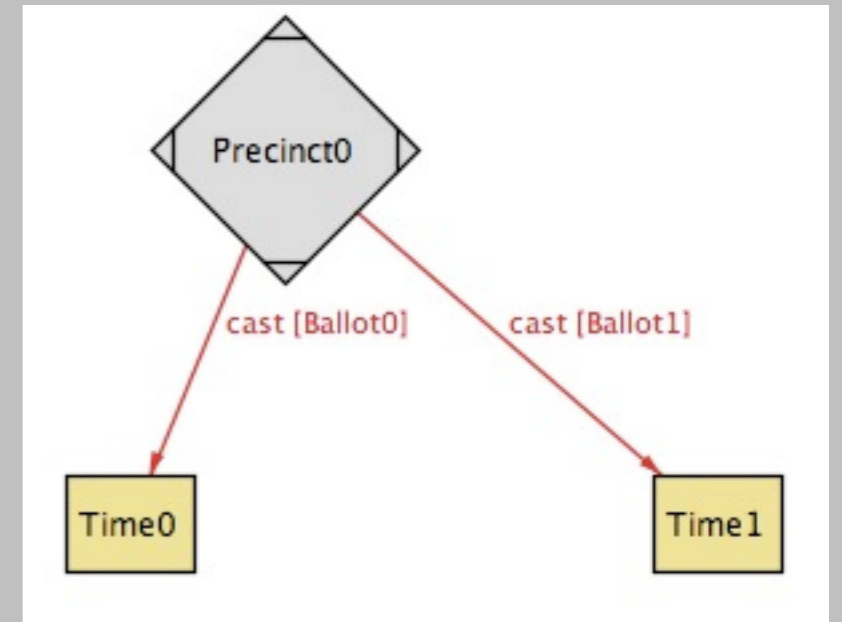
```
sig Candidate {}  
sig Voter { intent: Candidate }  
sig Ballot { choice: Candidate }  
sig Tally { ballots: set Ballot, count: Candidate -> one Int }  
sig Precinct {  
  registered: set Voter,  
  cast: Ballot -> Time,  
  voted: Voter -> Time  
}
```



cast: a ternary relation

funky idioms: stateful objects

```
sig Candidate {}  
sig Voter { intent: Candidate }  
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```



cast: a ternary relation

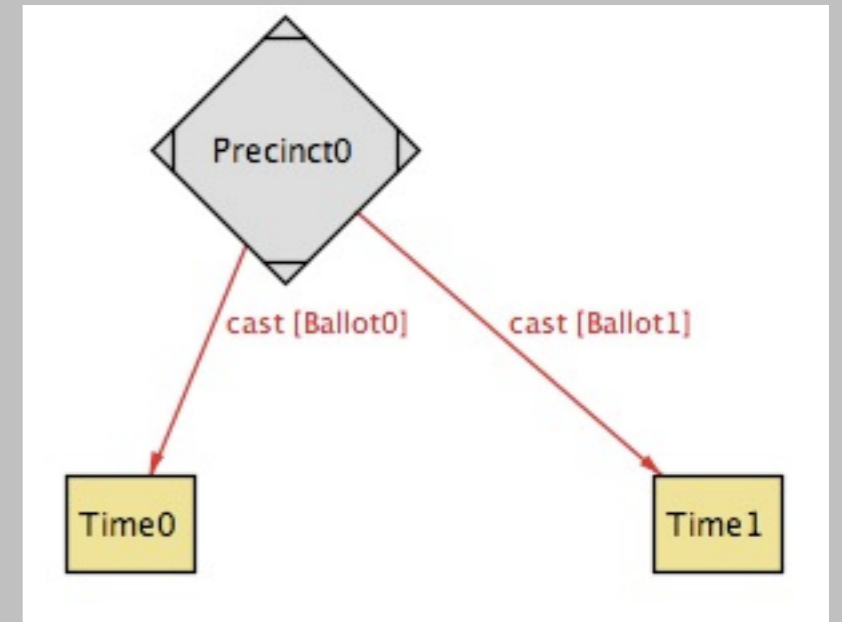
cast: Precinct -> Ballot -> Time

p.cast: Ballot -> Time

cast.t: Precinct -> Ballot

funky idioms: stateful objects

```
sig Candidate {}  
sig Voter { intent: Candidate }  
sig Ballot { choice: Candidate }  
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```

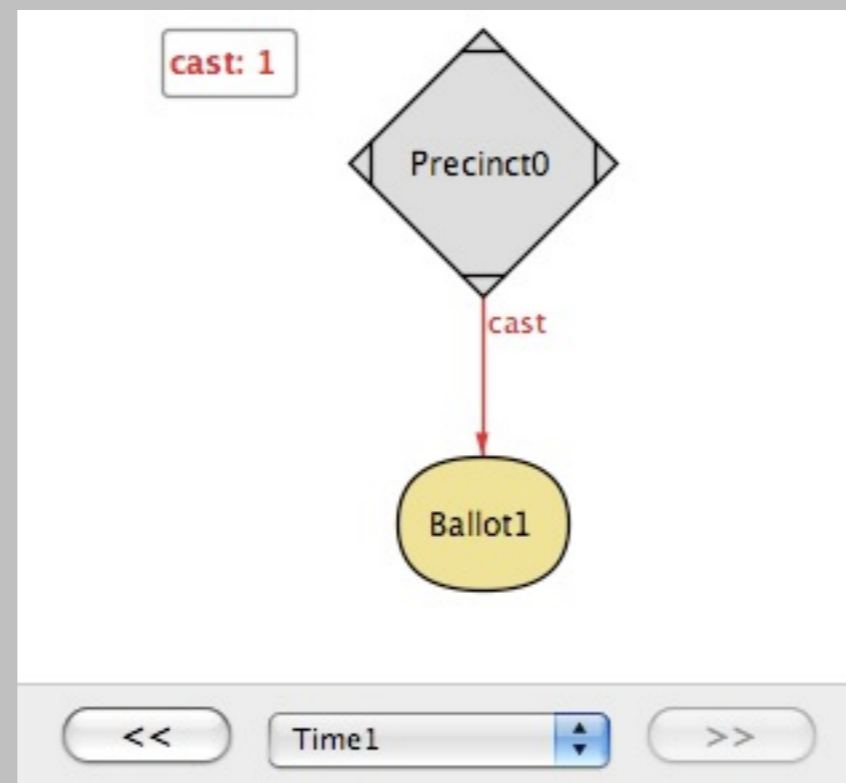
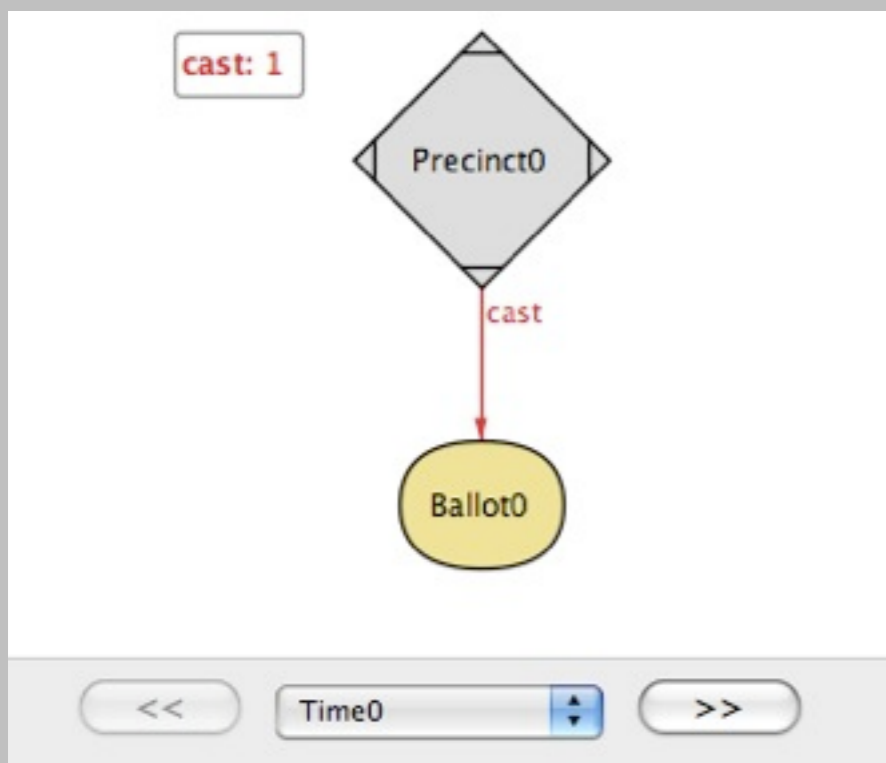


cast: a ternary relation

cast: Precinct -> Ballot -> Time

p.cast: Ballot -> Time

cast.t: Precinct -> Ballot



cast: projected onto Time

funky idioms: events

```
open util/ordering[Time]
```

```
sig Time {}
```

```
abstract sig Event {before, after: Time}
```

```
pred Event.follows (that: Event) { this.before = that.after }
```

```
fact {
```

```
  all t: Time - last | one before.t
```

```
-- one event starting in every time instant
```

```
  all e: Event | e.after = e.before.next
```

```
-- events execute in time order
```

```
}
```

funky idioms: events

```
open util/ordering[Time]
```

```
sig Time {}
```

```
abstract sig Event {before, after: Time}
```

```
pred Event.follows (that: Event) { this.before = that.after }
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fact {
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```
  all t: Time - last | one before.t
```

```
  all e: Event | e.after = e.before.next
```

```
}
```

a generic model of events

-- one event starting in every time instant

-- events execute in time order

funky idioms: events

```
open util/ordering[Time]
```

```
sig Time {}
```

```
abstract sig Event {before, after: Time}
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pred Event.follows (that: Event) { this.before = that.after }
```

```
fact {
```

```
  all t: Time - last | one before.t
```

```
  all e: Event | e.after = e.before.next
```

```
}
```

a generic model of events

-- one event starting in every time instant

-- events execute in time order

```
sig Cast extends Event { voter: Voter, ballot: Ballot, precinct: Precinct } {
```

```
  ballot.choice = voter.intent
```

-- marks ballot according to intent

```
  voter in precinct.registered - precinct.voted.before
```

-- registered but hasn't voted yet

```
  voted.after = voted.before + precinct -> voter
```

-- has now voted

```
  cast.after = cast.before + precinct -> ballot
```

-- ballot is cast

```
}
```

```
sig Count extends Event { tally: Tally } {
```

```
  tally.ballots = Precinct.cast.before
```

