

The Alloyed Joys of Software Engineering Research

Daniel Jackson

Keynote · ICSE 2017 · Buenos Aires





Alloy

alloy's cultural origins



Oxford, home of Z



Pittsburgh, home of SMV

lightweight formal methods

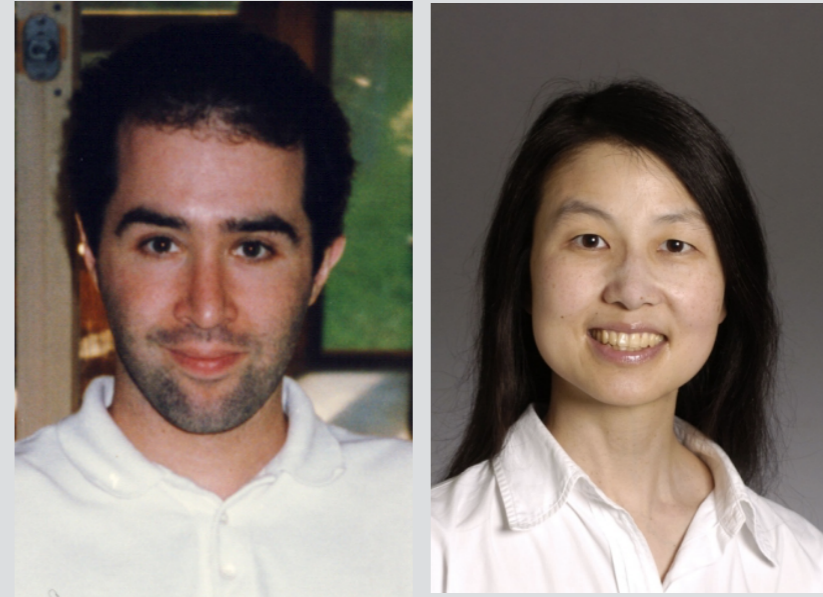
LIGHTWEIGHT FORMAL METHODS

Daniel Jackson and Jeannette Wing,
Carnegie Mellon University

Many benefits promised by formal methods are shared with other approaches. The precision of mathematical thinking relies not on formality but on careful use of mathematical notions. You don't need to know Z to think about sets and functions. Likewise, the linguistic advantages of a formal notation rely more on syntax than semantics.

Mechanical analysis, in contrast, is a benefit unique to formal approaches. An engineer's sketch can communicate ideas to other engineers, but only a detailed plan can be rigorously examined for flaws. Informal methods often provide some analysis, but since their notations are generally incapable of expressing behavior, the results of the analysis bear only on the properties of the artifact's description, not on the properties of the artifact itself.

IEEE Computer, 1996



traditional FM

full model of behavior
analysis to show no bugs

lightweight FM

model of critical aspect
analysis to find bugs



Alloy Analyzer 4.2 (build date: 2012-02-28 12:29 EST)

Executing "Run run\$1"

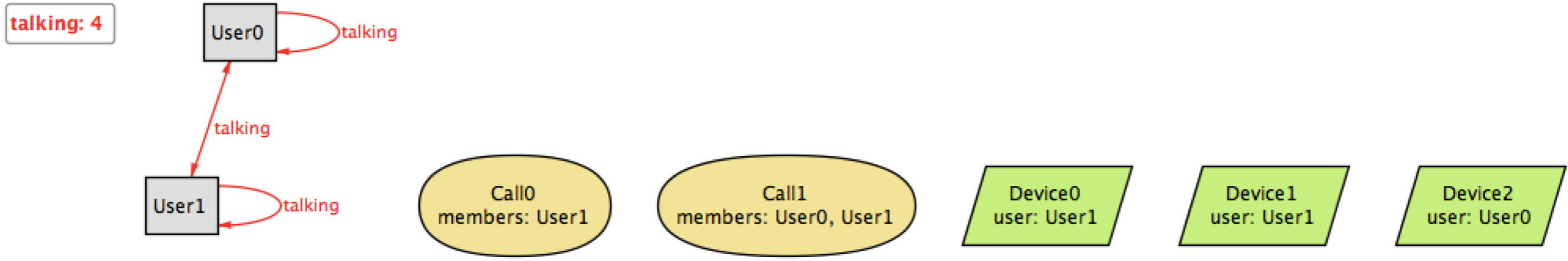
Solver=minisat(jni) Bitwidth=0 MaxSeq=0 SkolemDepth=1 Symmetry=20
403 vars. 36 primary vars. 758 clauses. 68ms.
Instance found. Predicate is consistent. 14ms.

```
sig Device {user: lone User}
sig Call {members: set User}
sig User {talking: set User}

fact {
  all u: User | u.talking = {u': User | some c: Call | u + u' in c.members}
  all u: User | some u.talking implies some user.u
}

run {
  #talking > 2
}
```

