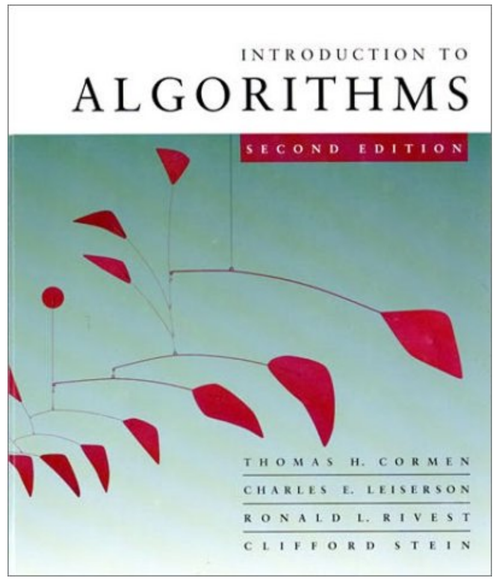
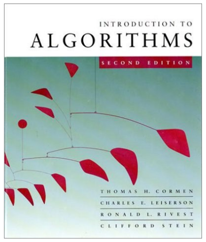
Introduction to Algorithms

6.046J/18.401J



Lecture 1

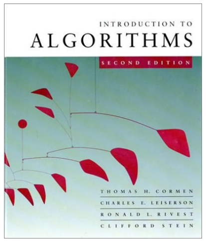
Prof. Piotr Indyk



Welcome to *Introduction to Algorithms*, Spring 2003

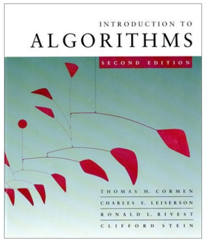
Handouts

1. Course Information
2. Calendar (also on the web)
3. PS0 (due tomorrow, 8 am)



Course information

1. Staff
2. Prerequisites
3. Lectures
4. Recitations
5. Handouts
6. Textbook (CLRS)
7. Website
8. Extra help (HKN)
9. Registration (by 8 am)
10. Problem sets
11. Describing algorithms
12. Grading policy
13. Collaboration policy

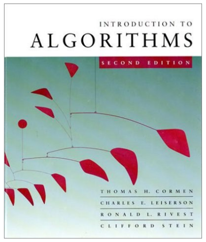


Analysis of algorithms

The theoretical study of computer-program performance and resource usage.

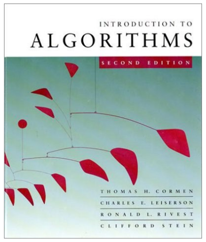
Also important:

- modularity
- maintainability
- functionality
- robustness
- user-friendliness
- programmer time
- simplicity
- extensibility
- reliability



Why study algorithms and performance?

- Performance often draws the line between what is feasible and what is impossible.
- Algorithmic mathematics provides a ***language*** for talking about program behavior.
- Speed is fun!



The problem of sorting

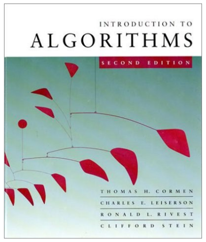
Input: sequence $\langle a_1, a_2, \ldots, a_n \rangle$ of numbers.

Output: permutation $\langle a'_1, a'_2, \ldots, a'_n \rangle$ such that $a'_1 \leq a'_2 \leq \cdots \leq a'_n$.

Example:

Input: 8 2 4 9 3 6

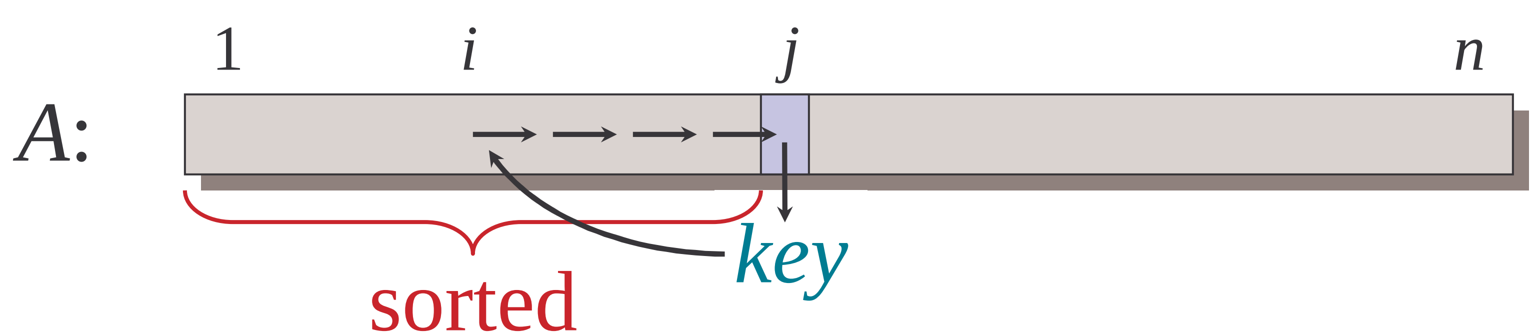
Output: 2 3 4 6 8 9

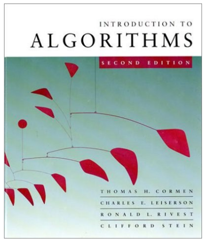


Insertion sort

“pseudocode”

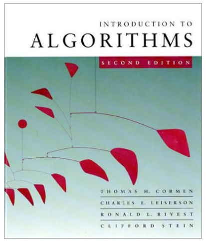
```
INSERTION-SORT (A, n)        ◁ A[1 . . n]
    for j ← 2 to n
        do key ← A[ j]
           i ← j – 1
           while i > 0 and A[i] > key
               do A[i+1] ← A[i]
                  i ← i – 1
           A[i+1] = key
```





Example of insertion sort

8 **2** 4 9 3 6



Example of insertion sort

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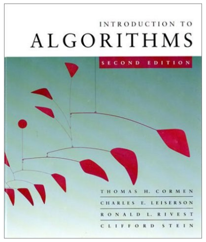
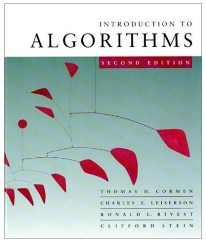