-- ballots tallied are those cast in all precincts

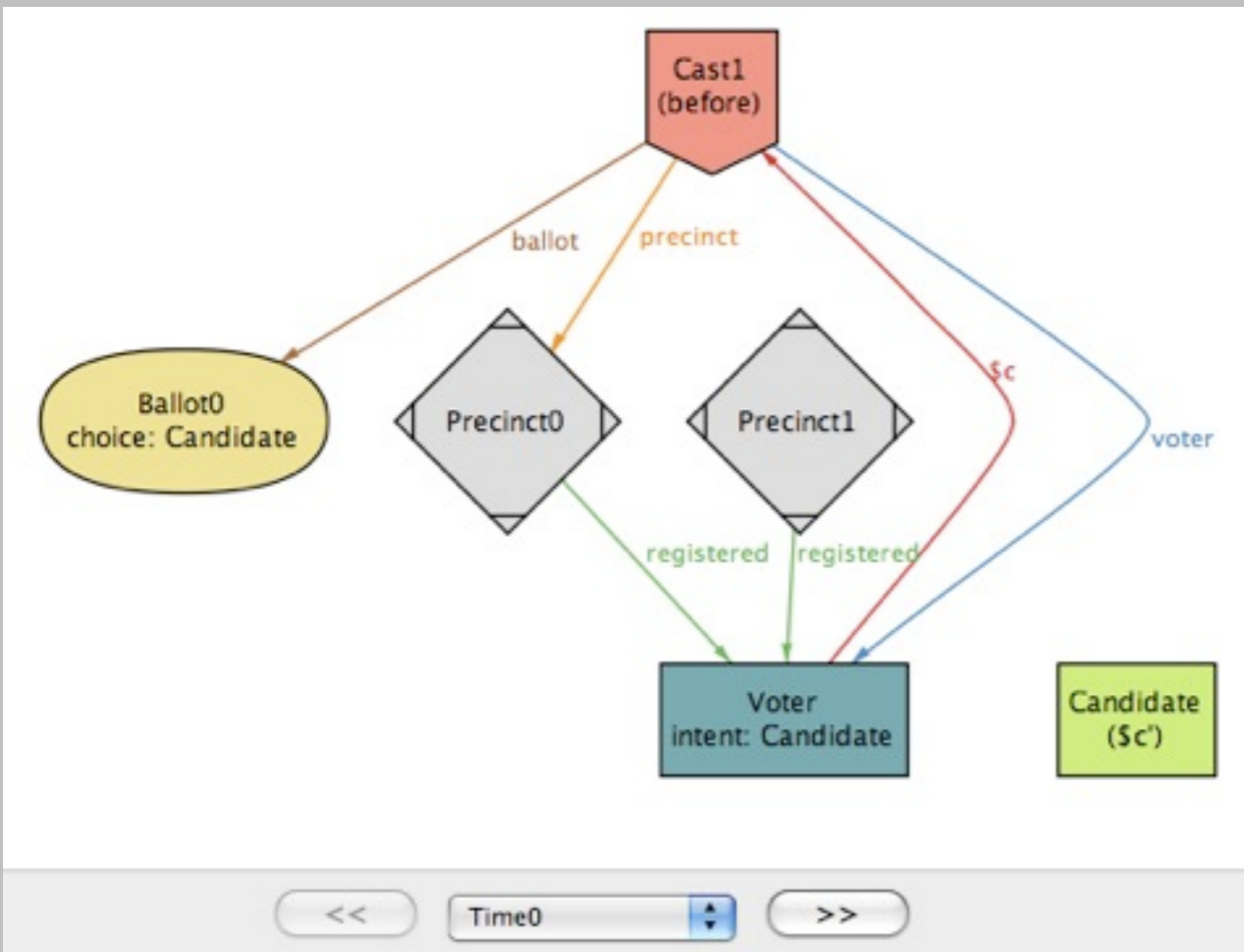
```
}
```

funky idioms: dynamics

```
fact {  
  no cast.first                                -- at first, ballot boxes are empty  
  all v: Voter | some c: Cast | c.voter = v    -- all voters cast ballots  
  no ca: Cast, co: Count | ca.follows [co]     -- no casting after counting  
}
```

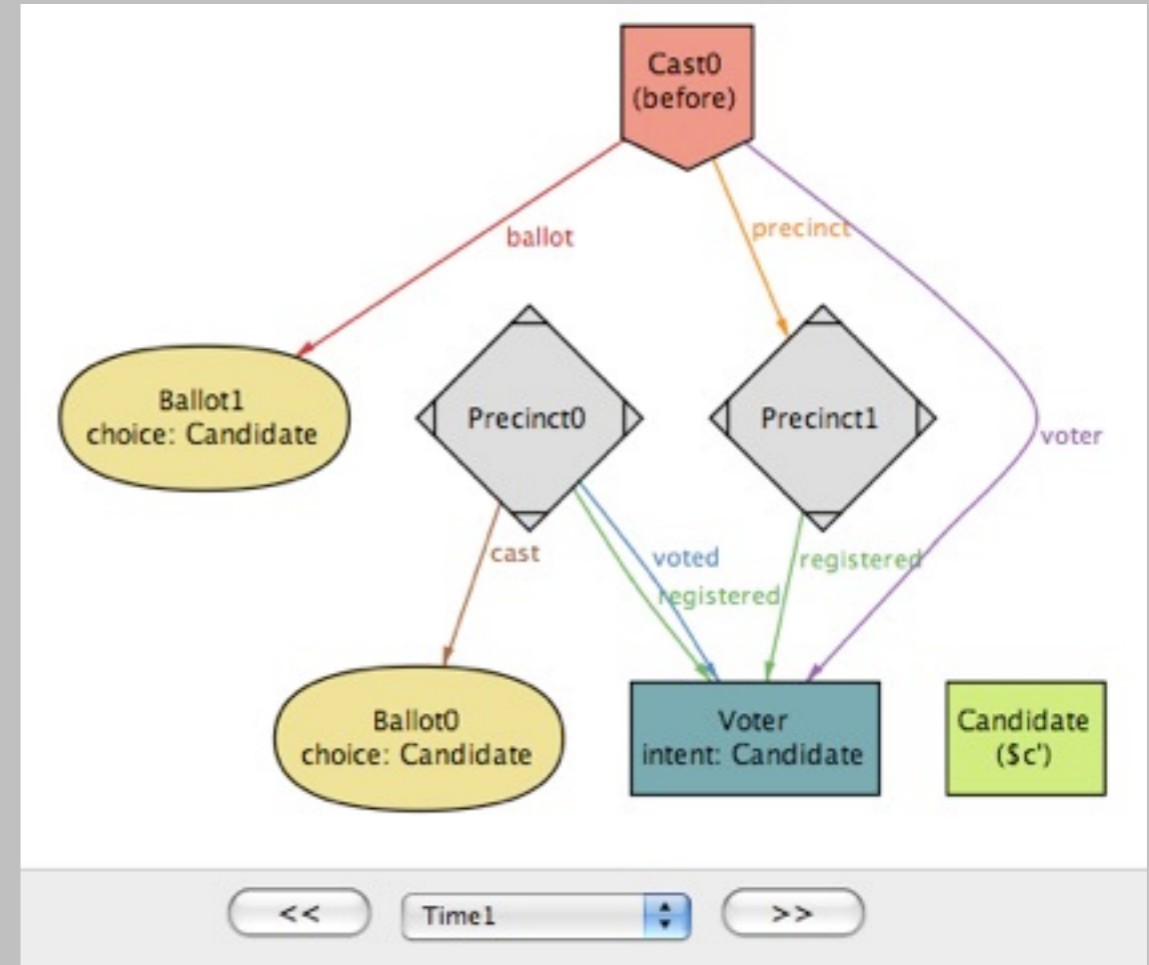
funky idioms: dynamics

```
fact {  
  no cast.first -- at first, ballot boxes are empty  
  all v: Voter | some c: Cast | c.voter = v -- all voters cast ballots  
  no ca: Cast, co: Count | ca.follows [co] -- no casting after counting  
}  
  
check { all e: Count | all c: Candidate | e.tally.count[c] = #intent.c } for 4 -- tally matches voter intent
```