Line 1, Column 1



alloy timeline

version

language

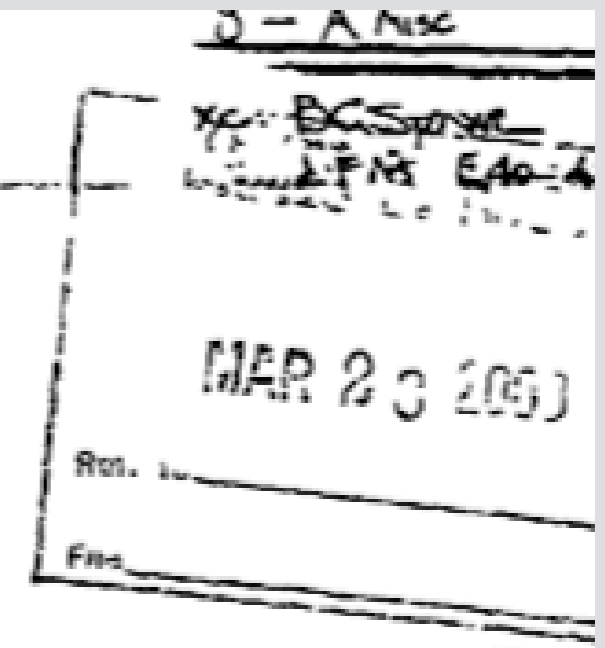
analysis

sample case study

Nitpick (1995)	relational calculus subset of Z	relation enumeration	IPv6 routing
Alloy 1 (1999)	+ navigation exps quantifiers	WalkSAT, Davis Putnam	intentional naming
Alloy 2 (2001)	+ non-binary relations signatures	Chaff, Berkmin symmetry, sharing	Unison filesync
Alloy 3 (2004)	+ subtyping overloading	atomization (bad)	Mondex smartcard
Alloy 4 (2007)	+ meta, sequences arithmetic	bounds better sharing	flash filesystem

the alloy constraint analyzer

President Charles M. Vest
Massachusetts Institute of Technology
Building 5
77 Massachusetts Avenue
Cambridge, MA 02139-4307



Re: Notice of Trademark Infringement

Dear President Vest:

We have just learned that software developers at MIT have named a software program "Alcoa". This unauthorized usage of the ALCOA trademark on software on the MIT internet site is an infringement of the trademark rights of Alcoa Inc (Alcoa).

Alcoa has been using the trademark and service mark ALCOA on a wide variety of goods and services throughout the world since 1926. Through extensive sales and advertising our trademark and trade name ALCOA is famous worldwide. It is well associated with metal alloys.

The software developers are knowingly using the **ALCOA** trademark and trade name. The last question on the corresponding FAQ page is:

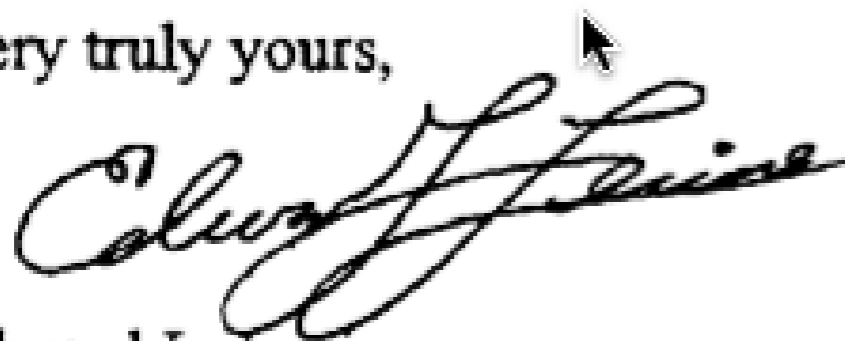
Is Alcoa endorsed by the Alcoa Corporation?

No, we just liked the name. The language, like an alloy, obtains its strength from a combination of ingredients, and, like many alloys, is lightweight. Running the tool is a bit like melting a metal: it heats things up (and sometime makes your structures fall apart :-).

You may also be aware that Alcoa is currently a participant in the Leaders for Manufacturing Program sponsored by the Sloan School of Management and the School of Engineering.

Thank you for your prompt attention to this matter.

Very truly yours,

A handwritten signature in cursive script that reads "Edward L. Levine". A mouse cursor is visible over the signature.