Example of insertion sort

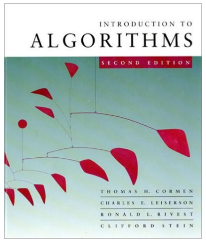
8	**2**	4	9	3	6
2	8	**4**	9	3	6



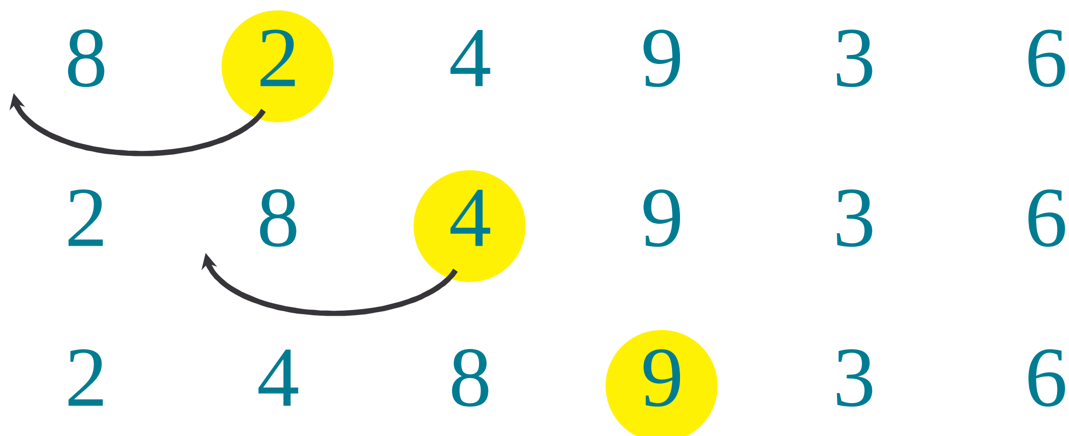
Example of insertion sort

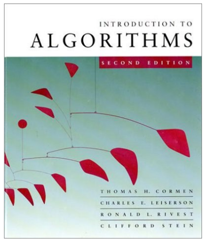
8 **2** 4 9 3 6

2 8 **4** 9 3 6



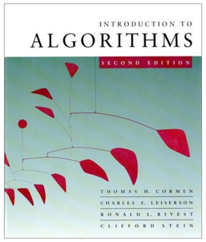
Example of insertion sort



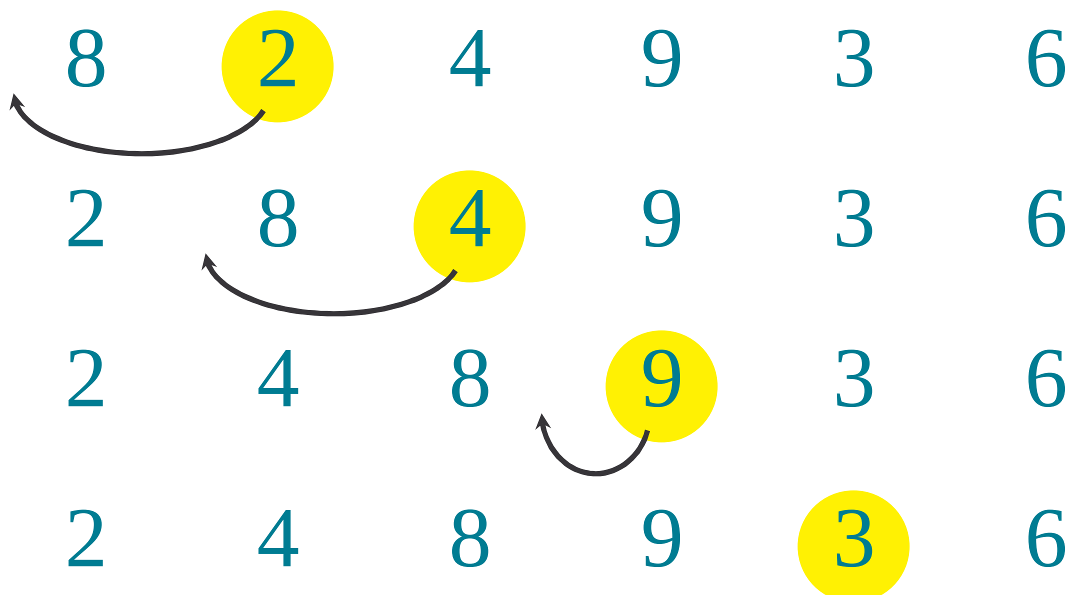


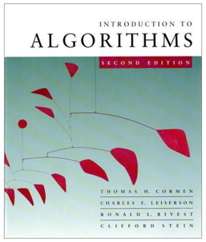
Example of insertion sort

8	2	4	9	3	6
2	8	4	9	3	6
2	4	8	9	3	6



Example of insertion sort





Example of insertion sort

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2	8	4	9	3	6
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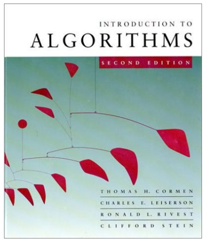
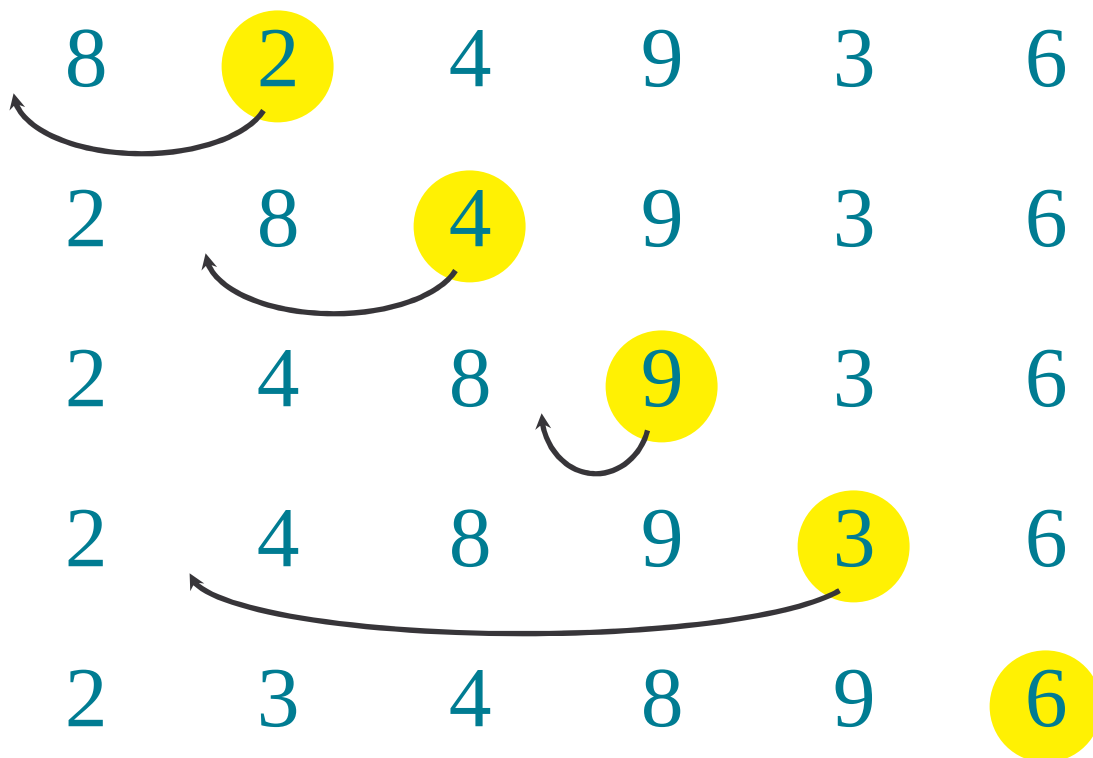
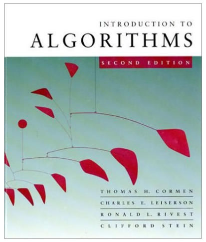


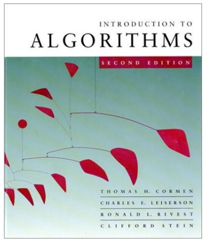
Example of insertion sort





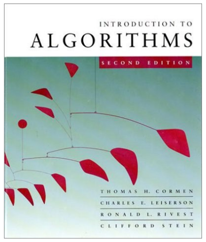
Example of insertion sort

8	2	4	9	3	6
2	8	4	9	3	6
2	4	8	9	3	6
2	4	8	9	3	6
2	3	4	8	9	6



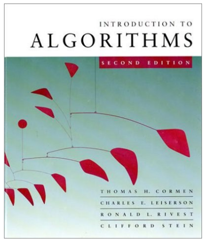
Example of insertion sort

8	**2**	4	9	3	6	
2	8	**4**	9	3	6	
2	4	8	**9**	3	6	
2	4	8	9	**3**	6	
2	3	4	8	9	**6**	
2	3	4	6	8	9	*done*



Running time

- The running time depends on the input: an already sorted sequence is easier to sort.
    - Parameterize the running time by the size of the input n
    - Seek upper bounds on the running time $T(n)$ for the input size n, because everybody likes a guarantee.



Kinds of analyses

Worst-case: (usually)

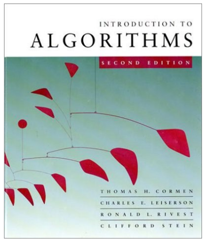
- $T(n)$ = maximum time of algorithm on any input of size n.

Average-case: (sometimes)

- $T(n)$ = expected time of algorithm over all inputs of size n.
- Need assumption of statistical distribution of inputs.

Best-case: (bogus)

- Cheat with a slow algorithm that works fast on *some* input.

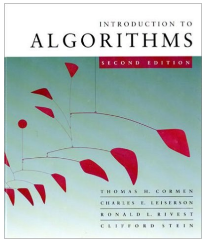


Machine-independent time

What is insertion sort's worst-case time?
- It depends on the speed of our computer:
  - relative speed (on the same machine),
  - absolute speed (on different machines).

BIG IDEA:
- Ignore machine-dependent constants.
- Look at ***growth*** of $T(n)$ as $n \to \infty$.

"Asymptotic Analysis"



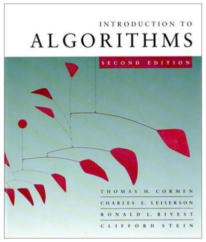
Θ-notation

Math:

$\Theta(g(n)) = \{ f(n)$: there exist positive constants c_1, c_2, and n_0 such that $0 \le c_1 g(n) \le f(n) \le c_2 g(n)$ for all $n \ge n_0 \}$

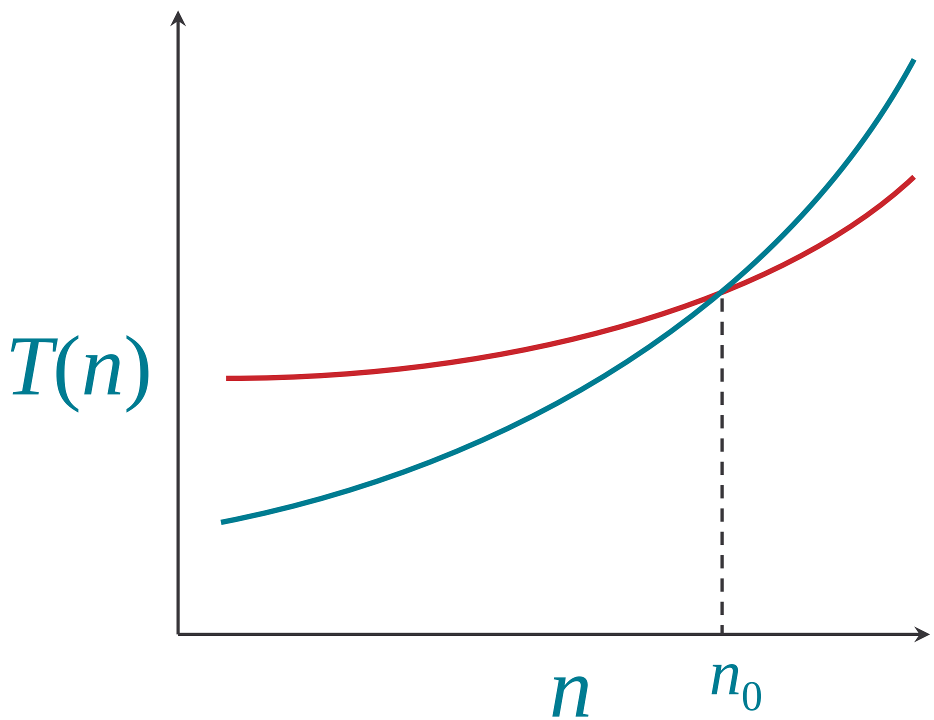
Engineering:

- Drop low-order terms; ignore leading constants.
- Example: $3n^3 + 90n^2 - 5n + 6046 = \Theta(n^3)$

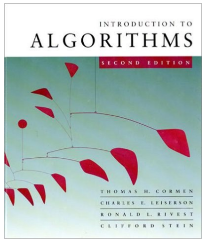


Asymptotic performance

When n gets large enough, a $\Theta(n^2)$ algorithm *always* beats a $\Theta(n^3)$ algorithm.