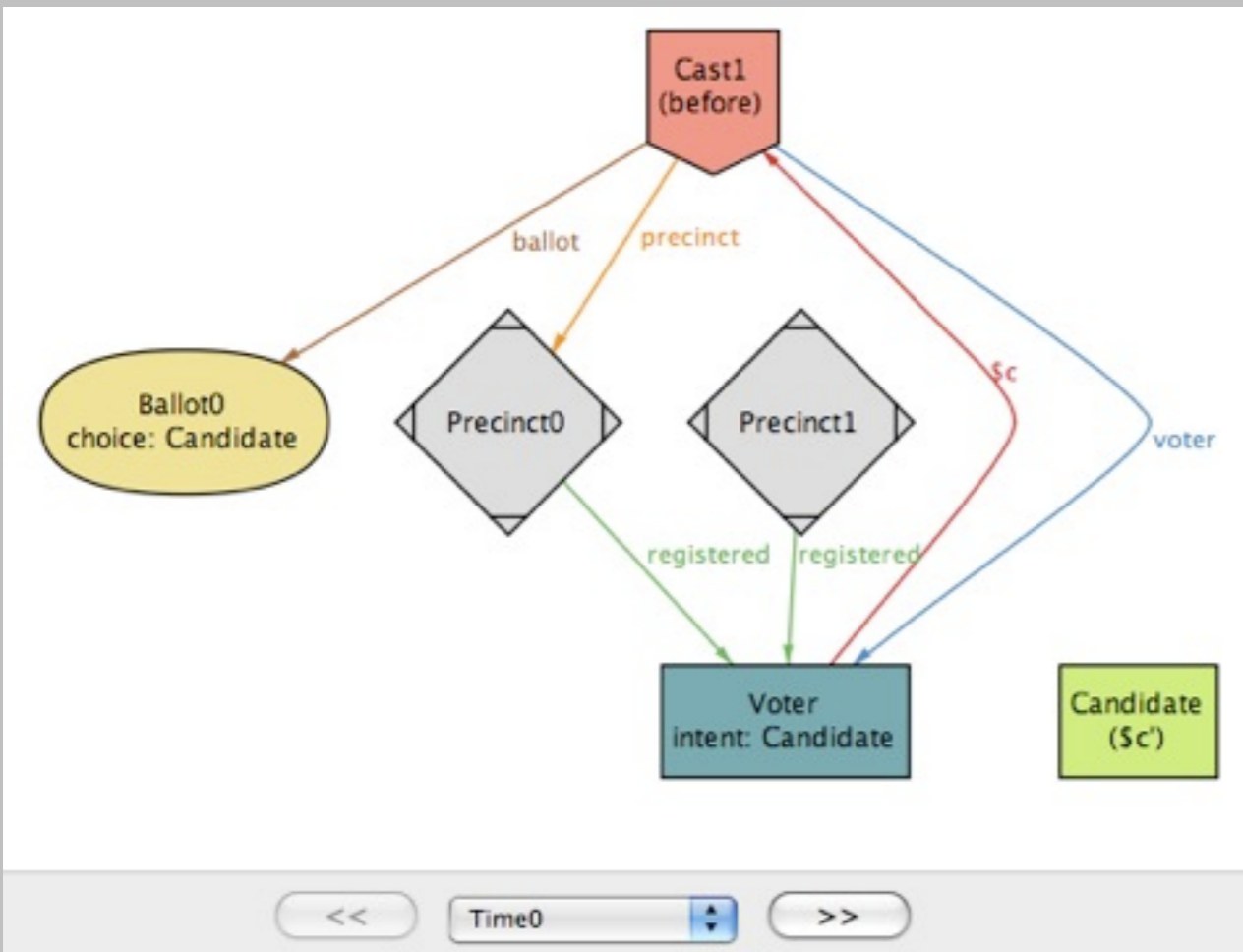



Voter casts Ballot0 in Precinct0

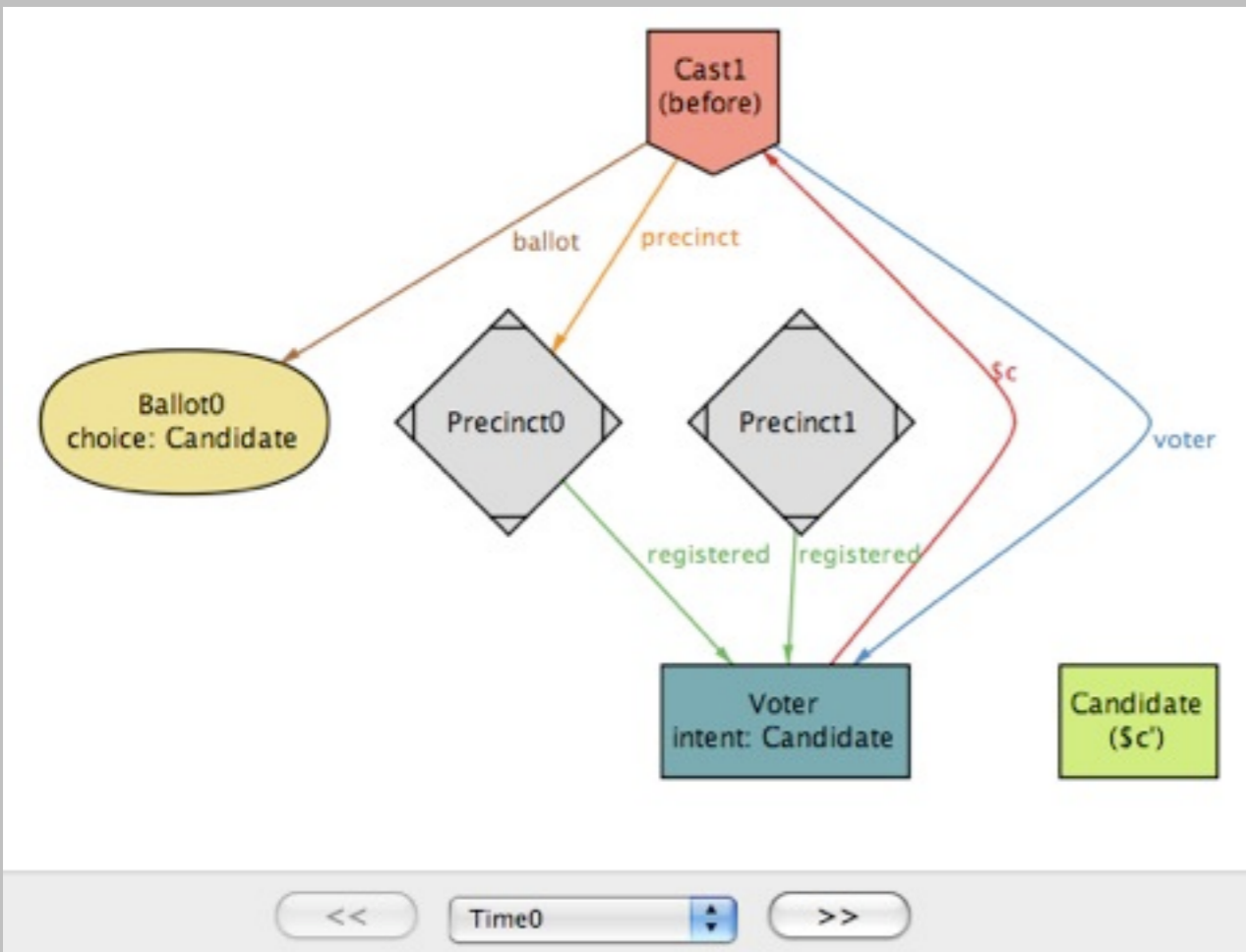
Voter casts Ballot1 in Precinct1



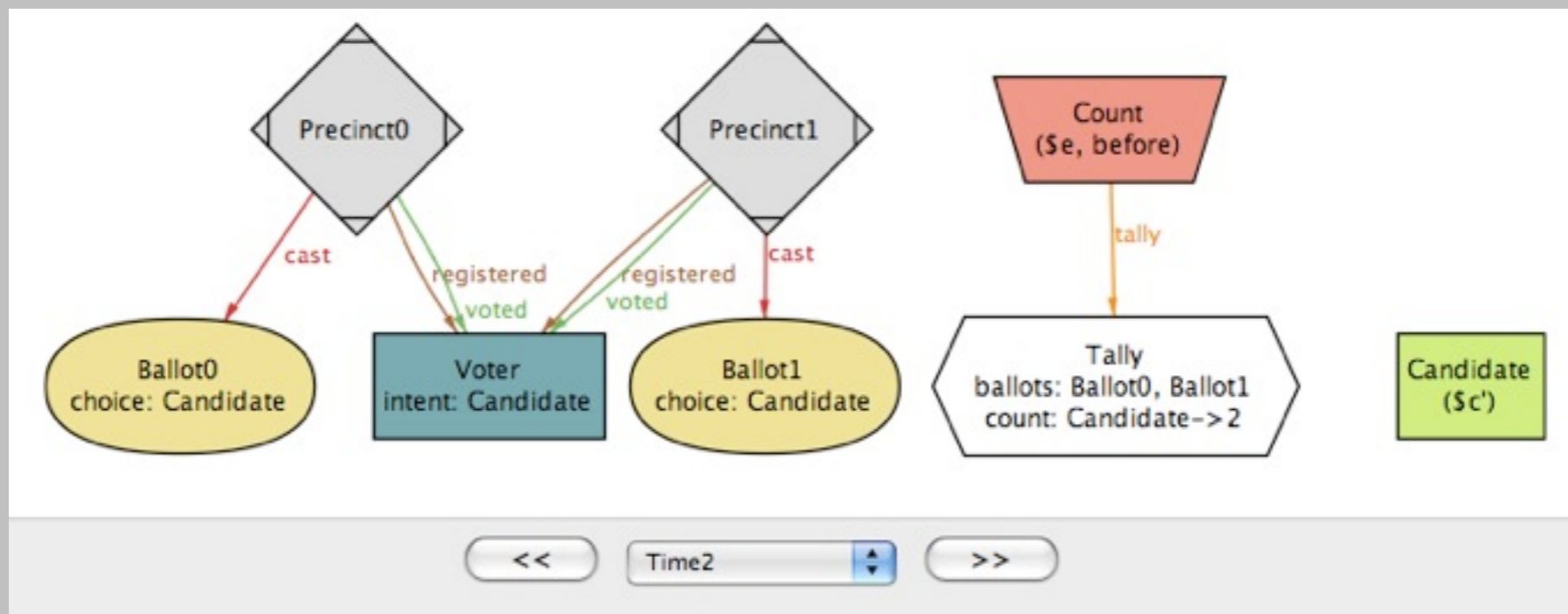
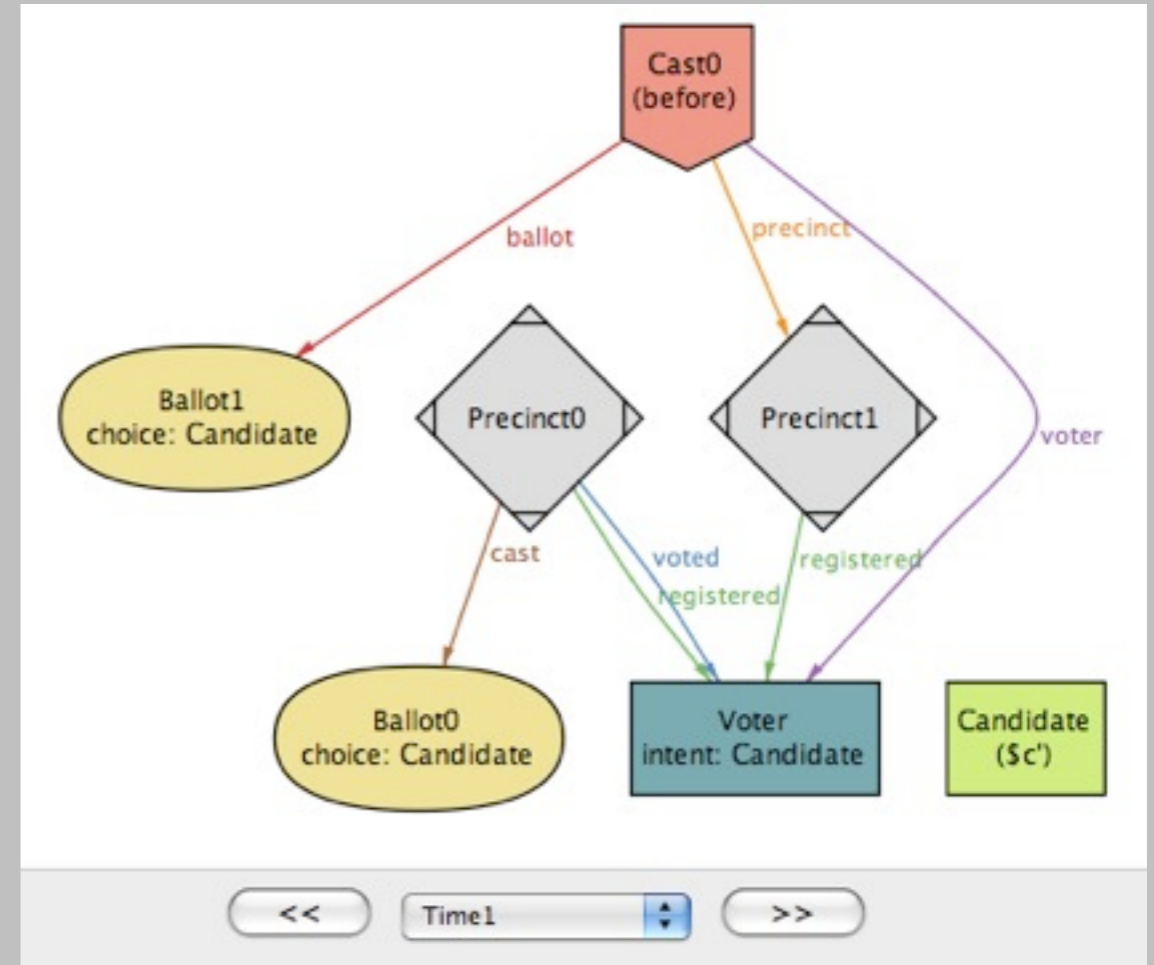
Voter casts Ballot0 in Precinct0