Edward L. Levine

Director – Intellectual Property Law

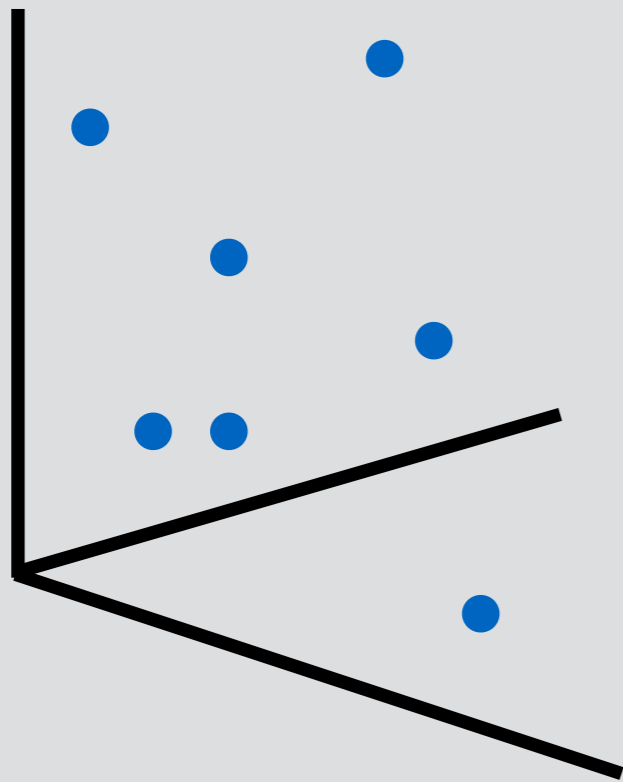
☎ (724)337-2759

FAX (724)337-5959

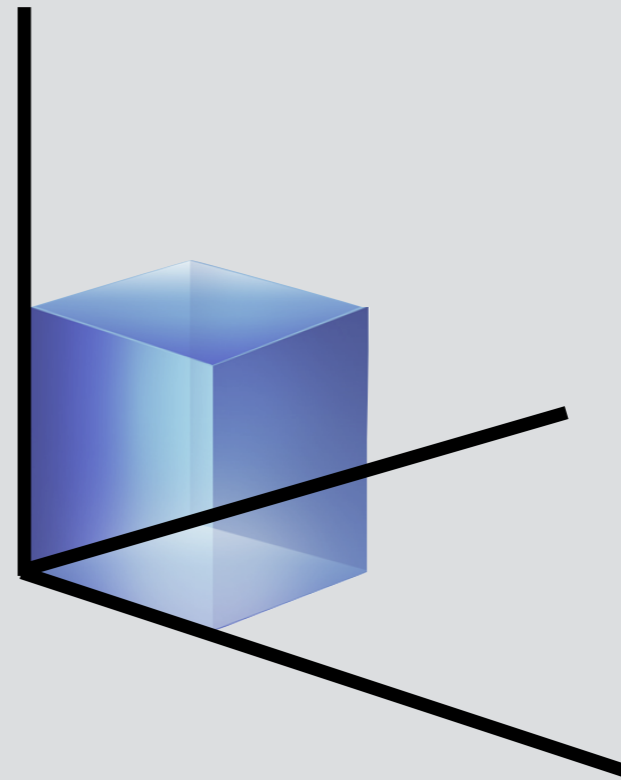
5

ideas

all small tests



traditional testing



bounded analysis

5 users, calls, devices
 2^{25} user-call, user-device relations
so $2^{50} = 10^{15}$ states

a signature style

sig Call {members: **set** User}

$Call = \mathbb{P} (id: CallId \times members: \mathbb{P} User)$
 $User = \mathbb{P} (id: UserId \times talking: \mathbb{P} User)$

traditional interpretation

$Call, User: \mathbb{P} Univ$
 $members: Call \leftrightarrow User$

Alloy interpretation

all $c, c': Call$ {**no** $c.members$ & $c'.members$ }

$\forall c, c' \in \mathbb{P} (id: CallId \times members: \mathbb{P} User) \mid \dots$

higher order
quantification:
ouch!

$\exists members: Call \leftrightarrow User \mid$
 $\forall c, c' \in Call \mid \dots$

first order
quantification:
solve with SAT

everything's a relation

sig Call {members: **set** User}

sig User {talking: **set** User}

some expressions:

c.members
members.u
u.talking
c.members.talking
u.talking = u'

navigation: dot is
just join, not
overloaded

no syntax
difference: fun vs
relation

no undefined
value, follows
Parnas

getting satisfaction

```
sig User {talking: set User}
```

```
check {no u: User | u in u.talking}
```

	<i>U0</i>	<i>U1</i>	<i>U2</i>
<i>U0</i>	1	0	0
<i>U1</i>	0	1	0
<i>U2</i>	0	0	0

u

{(U0)}

	<i>U0</i>	<i>U1</i>	<i>U2</i>
<i>U0</i>	0	1	0
<i>U1</i>	1	0	0
<i>U2</i>	0	0	0

talking

{U0,U1), (U1, U0)}

<i>U0</i>	0
<i>U1</i>	1
<i>U2</i>	0

u.talking

{(U0), (U1)}

getting satisfaction

add symmetry
breaking
predicates too

sig User {talking: **set** User}

check {**no** u: User | u **in** u.talking}

	<i>U0</i>	<i>U1</i>	<i>U2</i>			
<i>U0</i>	u0	t00	t01	t02	<i>U0</i>	$(u0 \wedge t00) \vee (u1 \wedge t10) \vee (u2 \wedge t20)$
<i>U1</i>	u1	t10	t11	t12	<i>U1</i>	$(u1 \wedge t01) \vee (u1 \wedge t11) \vee (u2 \wedge t21)$
<i>U2</i>	u2	t20	t21	t22	<i>U2</i>	$(u0 \wedge t02) \vee (u1 \wedge t12) \vee (u2 \wedge t22)$
	u	talking				u.talking

$u0 \Rightarrow (u0 \wedge t00) \vee (u1 \wedge t10) \vee (u2 \wedge t20) \quad \wedge$
 $u1 \Rightarrow (u1 \wedge t01) \vee (u1 \wedge t11) \vee (u2 \wedge t21) \quad \wedge$
 $u2 \Rightarrow (u0 \wedge t02) \vee (u1 \wedge t12) \vee (u2 \wedge t22)$