- We shouldn't ignore asymptotically slower algorithms, however.
- Real-world design situations often call for a careful balancing of engineering objectives.
- Asymptotic analysis is a useful tool to help to structure our thinking.



Insertion sort analysis

Worst case: Input reverse sorted.

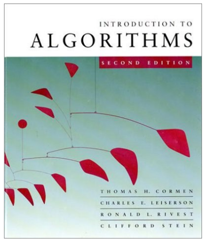
$$T(n) = \sum_{j=2}^{n} \Theta(j) = \Theta(n^2)$$ [arithmetic series]

Average case: All permutations equally likely.

$$T(n) = \sum_{j=2}^{n} \Theta(j/2) = \Theta(n^2)$$

Is insertion sort a fast sorting algorithm?

- Moderately so, for small n.
- Not at all, for large n.



Merge sort

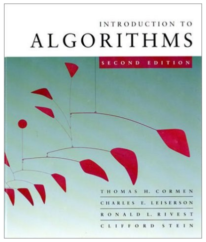
MERGE-SORT $A[1 \,..\, n]$

1. If $n = 1$, done.
2. Recursively sort $A[\,1 \,..\, \lceil n/2 \rceil\,]$ and $A[\,\lceil n/2 \rceil + 1 \,..\, n\,]$.
3. "***Merge***" the 2 sorted lists.

Key subroutine: MERGE

Merging two sorted arrays

20	12
13	11
7	9
2	1



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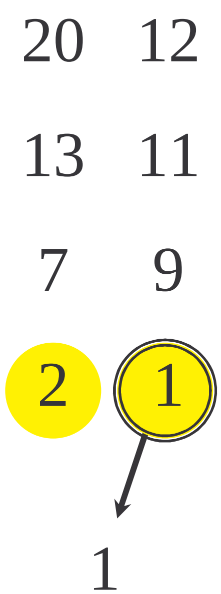
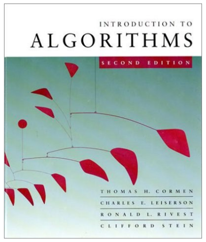
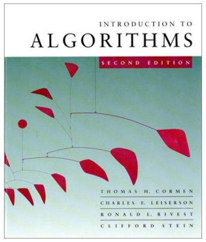




Merging two sorted arrays

20	12	‖	20	12
13	11	‖	13	11
7	9	‖	7	9
2	1	‖	2	

1



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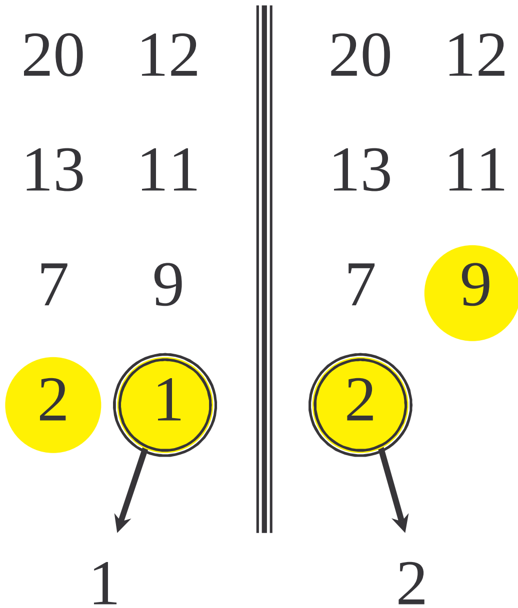

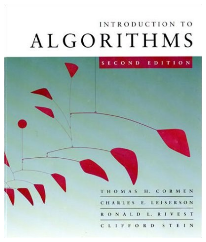

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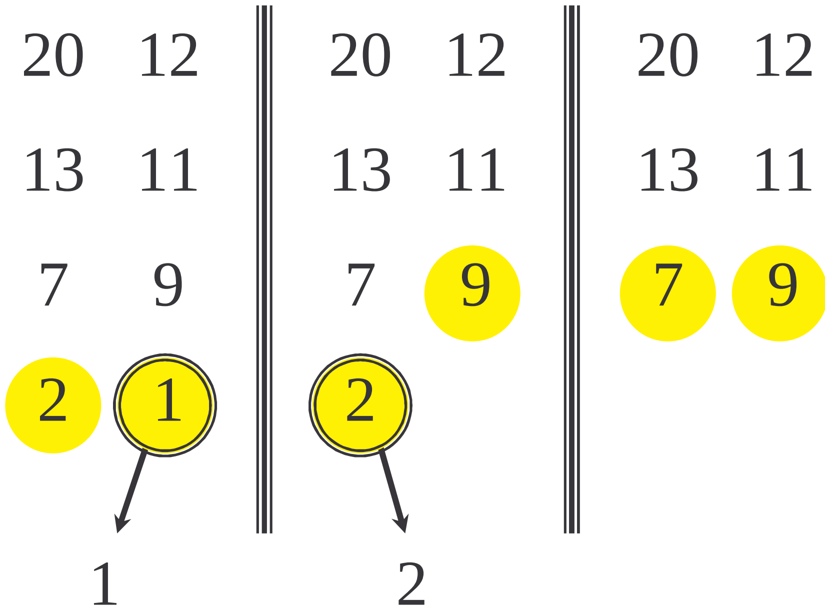

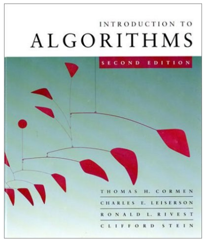

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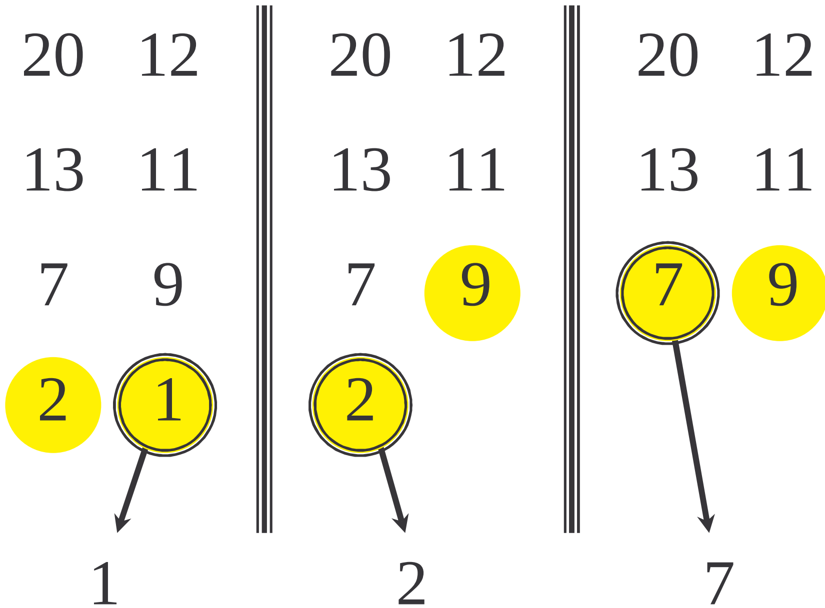

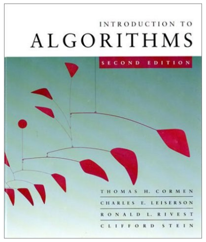

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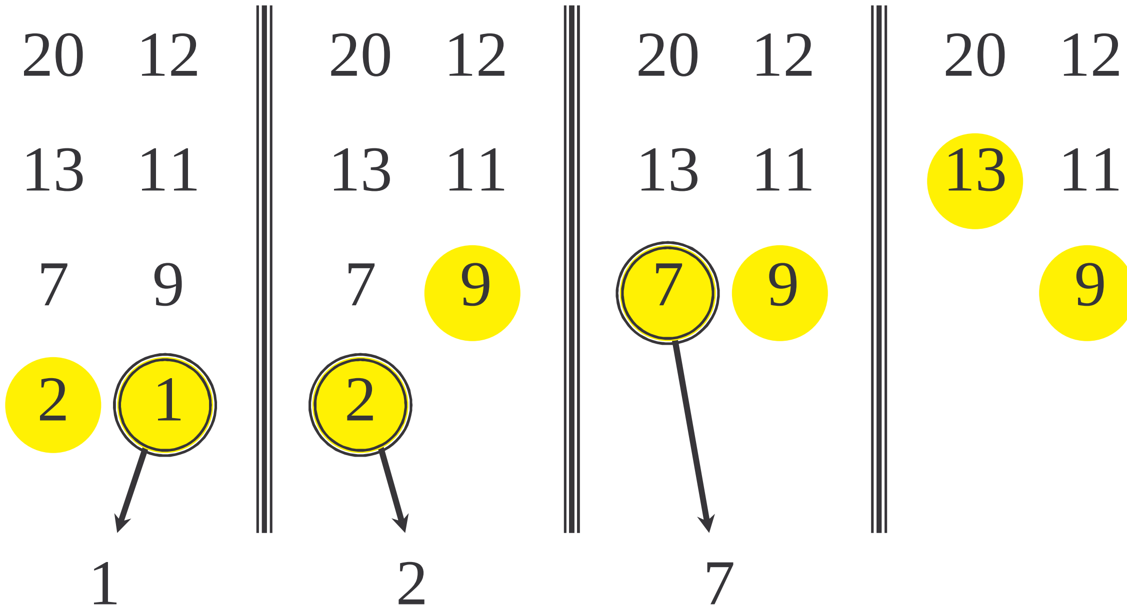