Voter casts Ballot1 in Precinct1



Voter casts Ballot0 in Precinct0



what else?

language

implicit subtyping
sequences, arithmetic
meta features

analysis

coverage by unsat core
partial instances
sharing, symmetry, etc

navigations are succinct

navigations are succinct

tally count for candidate c is $\# \text{ intent}.c$

navigations are succinct

tally count for candidate c is $\# \text{intent}.c$

in Z

$| \text{intent}^{\sim}(\{c\}) |$

navigations are succinct

tally count for candidate c is $\# \text{intent}.c$

in Z

$|\text{intent}^{-1}(\{c\})|$

in SQL

$\text{select count}(\ast) \text{ from voter where intent} = c$

navigations are succinct

tally count for candidate c is $\# \text{intent}.c$

in Z

$|\text{intent}^{-1}(\{c\})|$

in SQL

select count() from voter where intent = c*

in Java

aagh!

some applications

design analyses

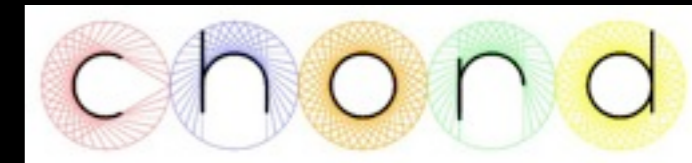
Mondex [Ramanandro, 2006]
reduced Z to Alloy & checked



flash file systems [Kang, 2008]
designs checked against POSIX



Chord [Zave, 2009]
a peer to peer protocol



typically a few billion cases and ...

kloc of Alloy
months of modeling
minutes of analysis

a typical story

Three features that distinguish Chord from many other peer-to-peer lookup protocols are its simplicity, provable correctness, and provable performance.

Ion Stoica et al. Chord: A Scalable Peer to Peer Lookup Service for Internet Applications, SIGCOMM 2001 (also TON, 2003)

a typical story

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Ion Stoica et al. Chord: A Scalable Peer to Peer Lookup Service for Internet Applications, SIGCOMM 2001 (also TON, 2003)

Modeling and analysis have shown that the Chord routing protocol is not correct according to its specification. Furthermore, not one of the six logical properties claimed as invariant is invariantly maintained by the protocol.

Pamela Zave. Invariant-Based Verification of Routing Protocols: The Case of Chord, 2009

code analysis studies

graph libraries (Taghdiri, 2006–2008)

Open JGraph, Quartz graph API

KOA Remote Voting System (Dennis, 2007)

found 19 methods not meeting specs

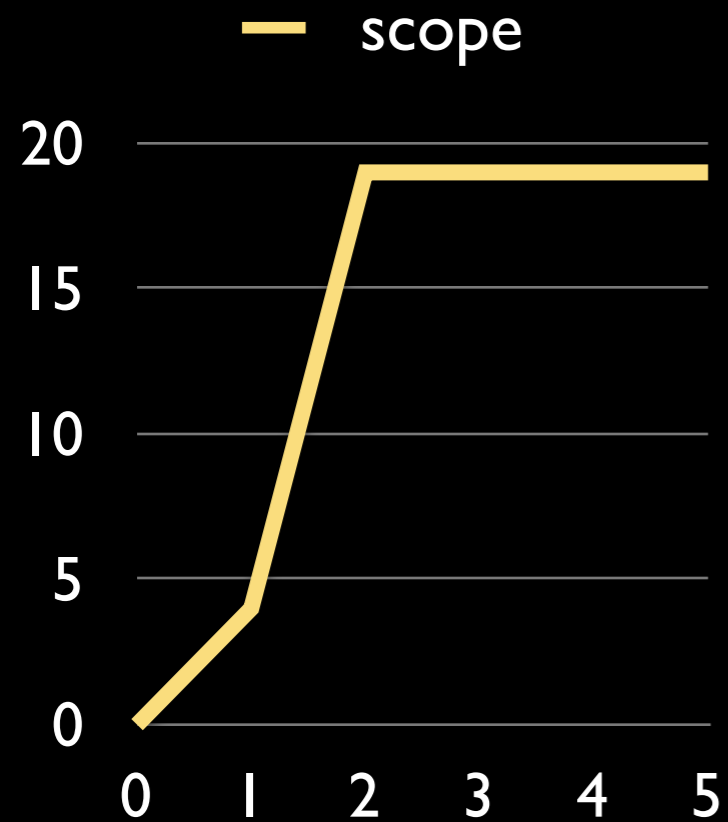
Burr Proton Therapy Center (Seater, 2008)

checked various dependability arguments

small scope hypothesis

analysis of KOA voting code
19 methods violating specs
how many bugs found in scope of k?
[Greg Dennis, 2008]