some u: User | u **in** u.talking

roll your own idiom

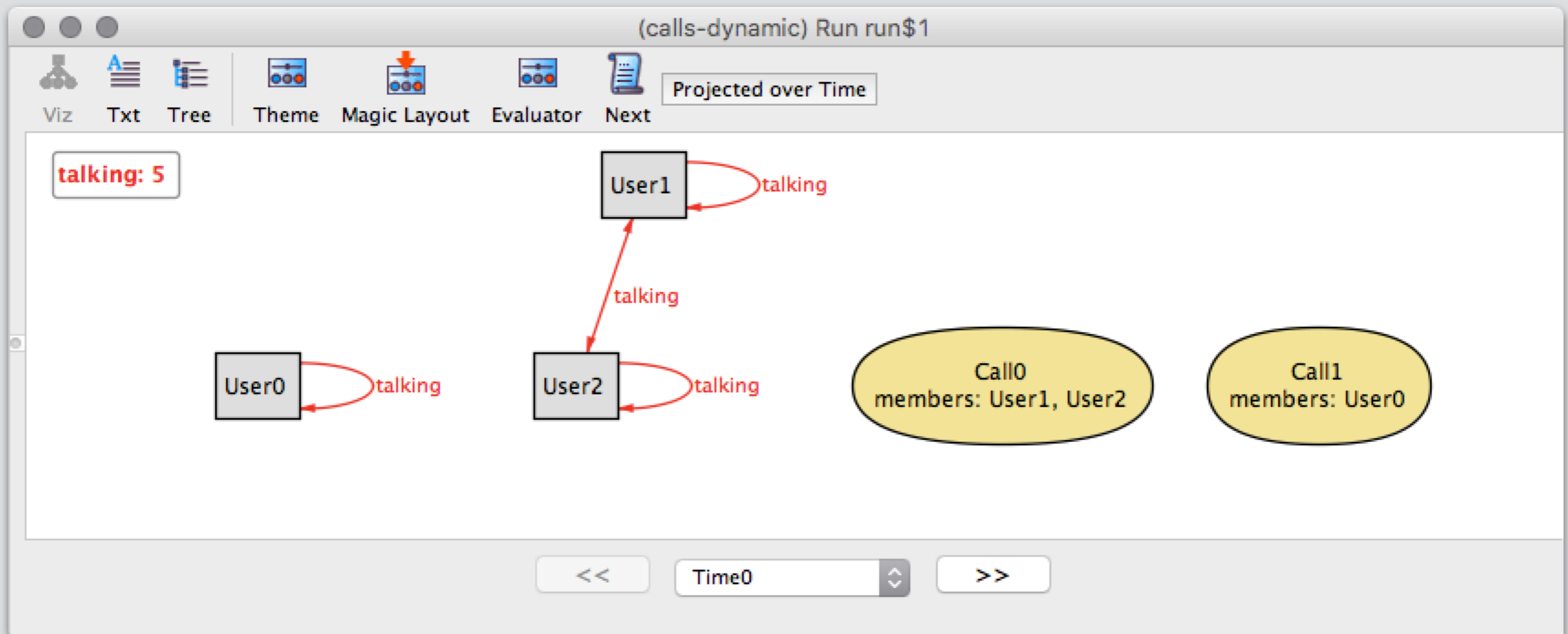
```
open util/ordering[Time]
```

```
sig Time {}
```

```
sig Call {members: User -> Time}
```

```
sig User {talking: User -> Time}
```

```
fact { all t: Time | let m = members.t | talking.t = ~m.m }
```



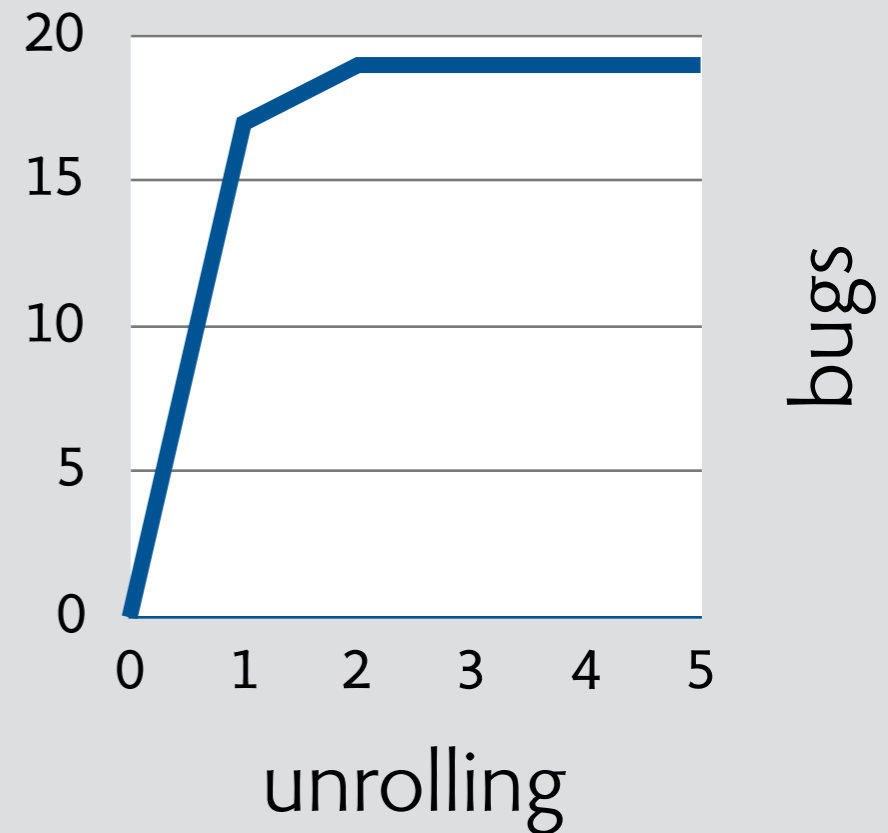
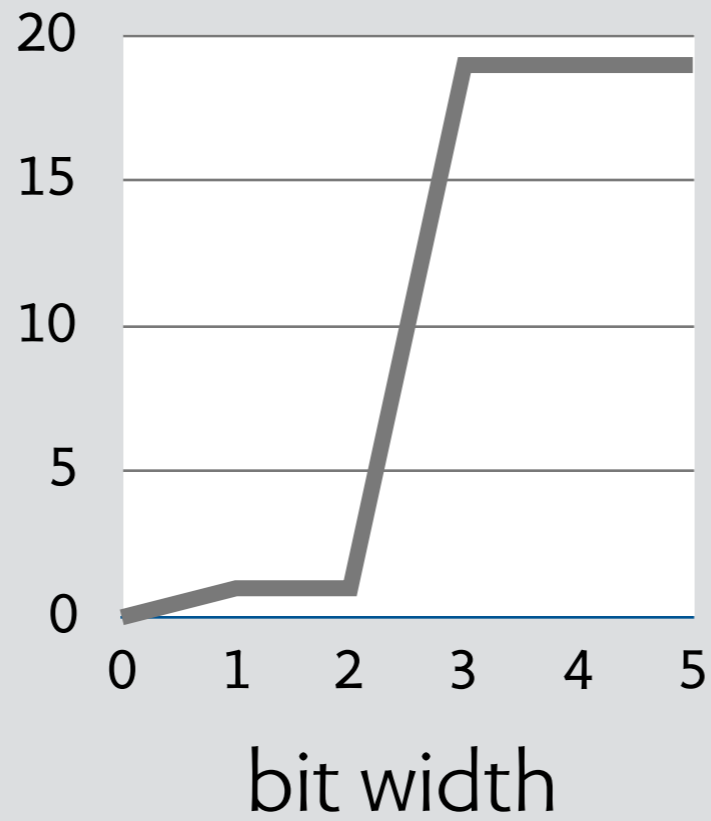
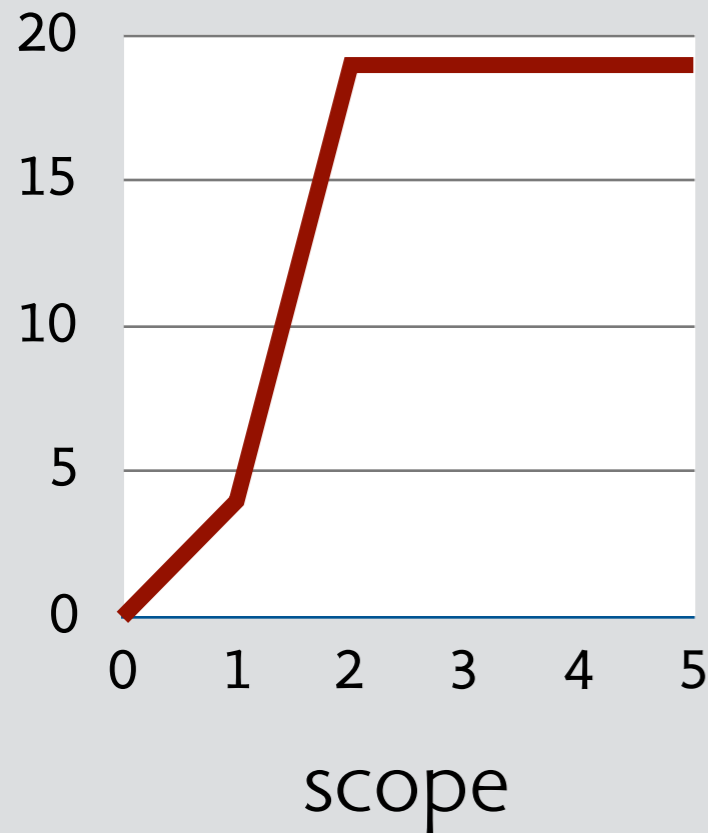
4