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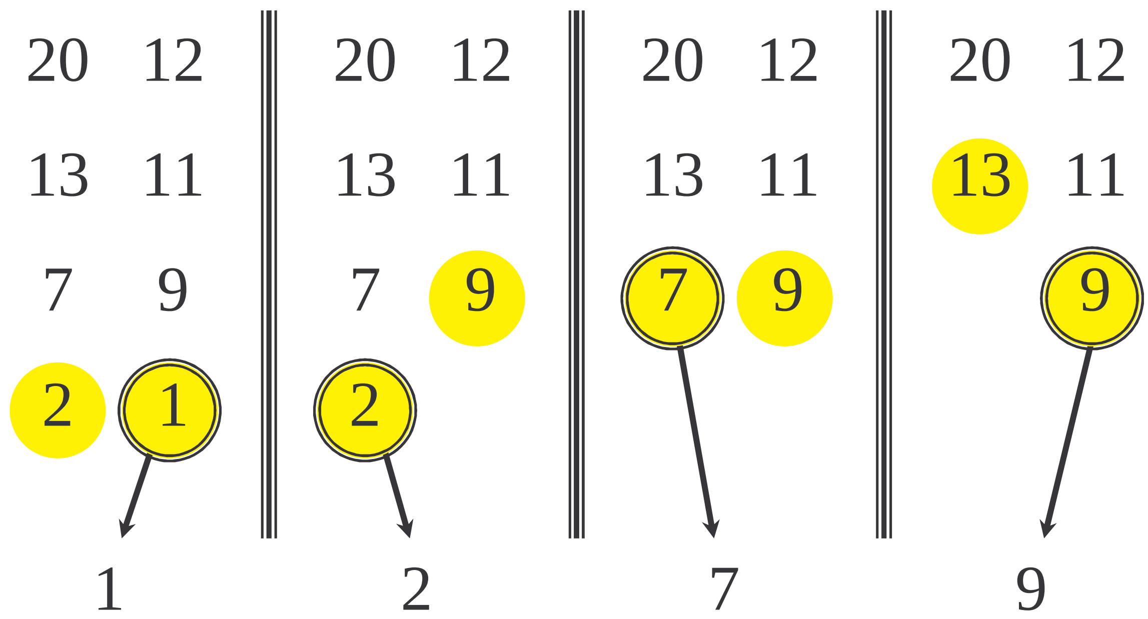

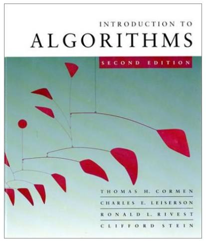

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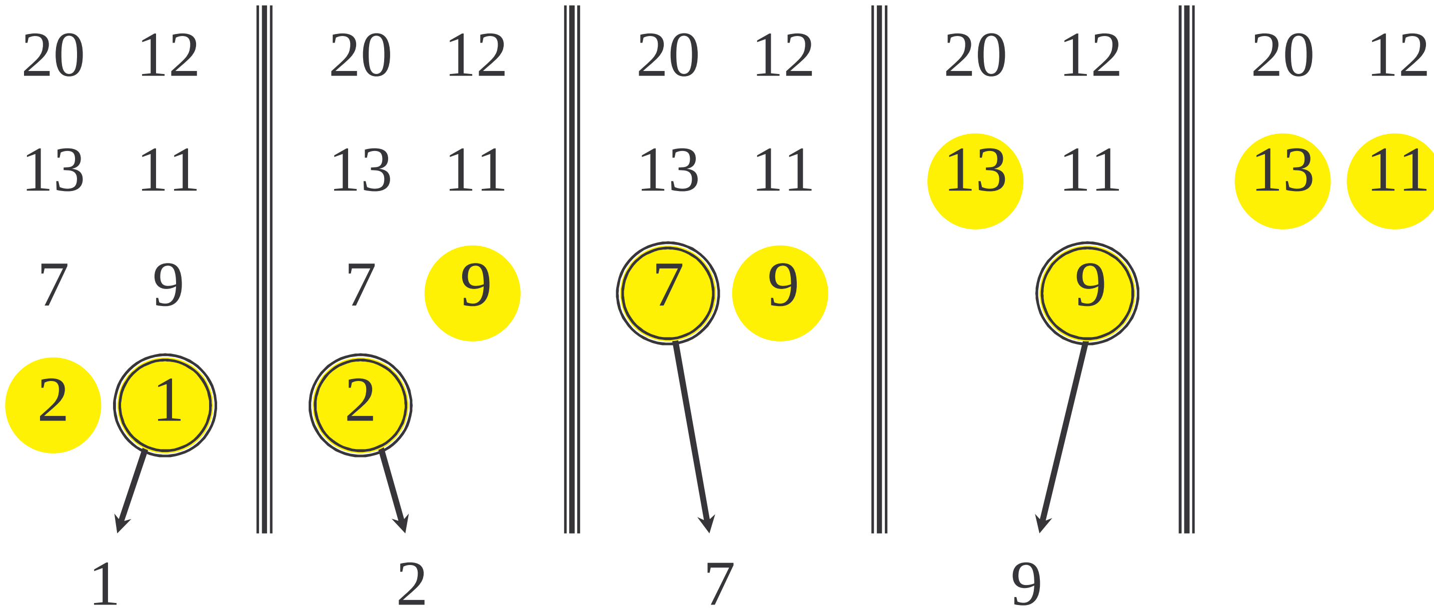

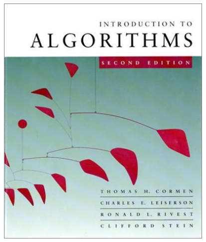

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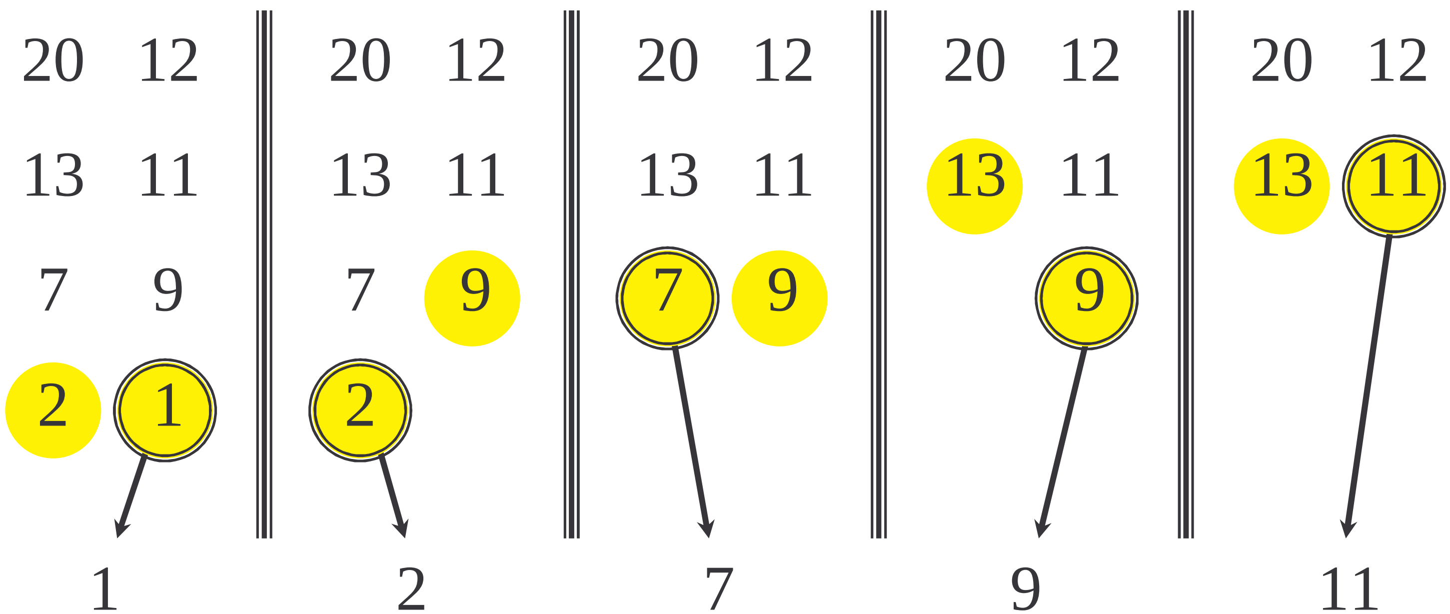