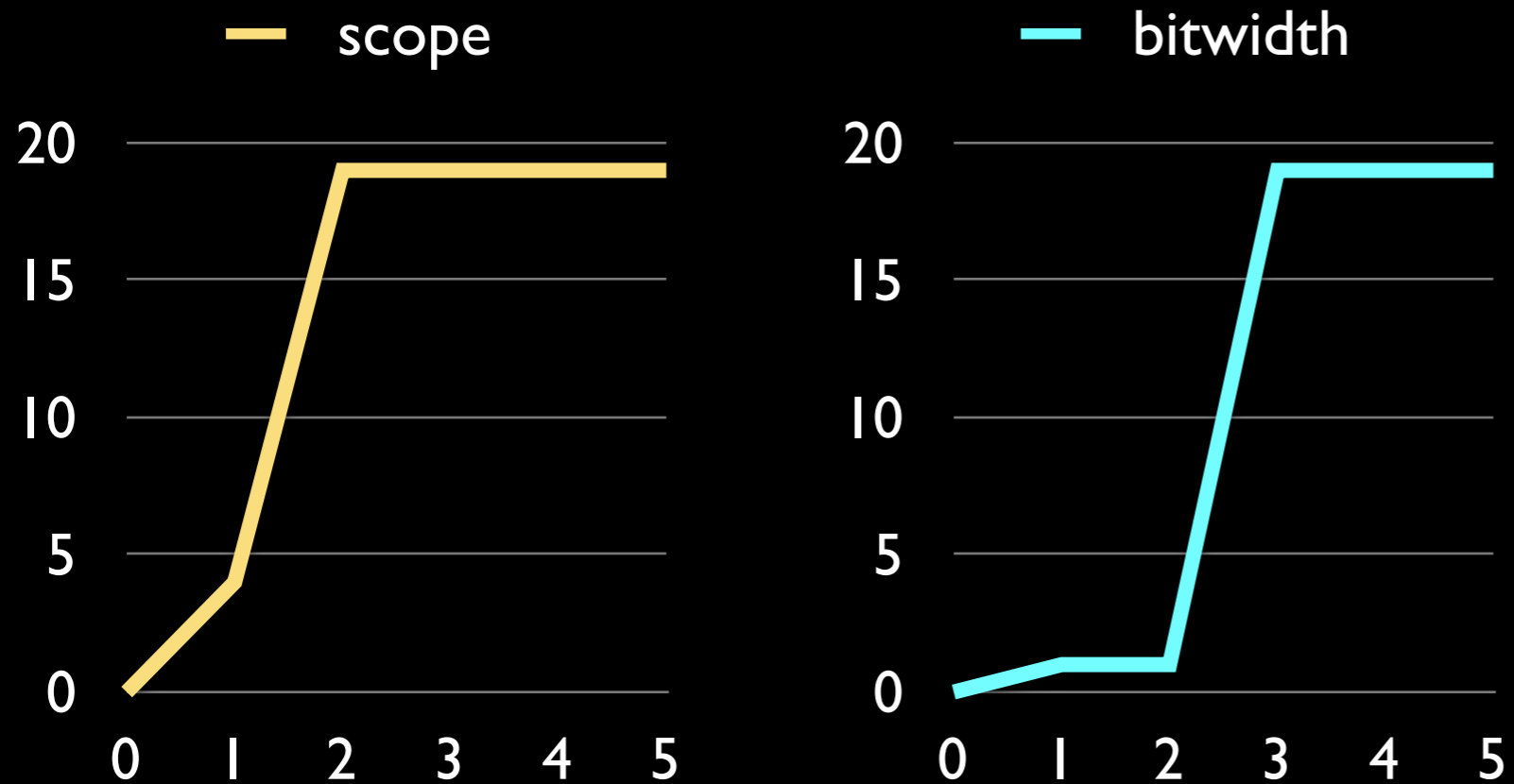
small scope hypothesis



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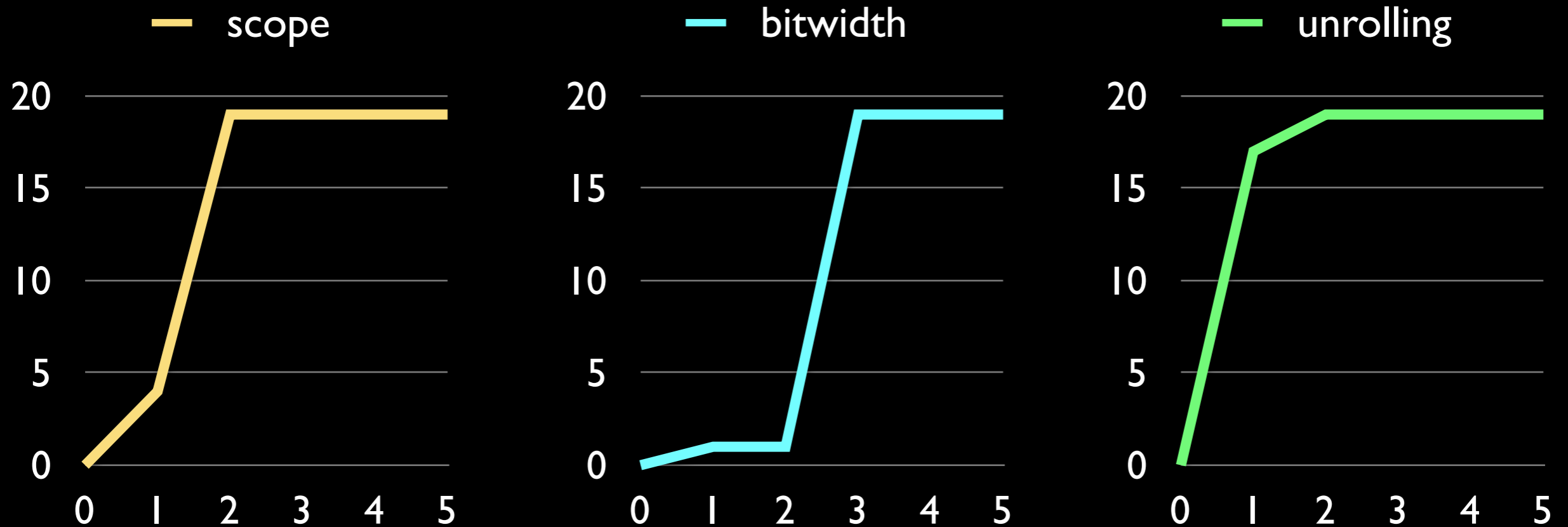
[Greg Dennis, 2008]

small scope hypothesis



analysis of KOA voting code
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[Greg Dennis, 2008]

small scope hypothesis



analysis of KOA voting code
19 methods violating specs
how many bugs found in scope of k?
[Greg Dennis, 2008]

code analysis: jforge plugin

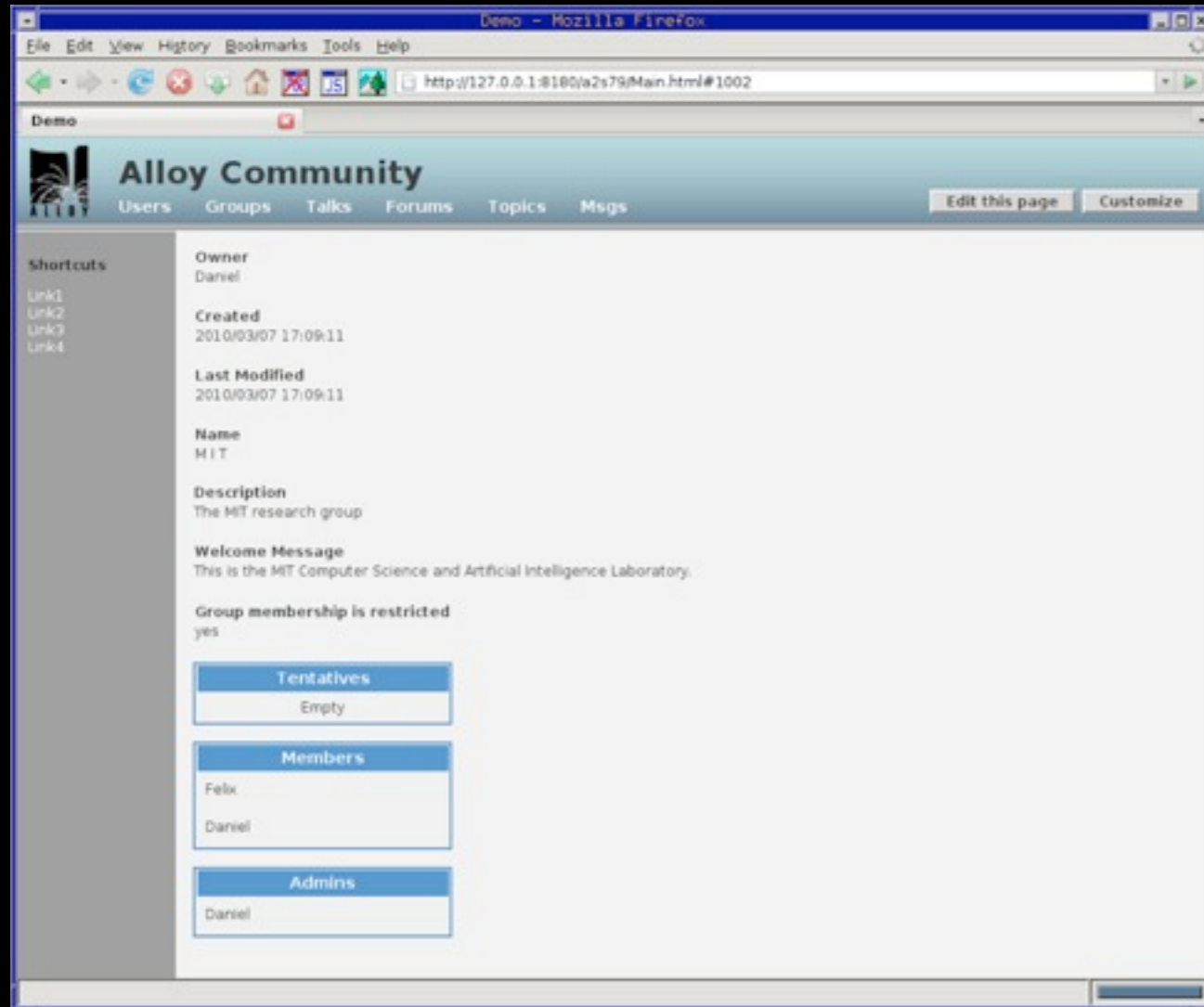
The screenshot shows an IDE window with the following components:

- Menu Bar:** File, Edit, Source, Refactor, Navigate, Search, Project, Run, Window, Help.
- Toolbar:** Includes icons for file operations, search, and execution.
- Left Panel (Project Explorer):** Shows a package hierarchy for `edu.mit.csail.sdg.annotations` with files like `AnnotationPrinter.java`, `AvlNode.java`, `AvlTree.java`, `BugClient.java`, `Bugs.java`, `CodeConverterTest.java`, `CyclicList.java`, `IntTree.java` (highlighted), `LinkedList.java`, `LinkedList2.java`, `MethodArgumentN...`, `ParserRunner.java`, `SpecFieldUse.java`, `TempRunner.java`, `TestBoot.java`, `TestBugClient.java`, `TestIntTree.java`, `TestRegressions.jav...`, `TestRunner.java`, and `TestSpecFieldUse.ja...`.
- Editor:** Displays the source code of `IntTree.java`. The `insert` method is highlighted, showing annotations: `@Requires`, `@Ensures`, and `@Modifies`. The code snippet is:

```
final void insert(Node z) {
    Node yy = null;
    for (Node x = root; x != null;) {
        yy = x;
        if (x.key > z.key)
            x = x.left;
        else
```
- Right Panel (JUnit):** Shows a list of methods including `delete(Node)`, `deleteFixUp(Node)`, and `rotateLeft(Node)`.
- Context Menu:** A menu is open over the `insert` method, listing actions such as `Cut`, `Copy`, `Paste`, `Delete`, `Source`, `Refactor`, `References`, `Declarations`, `Toggle Method Breakpoint`, `Run As`, `Debug As`, `Compare With`, `Replace With`, `Restore from Local History...`, `Properties`, and `JForge analysis`. The `JForge analysis` sub-menu is open, showing `Check and Simulate`, `Simulate`, and `Check`.
- Bottom Panel (Output Console):** Shows messages: `Simulation of insert: found a trace of an execution` and `Compliance check of insert: no counter examples found`.
- Status Bar:** Displays the current cursor position: `edu.mit.csail.sdg.annotations.test.IntTree.insert(Node z) : void - edu.mit.csail.sdg.annotations/src`.

plugin implemented by Kuat Yessenov; underlying Forge analyzer by Greg Dennis

code synthesis



WebAlloy [Chang 2009]
data model + policy in Alloy
generate entire site
gateway enforces policy

results

3 real websites in ~100
lines of Alloy each

declarative configuration

declarative configuration

*'We don't need hackers
to break systems because
they're falling apart
by themselves'*