outcomes

but does it work? tell us the truth!



are small scopes enough?



analysis of KOA voting code

19 methods violating specs

how many bugs found in scope of k?

[Greg Dennis, 2008]

most bugs in small scopes?

yes, but two caveats

integers are nasty: 'special' semantics

trace length must be set higher

why traces are tricky

in scope 5, call-user has ≤ 25 pairs

can check an operation on 2^{25} pre-states

but if initially empty, 25 steps to populate?

is first order enough?

converting Z (eg) to Alloy
generally straightforward

minimization may be OK

send packet to nearest neighbor?
easy: just say no shorter option

synthesis is higher order

find a program without bugs

$\exists p: \text{Program} \mid \forall s: \text{State} \mid S(p,s)$

this motivated Alloy* [Milicevic+]



Mondex smart card system
NatWest, Oxford U., Logica
[Ramananandro]



Tokeneer project
Praxis/NSA
50pp Z, 1200 lines Alloy
[Eunsuk Kang]

was purity a good idea?

on the one hand

breadth of domains
nice translation target
good for teaching logic

on the other hand

dynamic idioms are complex
frame conditions annoying

Just this year, students used Alloy for a broad range of unexpected topics including:

- checking theorems about groups
- generating Feynman Diagrams
- modeling Facebook privacy

Tim Nelson, talking about his Brown course, Logic for Systems

is declarative spec easy?

```
open util/ordering[Time]
sig Time {}

sig Call {members: User -> Time}
sig User {talking: User -> Time}

fact {
  all t: Time | members.t in Call lone -> User
  all t: Time | let m = members.t | talking.t = ~m.m
}

pred add [u: User, c: Call, t, t': Time] {
  members.t' = members.t + c->u
  u not in u.talking.t'
}

run add
```

don't end up
talking to
yourself

let's see what happens

The screenshot shows the Alloy Analyzer 4.2 interface. The title bar indicates the file path: `/Users/dnj/Filestore/Talks/fse17/models/add-call.als`. The menu bar includes `New`, `Open`, `Reload`, `Save`, `Execute`, and `Show`. The main editor displays the following Alloy code:

```
open util/ordering[Time]
sig Time {}

sig Call {members: User -> Time}
sig User {talking: User -> Time}

fact {
  all t: Time | members.t in Call lone -> User
  all t: Time | let m = members.t | talking.t = ~m.m
}

pred add [u: User, c: Call, t, t': Time] {
  members.t' = members.t + c->u
  u not in u.talking.t'
}

run add
```

The execution output on the right reads:

Alloy Analyzer 4.2 (build date: 2012-02-28 12:29 EST)

Executing "Run add"

Solver=minisatprover(jni) Bitwidth=0 MaxSeq=0 SkolemD
928 vars. 87 primary vars. 1636 clauses. 100ms.
No instance found. Predicate may be inconsistent. 13ms.
Core contains 2 top-level formulas. 27ms.