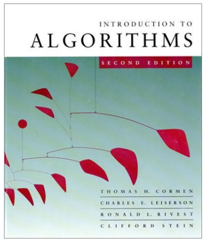

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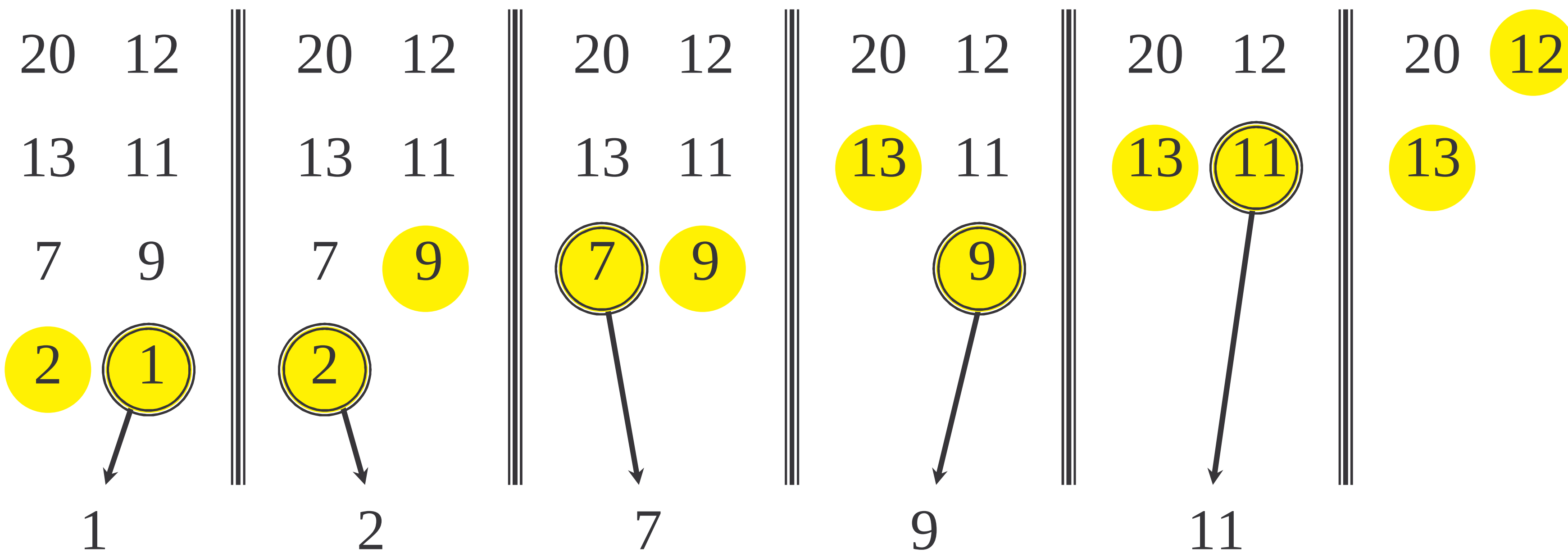

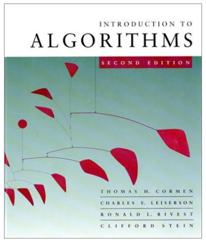

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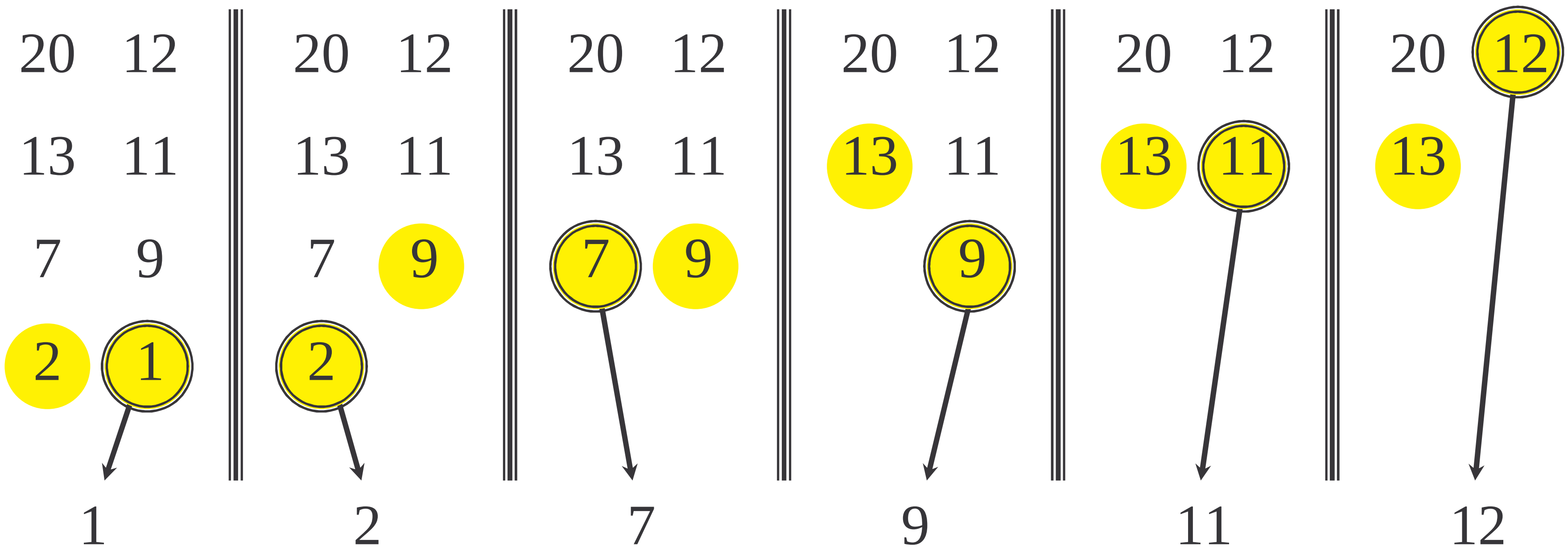
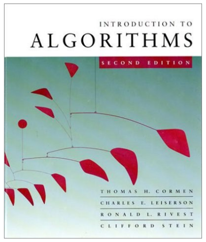
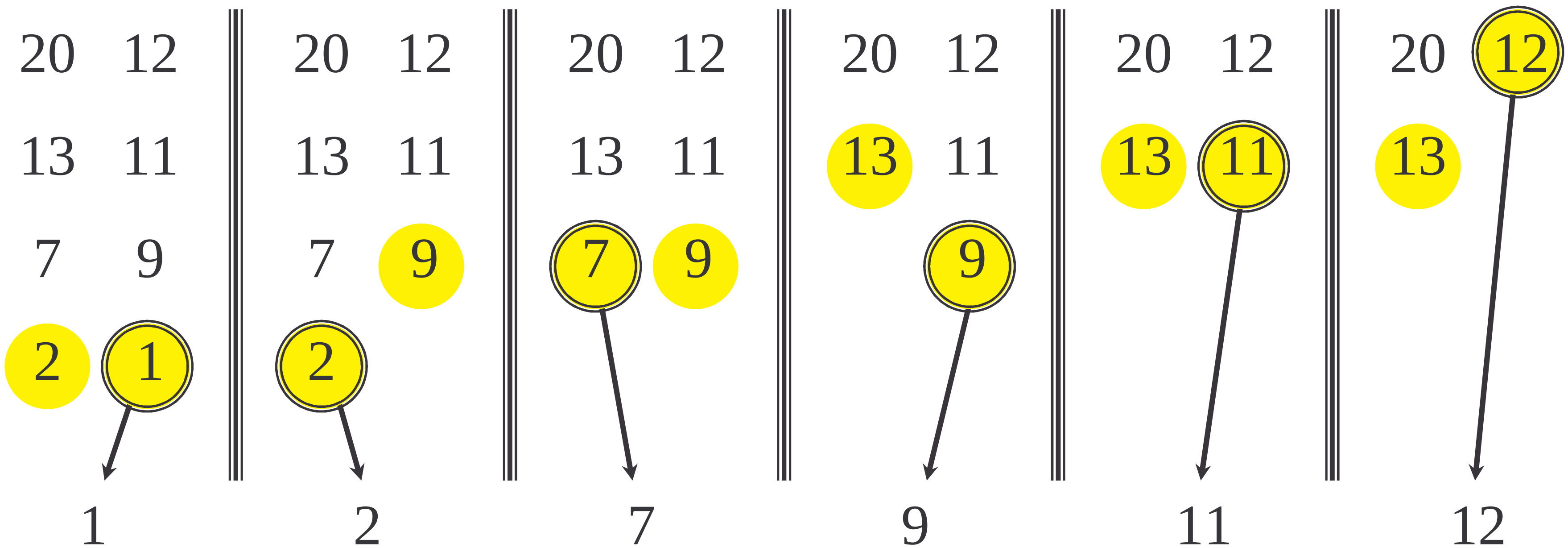
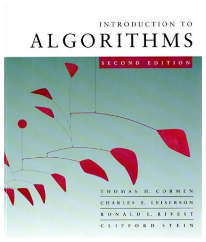




Merging two sorted arrays



Time = $\Theta(n)$ to merge a total of n elements (linear time).

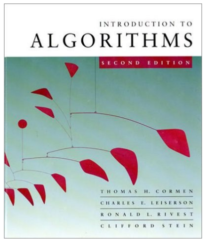


Analyzing merge sort

$T(n)$	**MERGE-SORT** $A[1 \,..\, n]$
$\Theta(1)$	1. If $n = 1$, done.
$2T(n/2)$	2. Recursively sort $A[\,1 \,..\, \lceil n/2 \rceil\,]$ and $A[\,\lceil n/2 \rceil + 1 \,..\, n\,]$.
$\Theta(n)$	3. ***"Merge"*** the 2 sorted lists

Abuse (arrows pointing to $\Theta(1)$ and $2T(n/2)$)

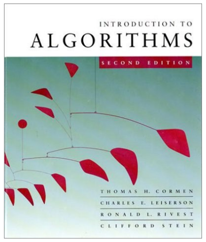
Sloppiness: Should be $T(\lceil n/2 \rceil) + T(\lfloor n/2 \rfloor)$, but it turns out not to matter asymptotically.