Peter Neumann

*bad configurations responsible for
80% of USAF vulnerabilities*

NSA

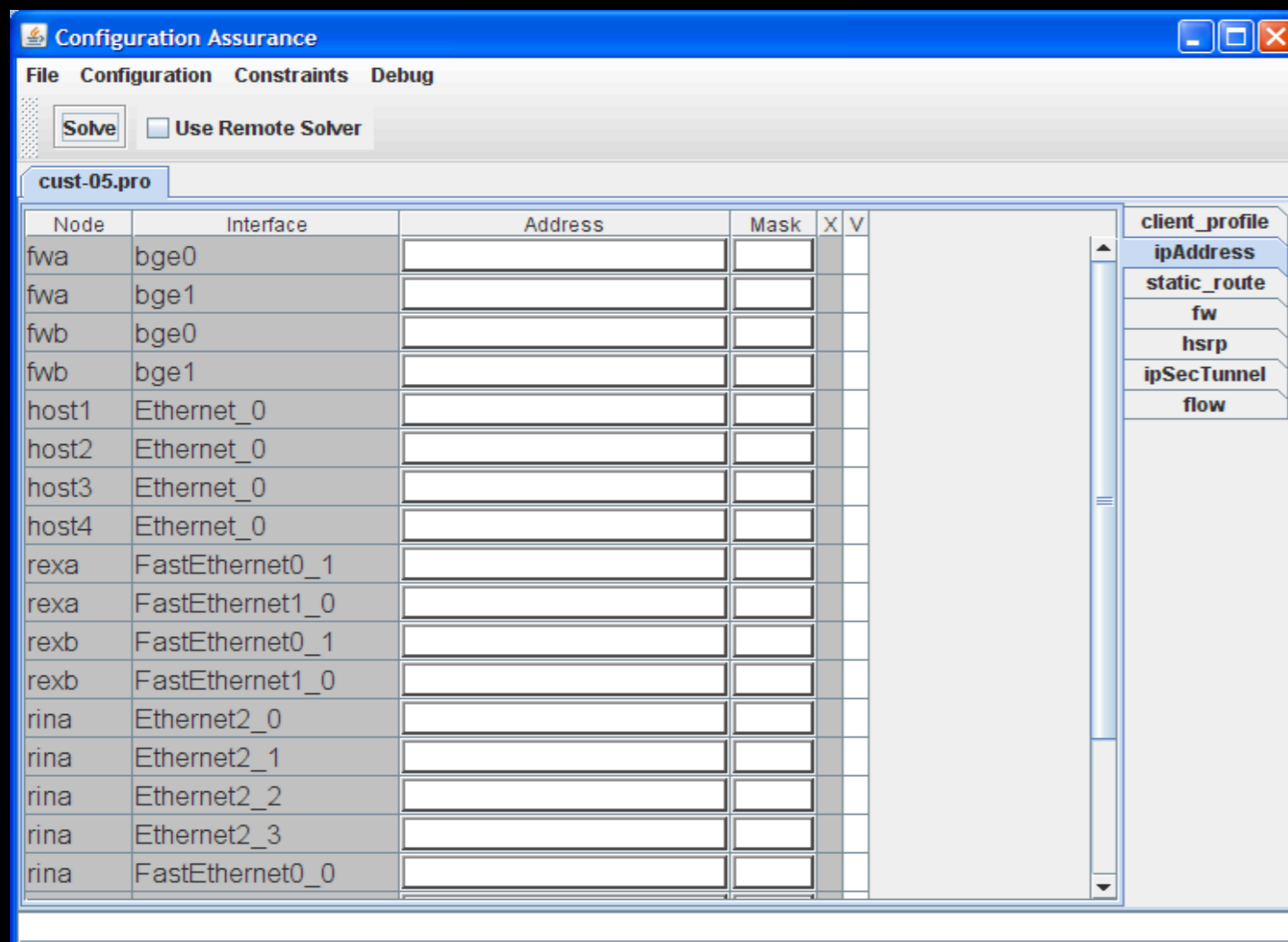
declarative configuration

ConfigAssure

[Telcordia, 2008]

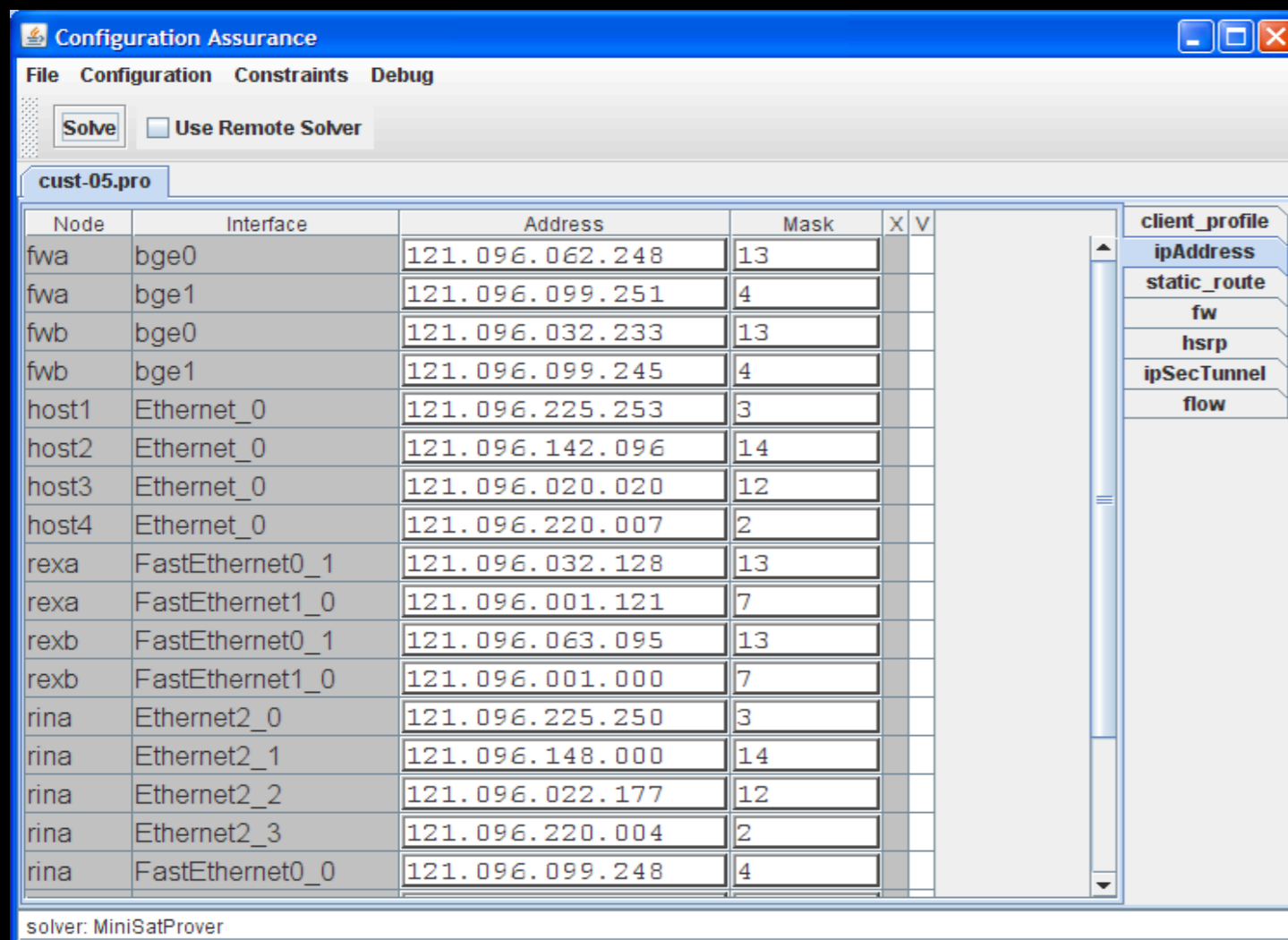
*use Kodkod to find
network configuration
that meets rules & goals*

declarative configuration



ConfigAssure
[Telcordia, 2008]
use Kodkod to find
network configuration
that meets rules & goals

declarative configuration



The screenshot shows the Configuration Assurance application window. The main area displays a table with columns for Node, Interface, Address, Mask, X, and V. The table lists configurations for various nodes like fwa, fwb, host1, host2, host3, host4, rexa, rexb, rina, and rina. The status columns X and V are currently empty. A sidebar on the right lists configuration categories: client_profile, ipAddress, static_route, fw, hsrp, ipSecTunnel, and flow. The status bar at the bottom indicates the solver is MiniSatProver.

Node	Interface	Address	Mask	X	V
fwa	bge0	121.096.062.248	13		
fwa	bge1	121.096.099.251	4		
fwb	bge0	121.096.032.233	13		
fwb	bge1	121.096.099.245	4		
host1	Ethernet_0	121.096.225.253	3		
host2	Ethernet_0	121.096.142.096	14		
host3	Ethernet_0	121.096.020.020	12		
host4	Ethernet_0	121.096.220.007	2		
rexa	FastEthernet0_1	121.096.032.128	13		
rexa	FastEthernet1_0	121.096.001.121	7		
rexb	FastEthernet0_1	121.096.063.095	13		
rexb	FastEthernet1_0	121.096.001.000	7		
rina	Ethernet2_0	121.096.225.250	3		
rina	Ethernet2_1	121.096.148.000	14		
rina	Ethernet2_2	121.096.022.177	12		
rina	Ethernet2_3	121.096.220.004	2		
rina	FastEthernet0_0	121.096.099.248	4		

ConfigAssure
[Telcordia, 2008]
use Kodkod to find
network configuration
that meets rules & goals