A yellow callout box points to the line `all t: Time | let m = members.t | talking.t = ~m.m` in the code, containing the text: **this definition makes everyone self talkers**.

At the bottom left, the status bar shows: `Line 12, Column 42`.

so what's the story?

declarative specification

can be magical
often very succinct
nice separation of concerns

can be maddening
harder to learn than I knew
even harder to debug
unsat core not enough

20

projects

extending Alloy

expressiveness

Alloy*: higher-order quantifiers [Milicevic+]

temporal constructs

DynAlloy [Frias+], [Macedo+]

better scenarios

target instances [Cunha+]

Aluminum: minimal instances [Nelson+]

performance

separating configurations [Macedo+]

exploit previous analyses: Titanium [Bagheri+]

translation optimizations [Marinov+]

platforms

Eclipse [LeBerre], web client [Cunha+]

tools built on Alloy

code analysis

Forge [Dennis+], TACO [Galeotti+]

architecture

design space exploration [Bagheri+]

architectural styles [Garlan+]

security

Margrave: policy analysis [Fisler+]

Poirot: vulnerabilities due to platform choice [Kang+]

software defined networking

Flowlog [Nelson+]

checking theorems

Nitpick for Isabelle [Blanchette]

a small sample of
amazing tools
people have built

some favorite applications of Alloy

web security [Akhawe+]

reusable model of web platform

found 2 known and 3 new vulnerabilities

in all cases,
it's more than
finding bugs

networking [Zave]

showed Chord violates all its invariants

designed a new version + invariant

dependability cases [UW PLSE]

end-to-end analysis of neutron therapy

memory models [Torlak+; Wickerson+, Dodds+, Lustig+]

validate and develop new memory models

3

lessons

invest in your tool

```
sig User {device: Device, calls: set Call}{  
  no device implies no calls  
  this in calls.users  
}
```

look Ma, no
semicolons!

before she went
to jail

Node Color Palette:

Martha

Use original atom names:

Edge Color Palette:

Classic

Font Size:

12

Hide private sigs/relations:

Hide meta sigs/relations:

be nice (and objective)

a stupid thing I wrote:

"[In Z,] since declared sets cannot be used in subsequent declarations, simple multiplicity constraints must be written as additional textual formulas. The resulting specification is cluttered and unnatural."

understandably aggrieved reviewer:

I suppose that I shouldn't be irritated by the final sentence in this quote, but I am: what is the measure of what is natural? Anyway, the whole thing is complete tosh...

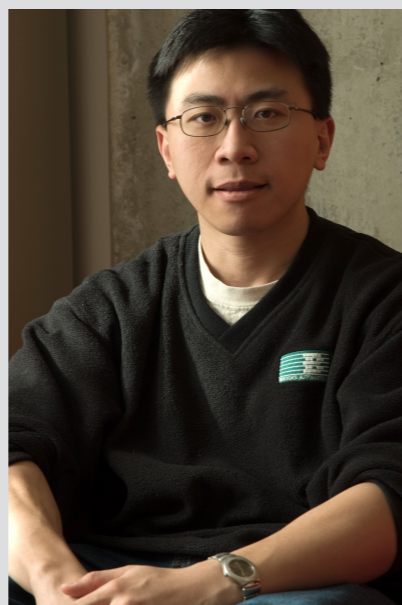
tosh

/täSH/ 

noun **BRITISH** *informal*

rubbish; nonsense.
"it's sentimental tosh"

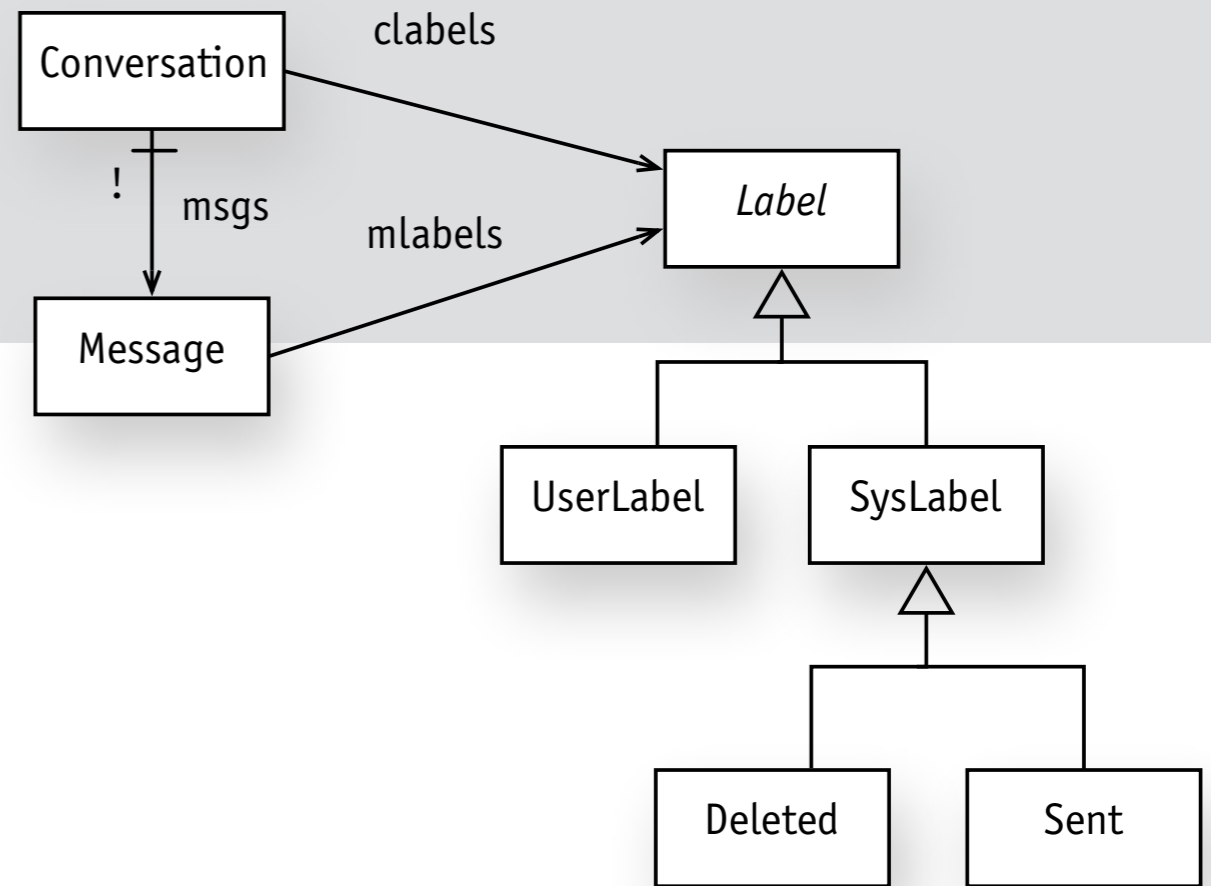
get lucky!