Recurrence for merge sort

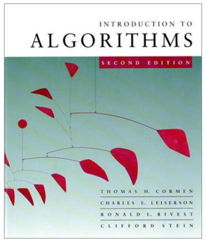
$$T(n) = \begin{cases} \Theta(1) \text{ if } n = 1; \\ 2T(n/2) + \Theta(n) \text{ if } n > 1. \end{cases}$$

- We shall usually omit stating the base case when $T(n) = \Theta(1)$ for sufficiently small n, but only when it has no effect on the asymptotic solution to the recurrence.
- CLRS and Lecture 2 provide several ways to find a good upper bound on $T(n)$.



Recursion tree

Solve $T(n) = 2T(n/2) + cn$, where $c > 0$ is constant.



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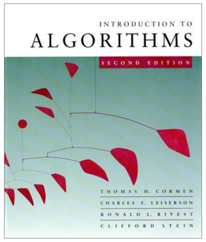
Solve $T(n) = 2T(n/2) + cn$, where $c > 0$ is constant.

$$T(n)$$



Recursion tree

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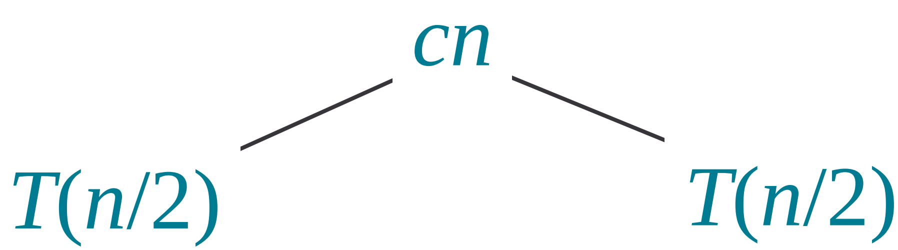

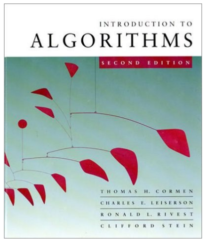

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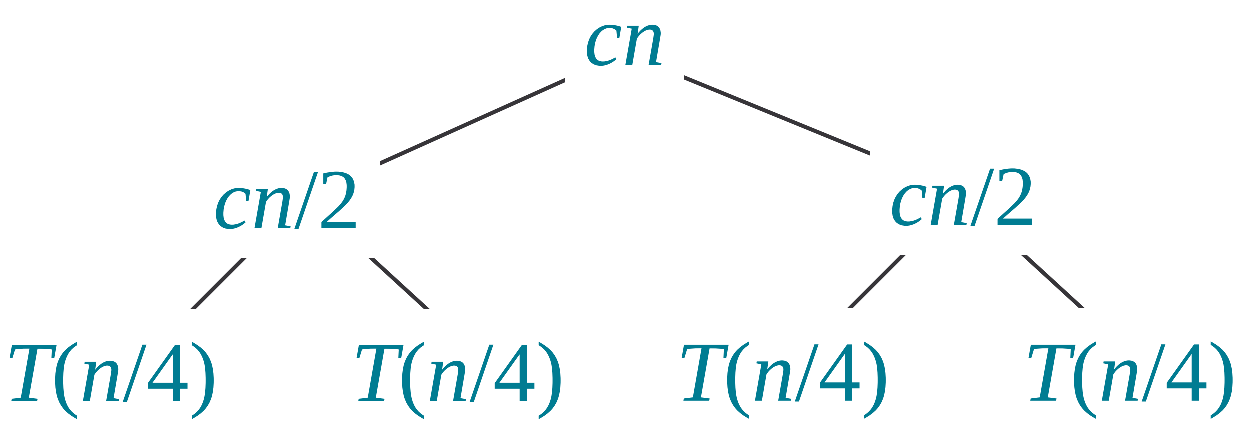

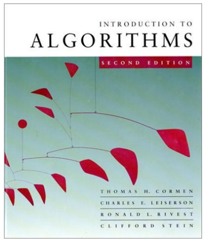

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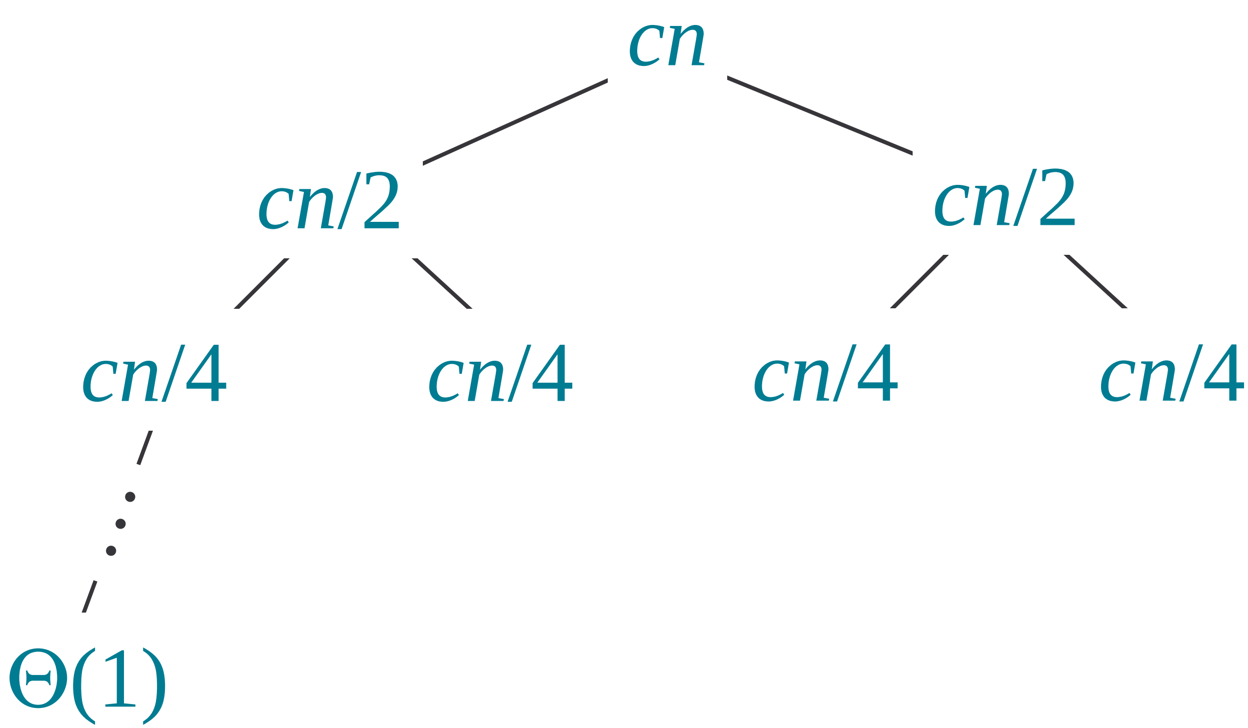

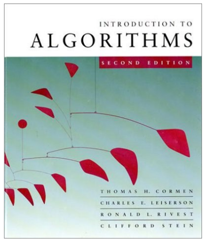

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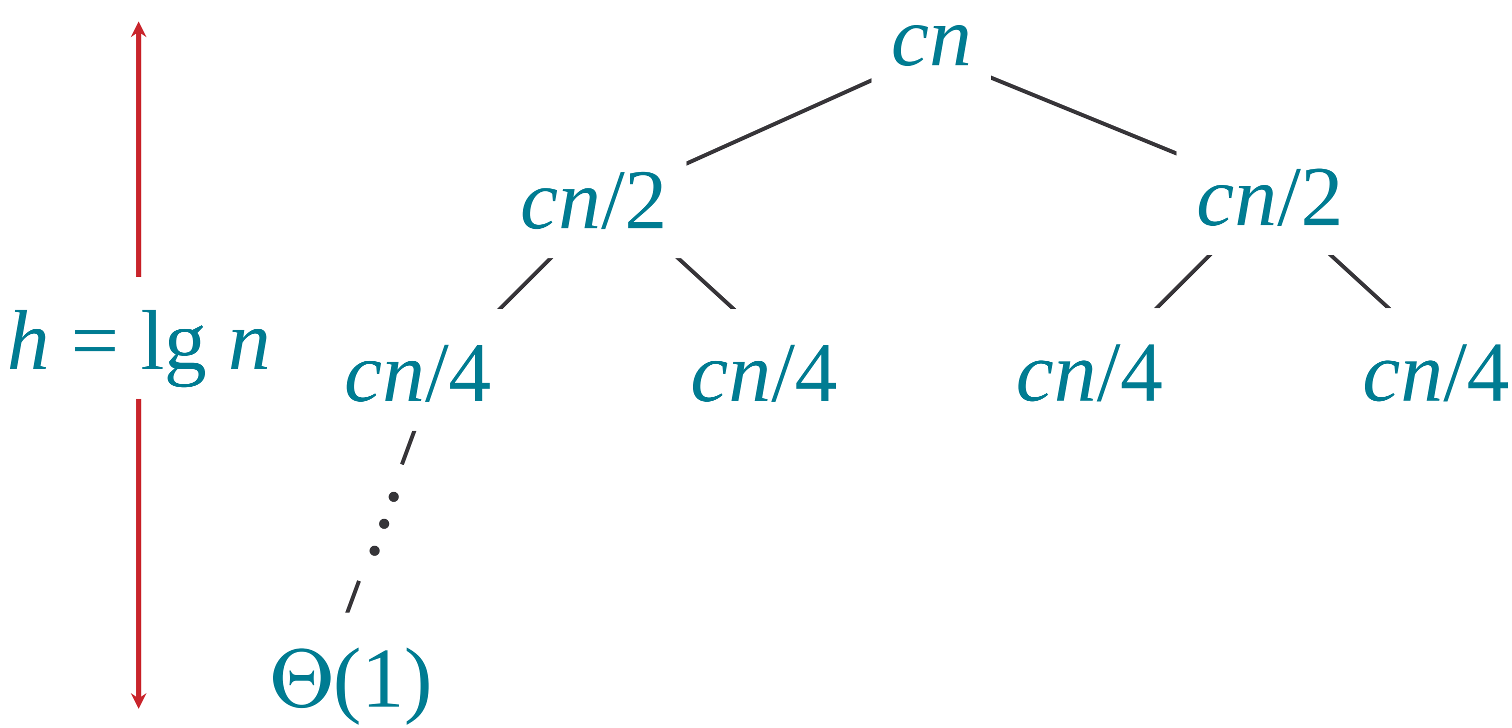