some other uses of Alloy

bytecode security checking
Reynolds

ontology analysis
Dong

mixed programming
Rayside et al

architectural styles
Kim & Garlan, Sullivan

product-line test generation
Uzuncaova et al

HOL to Alloy
Blanchette

memory model analysis
Torlak

UML to Alloy
Dingel & Zito, Bordbar, Roquette

security policy analysis
Krishnamurthi & Fislser, Hassan

multiobjective optimization
Rayside & Estler

a new direction

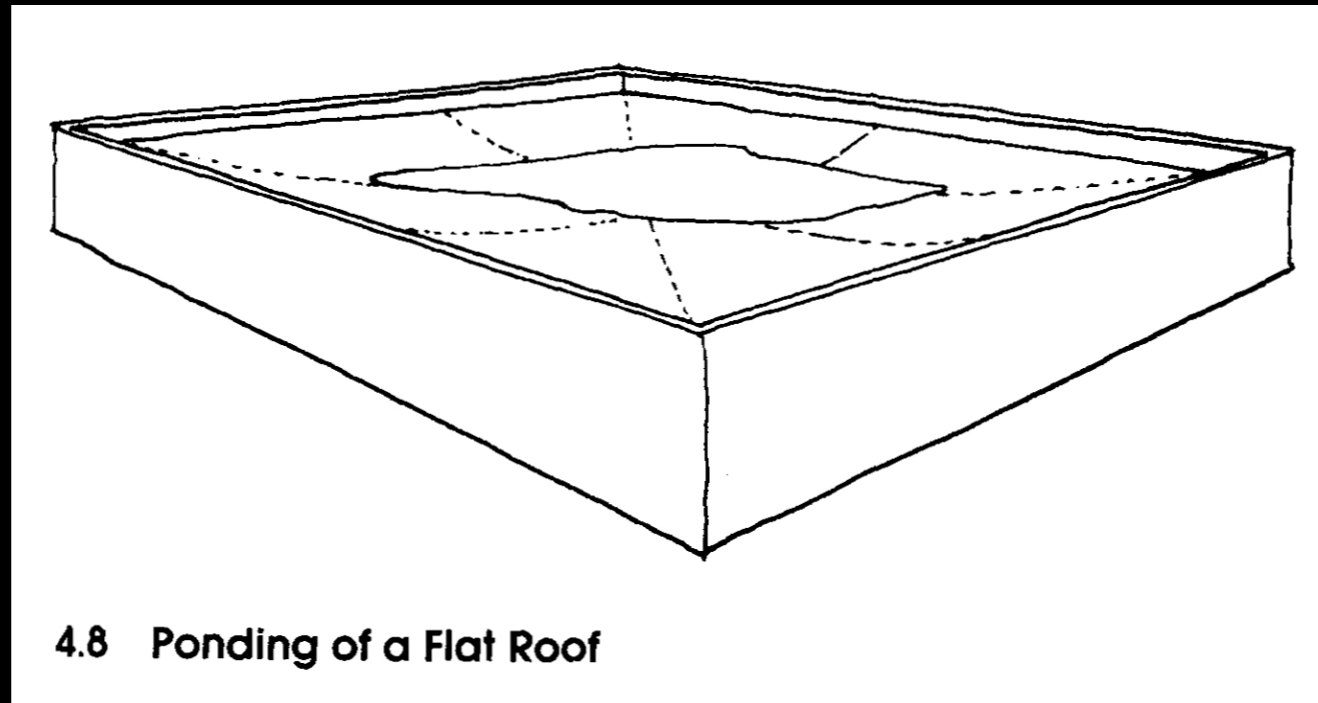
kemper arena, kansas city, 2007



kemper arena, 1979

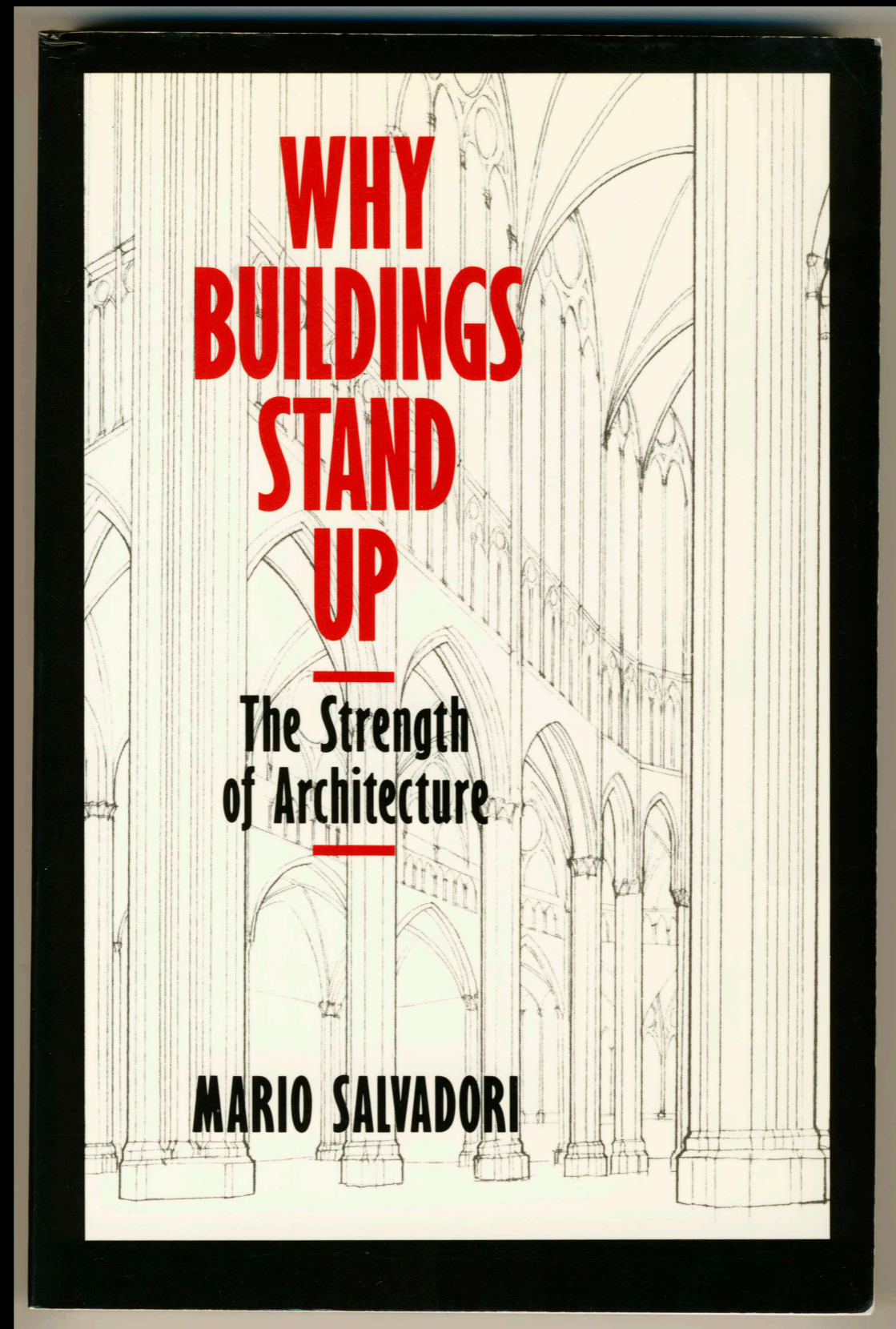


what happened?

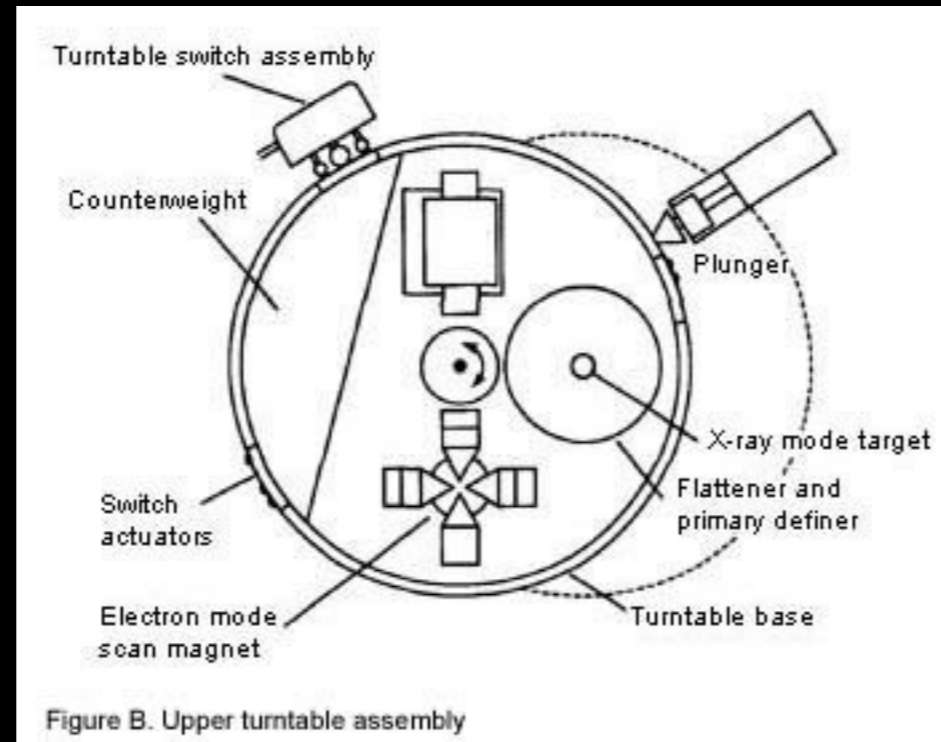


Levy & Salvadori, *Why Buildings Fall Down*

failure = flawed success story



Therac 25



AECL fault tree analysis (1983)
did not include software

$P(\text{computer selects wrong energy}) = 10^{-11}$

Leveson & Turner (1993)
race conditions, lack of interlocks, etc

dependability arguments

solve the real problem

not just software: operators & peripherals too
top-to-bottom, end-to-end

be realistic

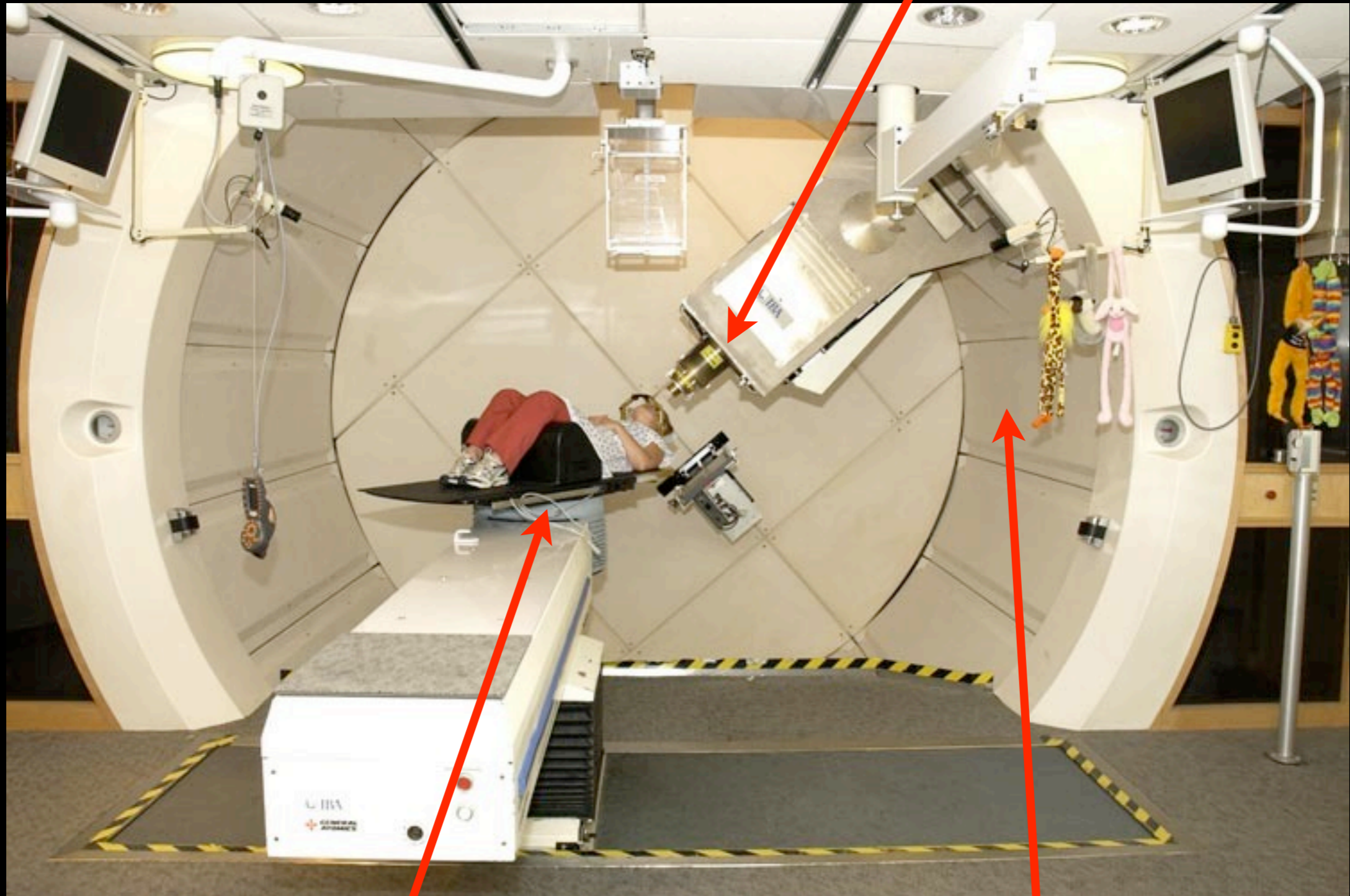
not 'full correctness': prioritize requirements