3

thoughts

human factors



more emphasis needed
especially in formal methods

what I eventually figured out
abstraction is really hard

most programmers can't draw an ER diagram
usual educational approaches don't work

what if I'd studied this 20 years ago?

might not have changed Alloy

but might have changed my research direction?

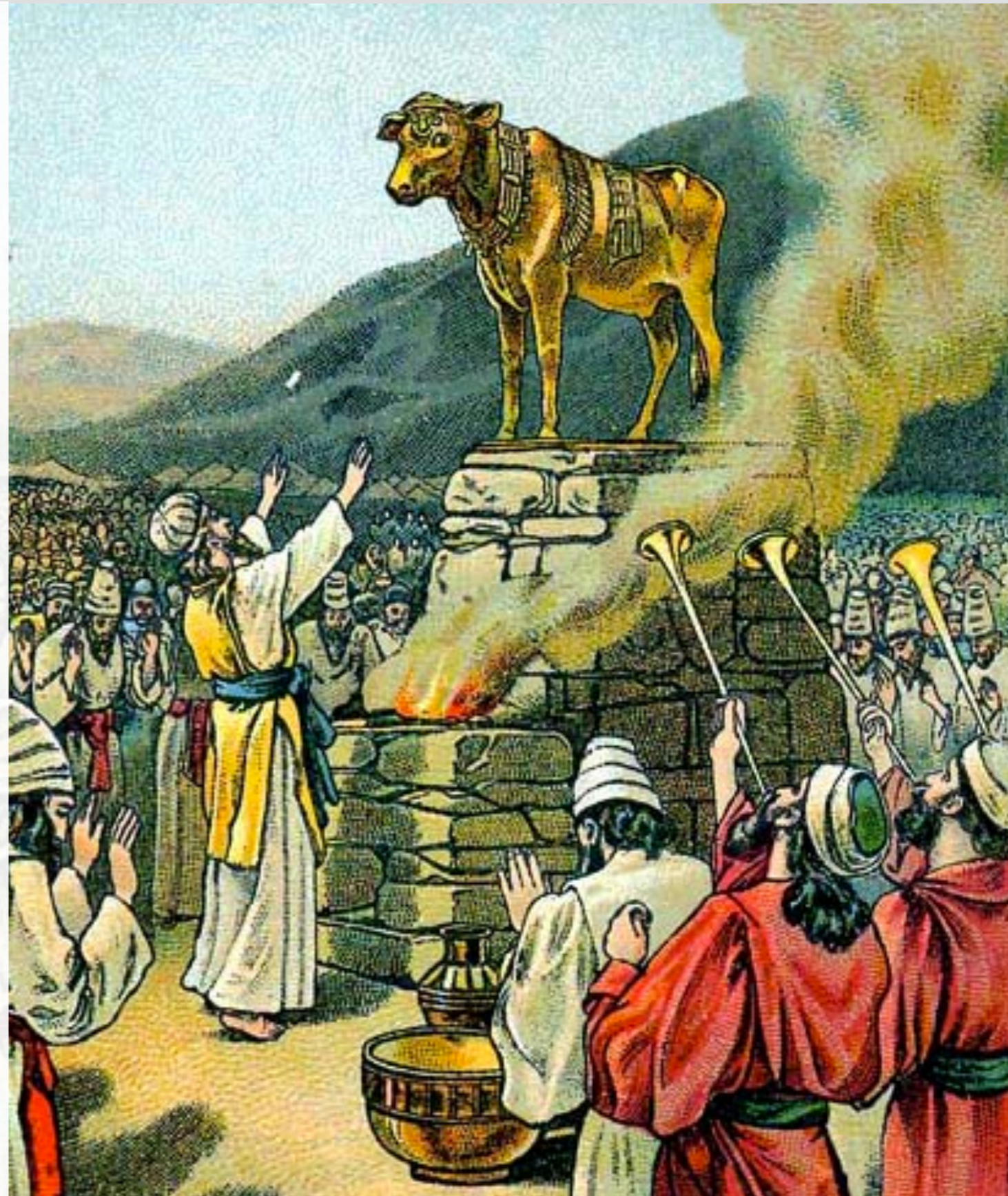
on empiricism

empirical research
exciting & powerful

empirical validation
as sole arbiter: a mistake

has not
upped field's reputation
resolved old disputes
made papers compelling

but has
inhibited novel work
devalued design research



serving industry?