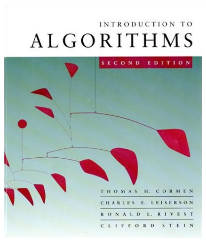

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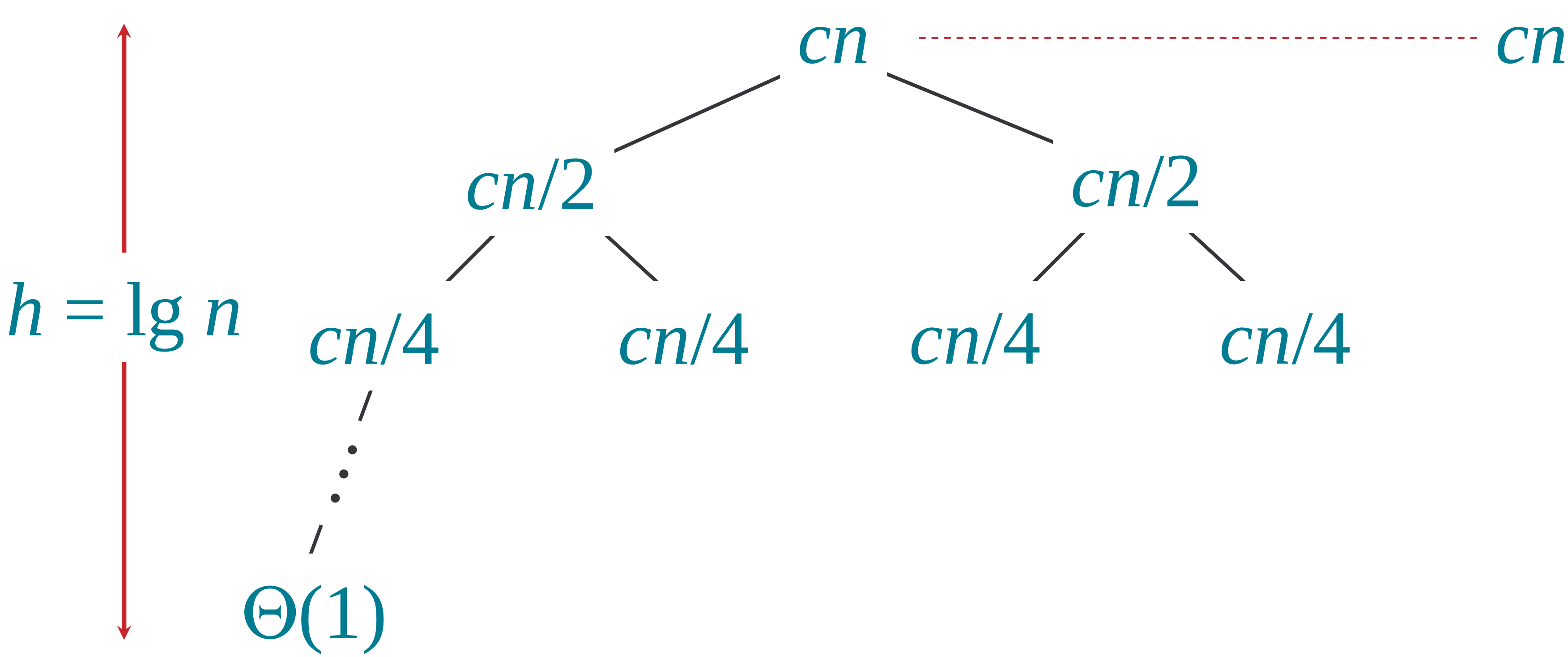

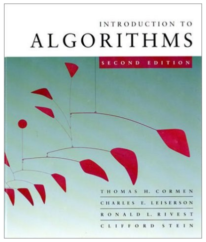

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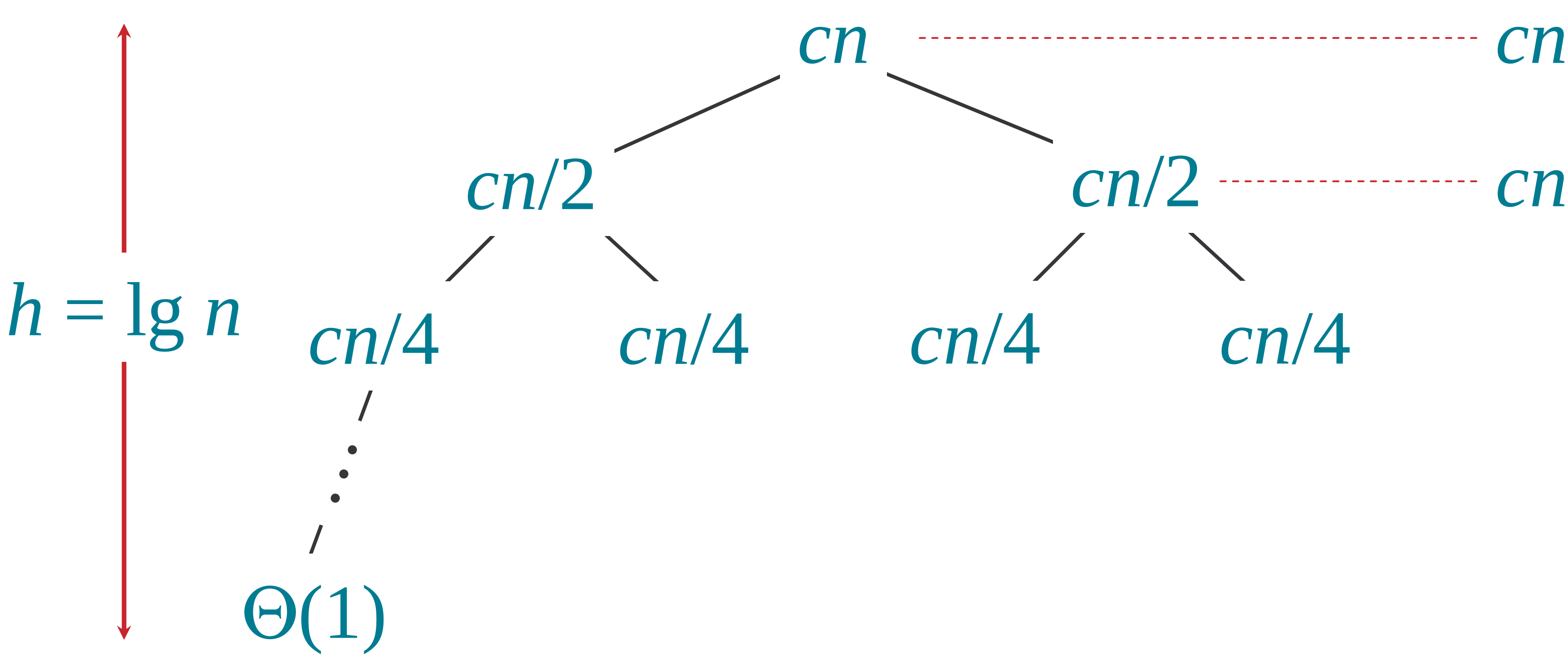