don't be a masochist

use a safe language

design to localize critical properties

beam line



robotic couch

rotating gantry

a dependence diagram

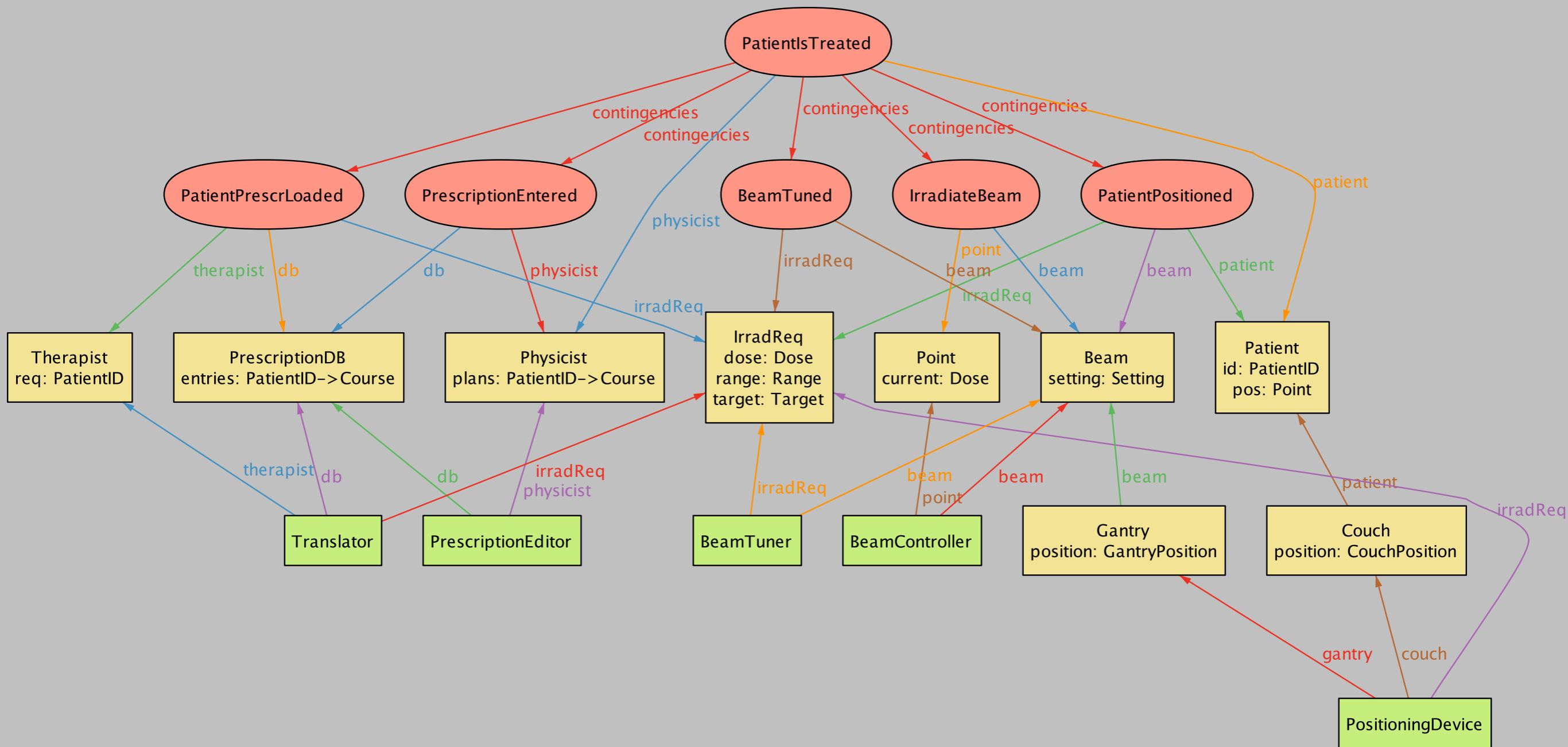


diagram by Eunsuk Kang

reflections

the future of modelling

the future of modelling

it will happen, really
domain-specific approaches
Rails is a modelling language

the future of modelling

it will happen, really
domain-specific approaches
Rails is a modelling language

animation now, analysis later
exploration, test case generation
then maybe checking

the future of modelling

it will happen, really
domain-specific approaches
Rails is a modelling language

animation now, analysis later
exploration, test case generation
then maybe checking

relational modelling
non-transparent ORMs must go

what I learnt from this project

what I learnt from this project

what worked

drive research by case studies

robust tools: x10 cost, but worth it

semantics is a verb: helps keep it simple

what I learnt from this project

what worked

drive research by case studies

robust tools: x10 cost, but worth it

semantics is a verb: helps keep it simple

what didn't work

language design is a hard sell

software engineering is a really hard sell

easy to get new idea funded once popular

what I learnt from this project

what worked

drive research by case studies

robust tools: x10 cost, but worth it

semantics is a verb: helps keep it simple

what didn't work

language design is a hard sell

software engineering is a really hard sell

easy to get new idea funded once popular

a mistake?

keeping Alloy general & pure

limited impact, except in teaching

for more info

<http://sdg.csail.mit.edu>

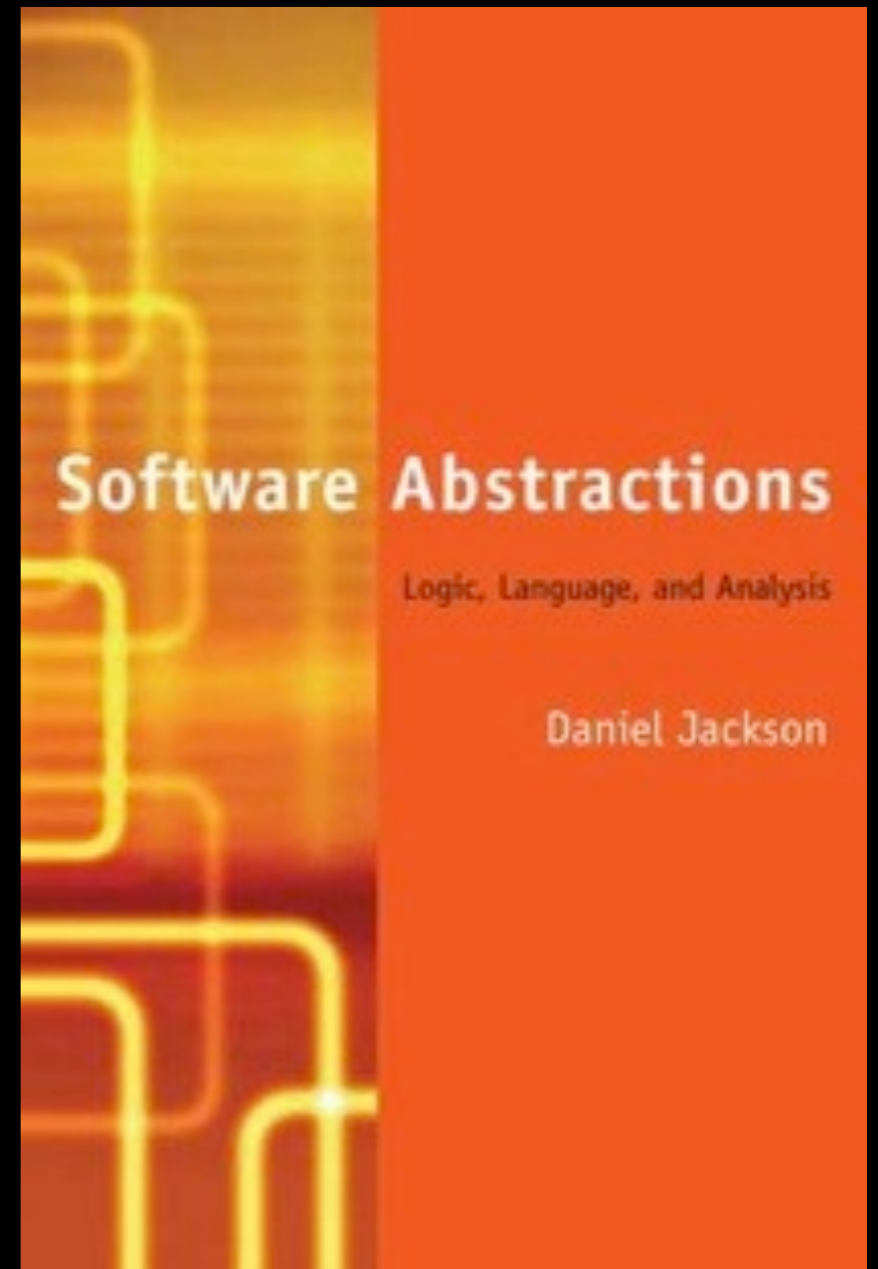
other fun things my group does

<http://alloy.mit.edu>

community, downloads, tutorial

<http://softwareabstractions.org>

sample chapters, models



acknowledgments



content of this talk especially

Alloy 4: Felix Chang

Kodkod engine: Emina Torlak

code analysis: Greg Dennis, Mana Taghdiri, Mandana Vaziri

Alloy 3: Ilya Shlyakhter, Manu Sridharan

not shown in photos: Sarfraz Khurshid, Mandana Vaziri, Manu Sridharan, Ilya Shlyakhter