industrial collaborations provide

source of new problems
deeper understanding of old problems
new approaches (XP, agile, etc)
opportunity to try research ideas

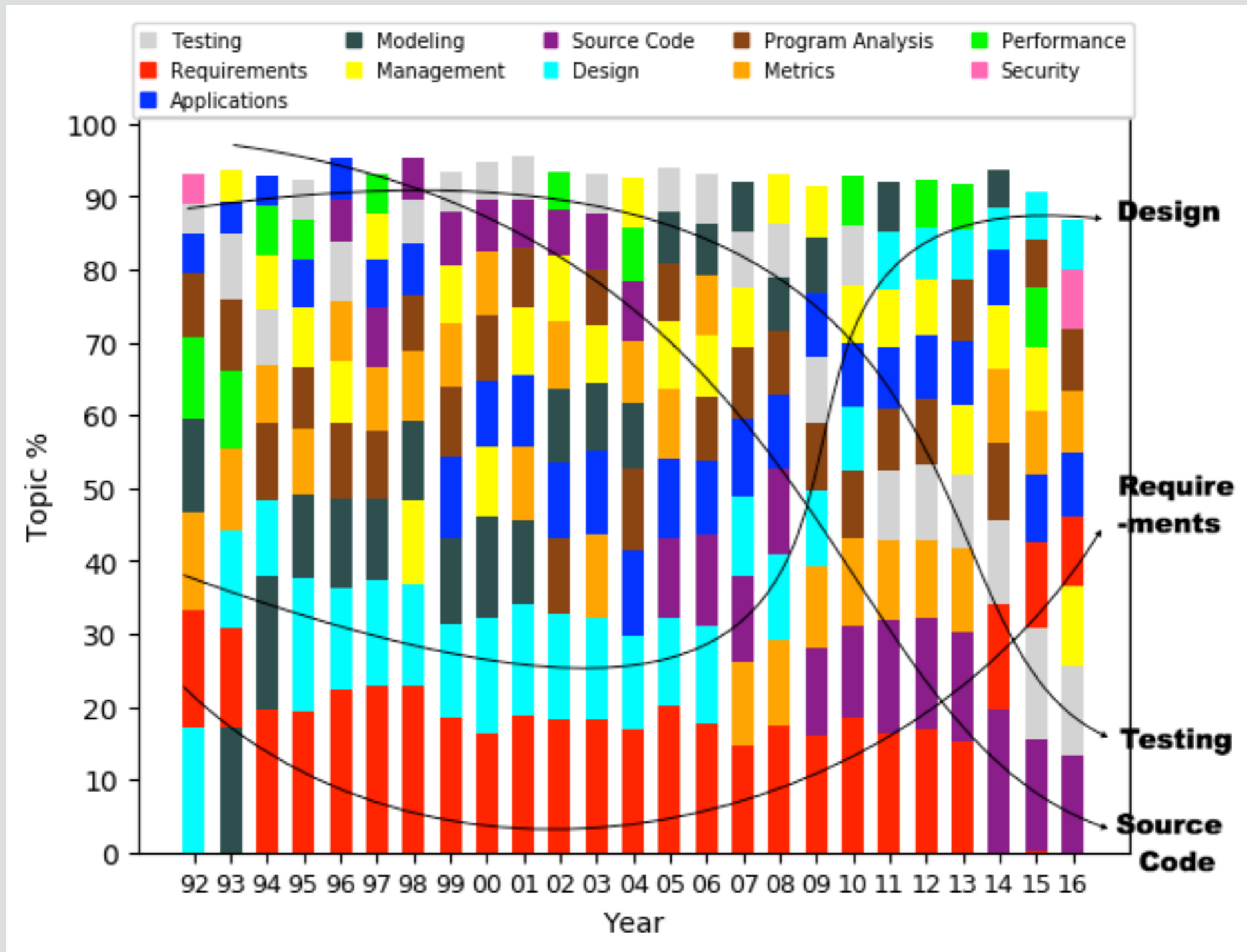
but increasingly seems that

SE researchers see their role as serving industry
addressing immediate problems

this leads to

overemphasis on code & test
lack of long-term thinking

a consequence



from Mathew, Agrawal & Menzies

more info at <http://alloy.mit.edu>

[about](#) | [community](#) | [download](#) | [documentation](#) | [book](#) | [applications](#) | [people](#) | [thanks](#)

alloy: a language & tool for relational models

about alloy

Alloy is a language for describing structures and a tool for exploring them. It has been used in a wide range of [applications](#) from finding holes in security mechanisms to designing telephone switching networks.

An Alloy model is a collection of constraints that describes (implicitly) a set of structures, for example: all the possible security configurations of a web application, or all the possible topologies of a switching network. Alloy's tool, the [Alloy Analyzer](#), is a solver that takes the constraints of a model and finds structures that satisfy them. It can be used both to explore the model by generating sample structures, and to check properties of the model by generating counterexamples. Structures are displayed graphically, and their appearance can be customized for the domain at hand.

At its core, the Alloy language is a simple but expressive logic based on the notion of relations, and was inspired by the Z specification language and Tarski's relational calculus. Alloy's syntax is designed to make it easy to build models incrementally, and was influenced by modeling languages (such as the object models of OMT and UML). Novel features of Alloy include a rich subtype facility for factoring out common features and a uniform and powerful syntax for navigation expressions.

The Alloy Analyzer works by reduction to SAT. Version 4 was a complete rewrite that included [Kodkod](#), a new model finding engine that optimizes the reduction, and a new front end.

contact us!

news

A [Japanese translation](#) of book published!

Revised edition of book now out!
Available from [MIT Press](#).