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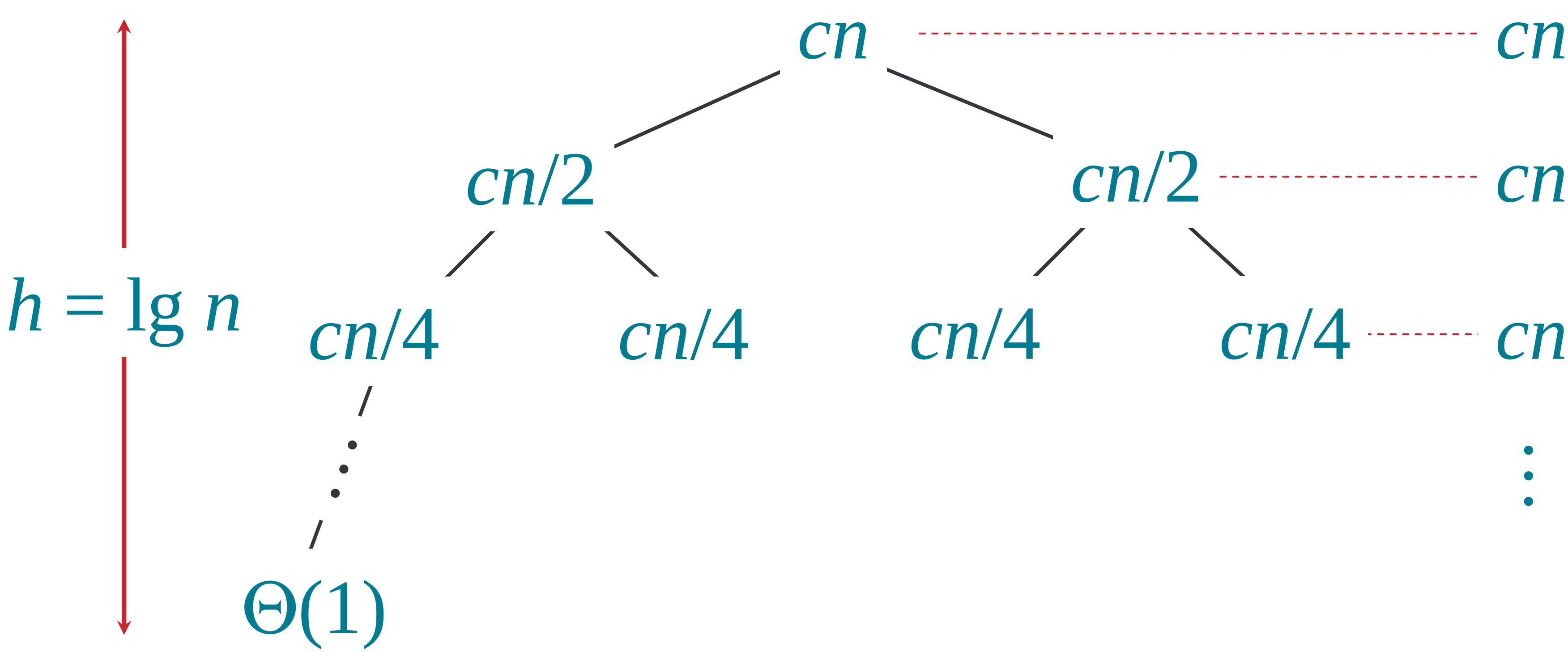

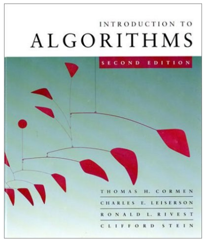

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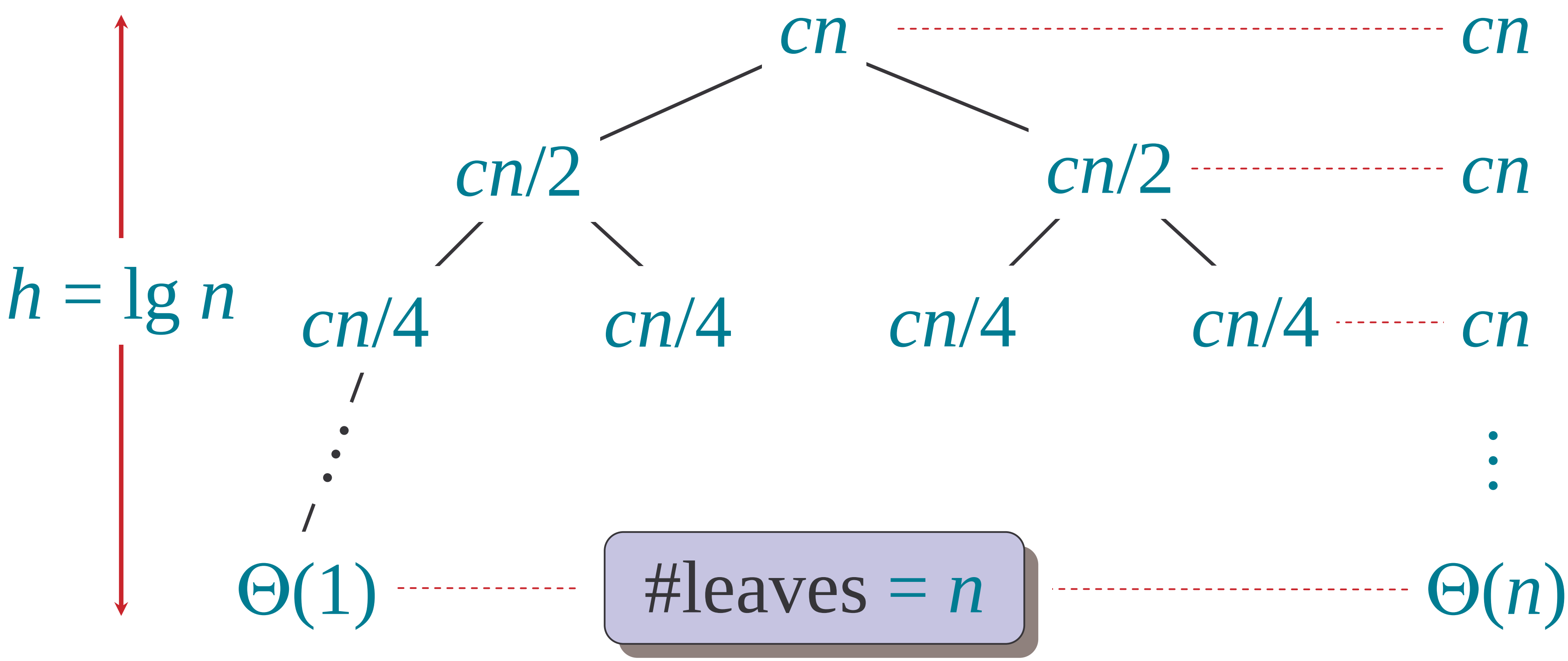

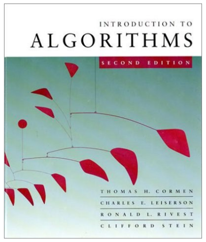

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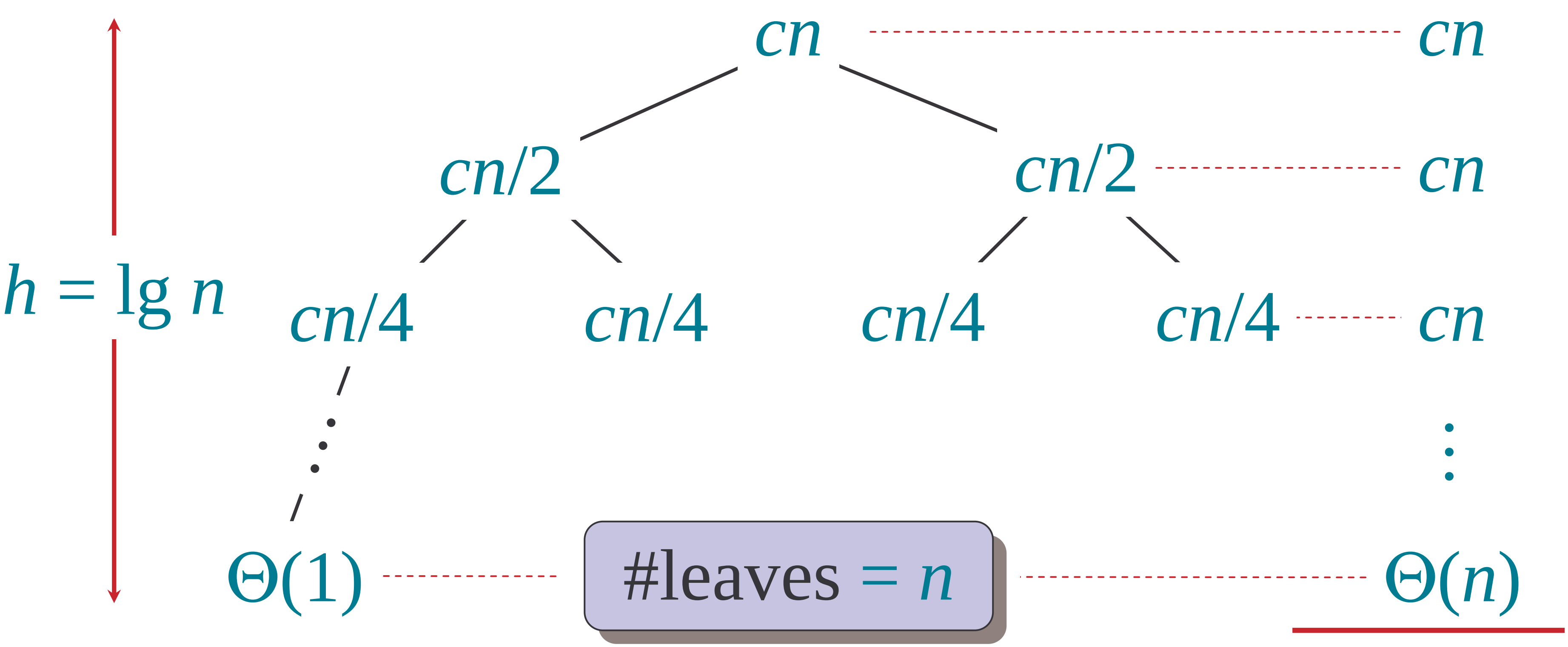
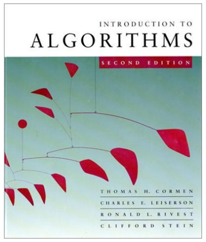


Total = $\Theta(n \lg n)$



Conclusions

- $\Theta(n \lg n)$ grows more slowly than $\Theta(n^2)$.
- Therefore, merge sort asymptotically beats insertion sort in the worst case.
- In practice, merge sort beats insertion sort for $n > 30$ or so.
- Go test it out for yourself!